



Microprogramming 21MX Computers

operating and reference manual



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This manual is a complete reference source for microprogramming the Hewlett-Packard 21MX Computer Series. With the facilities of the HP 12978A Writable Control Store the user can expand the already powerful capability of his 21MX Series Computer by adding custom-tailored instructions to the existing set of microprogrammed basic instructions.

The HP 12978A Writable Control Store is provided with two options. The 12978A option 001 provides software that operates in the DOS-III operating system. The 12978A option 002 provides software that operates in the Basic Control System. Refer to Section VI of this manual for a complete description of the options.

This manual is written for an individual who already has considerable experience as an assembly language programmer. HP 21MX Computer Series microprogramming is no more complex than normal assembly language programming on larger computers. Thus, with little more investment than learning a new assembly language, large computer capability can be had for small computer expense.

RELATED DOCUMENTATION

It is assumed that the microprogrammer has read the **HP 21MX Computer Series Reference Manual (HP 02108-90002)** and that he knows how to use his operating system, DOS-III (HP 24307B), or the Basic Control System (HP 20855A). These operating systems are described in the following publications:

HP 24307B DOS-III Disc Operating System (HP 24307-90006)
Basic Control System (HP 02116-9017)

During the process of writing, debugging, and using a microprogram, the user should also have access to and be familiar with the following additional publications.

The assembler used with the DOS-III-B system is described in:

Assembler Reference Manual (HP 24307-90014)

The assembler used with the Basic Control System is described in:

HP Assembler (HP 02116-2014)

The 21MX computer is described in:

HP 21MX Computer Series Operator's Manual (HP 02108-90004)
HP 21MX Computer Series Installation and Service Manual (HP 02108-90006)

The HP 12909B pROM Writer, which is used in conjunction with the six mask tapes produced by the Micro Debug Editor, and installing pROMs is described in:

HP 12909B pROM Writer Operating and Reference Manual (HP 12909-90009)
HP 12909B Programmable ROM Writer Interface Kit Installation and Service Manual (HP 12909-90005)

HOW TO USE THIS MANUAL

This manual is intended to be used in the following way:

- a. Read Section I for the introduction to user microprogramming.
- b. Study Section II to learn the structure of the system that is being controlled by microprogramming. Section II explains the relationship between the Control Section and the other sections of the computer.
- c. Become familiar with the reference material in Sections IV, V, and VI so that when the time comes to use the material, it may be found easily. These sections describe the microprogramming language, the Micro-assembler, the Micro Debug Editor, and the 12978A Writable Control Store.
- d. Study Section III to learn how to write a microprogram.

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INTRODUCTION TO USER MICROPROGRAMMING

SECTION

I

The Control Section of a computer contains circuitry which decodes each machine instruction and then executes the required sequence of operations. Machine instructions can be decoded and executed by either a conventional Control Section or a microprogrammed Control Section.

1-1. CONVENTIONAL CONTROL SECTION

In a conventional computer Control Section, specific hardware is dedicated to each function performed by the instruction set. The major advantage of this specially designed hardware is speed for the instruction set. The major disadvantage is the loss of flexibility for special applications or for enhancements. Changes and additions to hardware components are required to implement changes and additions to existing capabilities.

This is no problem for a conventional computer if no new machine instructions are required. The hardware has been designed to minimize timing for the instruction set. Rarely however, does a computer manufacturer produce an instruction set that fully meets the requirements of most potential users. Hence, the manufacturer must either focus his attention on one group of users (specialize) or widen his scope and generalize the hardware design to meet the needs of a number of user groups. In the latter case, the user must modify his discipline to some extent to meet the limitations of his hardware.

1-2. MICROPROGRAMMED CONTROL SECTION

In the microprogrammed computer, all distinct logical functions are separated from the sequence in which those functions are performed. Hardware redundancy is thus reduced. The logical functions are defined by a bit pattern or micro-instruction held in Control Store. Each machine instruction in Main Memory is performed by a sequence of micro-instructions in Control Store that defines the logical functions to be performed. This sequence of micro-instructions is called a microprogram and is often referred to as firmware, because it lies somewhere between hardware and software in origin and permanence.

Software can execute much faster with the application of microprogramming. This speed is achieved by two factors: the ratio of Control Store speed over Main Memory and the relative flexibility of a micro-instruction over normal machine instructions. The HP 21MX Control Store, where micro-instructions reside, cycles more than twice as fast as Main Memory, where normal machine instructions reside. Control Store words are 24 bits whereas Main Memory words are 16 bits. In addition, micro-instructions have access to many internal registers and logical functions that Main Memory programs cannot use.

For example, the 21MX floating point software subroutines were identified as being very time consuming. They were then microcoded by a Hewlett-Packard microprogrammer and made available in Read Only Memory to users. Implementation of the floating point firmware requires no change to user programs. The microprogrammed floating point instructions run about 20 times faster than the corresponding software subroutines.

As in the floating point microprogram, the user can study his software, determine the most time consuming functions performed, and then microprogram those functions, that is, execute them in Control Store using a single Main Memory instruction instead of a sequence of Main Memory instructions. Any software that uses those microprogrammed functions will execute at a higher speed.

1-3. LIMITATIONS OF HP 21MX MICROPROGRAMMING

The user should be aware of the following limitations imposed by HP 21MX microprogramming:

- a. Since the origin of a microprogram is specified during micro-assembly, HP 21MX microprograms are not relocatable.
- b. Since there is only one register available to the microprogrammer to save subroutine return addresses, the HP 21MX design allows for no more than one logical microprogram subroutine level. This limitation can be circumvented by using other registers or Main Memory to simulate subroutine nesting.
- c. The microprocessor cannot be interrupted. If the microprogram execution time exceeds the interval between interrupts (75 μ s. is the maximum interval

allowed by Hewlett-Packard instruction set microprograms), the microprogram must test for pending interrupts or they can be lost. When a pending interrupt is detected, the microprogram must yield control to the interrupt handler. For a discussion of microprogram interrupt handling, refer to sections 3-32 and 3-33 in this manual.

1-4. SUMMARY

The advantages of microprogrammed control are:

- a. The user can use a fully-supported general purpose computer to aid in the generation and debugging of extensions to the computer's own instruction set.
- b. The user can speed up the overall execution time of his software by microprogramming its most time consuming or repetitious routines.
- c. The user can implement enhancements of the instruction set and special purpose processors produced by the manufacturer with little impact on his existing software.

THE MICROPROGRAMMABLE COMPUTER

SECTION

II

To successfully implement microprograms, the assembly language programmer must learn more about the computer. This section of the manual is the introduction to the structure of the computer. A functional block diagram of the microprogrammable machine is provided in Appendix D. This diagram describes what paths data can follow. Control commands or micro-instructions spell out what paths the data does follow and what modifications and tests are performed in the process.

Functionally, a computer consists of four major sections:

- Control
- Main Memory
- Input and Output
- Arithmetic and Logic

and data are stored in the Main Memory. Parameters, status, commands, and processor results (data) are exchanged with external devices such as teleprinters, magnetic tape units, and line printers via the Input and Output (I/O) section. Add, subtract, and other mathematical functions and shift, "or", "and", and other logical functions are performed in the Arithmetic and Logic section. The Front Panel registers and switches provide direct operator communication.

Each section executes under the direction of the Control Section by means of a microprogram. The Control Section reads the user's program stored in Main Memory and directs the appropriate hardware in each of the other sections.

Figure 2-1 shows the four major sections of the computer.

2-1. RELATIONSHIP BETWEEN SECTIONS

These four sections and the Front Panel are interconnected by a network of signal paths. Data processing programs

2-2. CONTROL SECTION

To write a microprogram an understanding of the Control Section is required. The Control Section takes an instruction from Main Memory and stores it into the Instruction

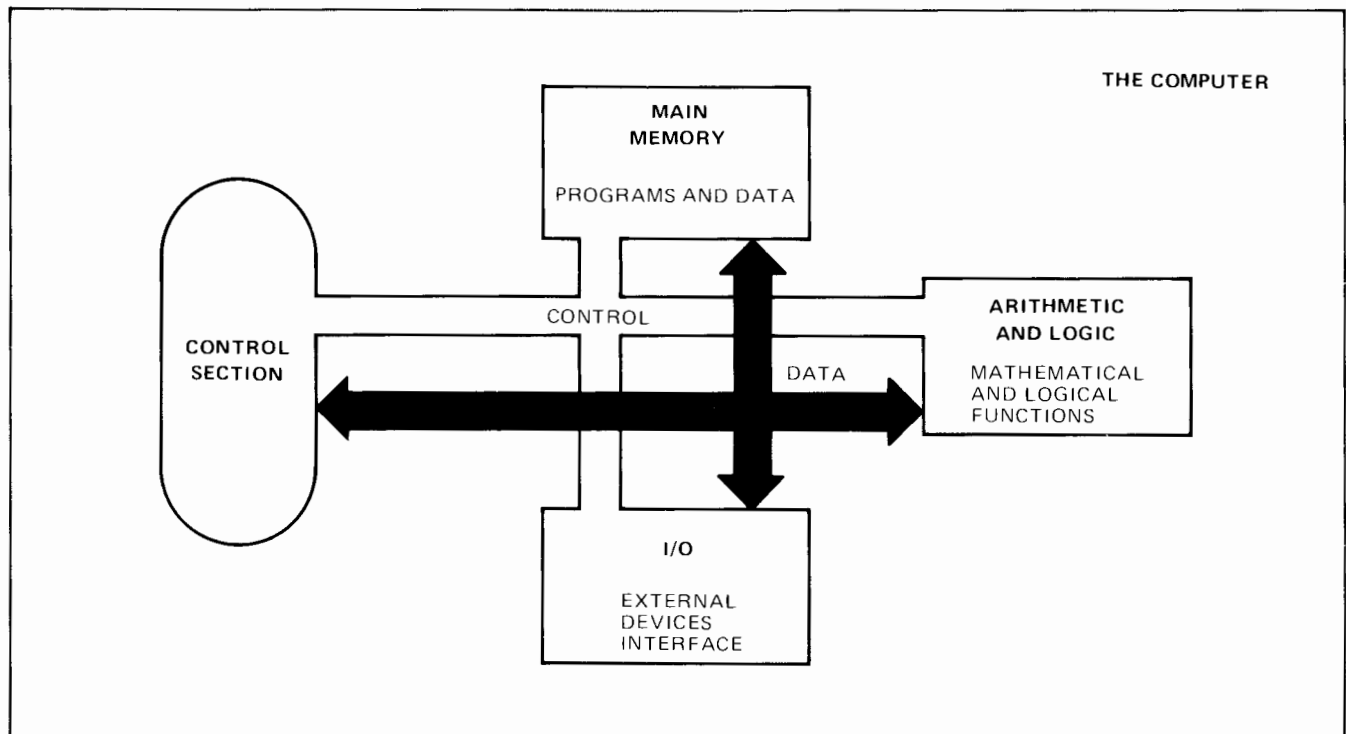


Figure 2-1. Four Major Computer Sections

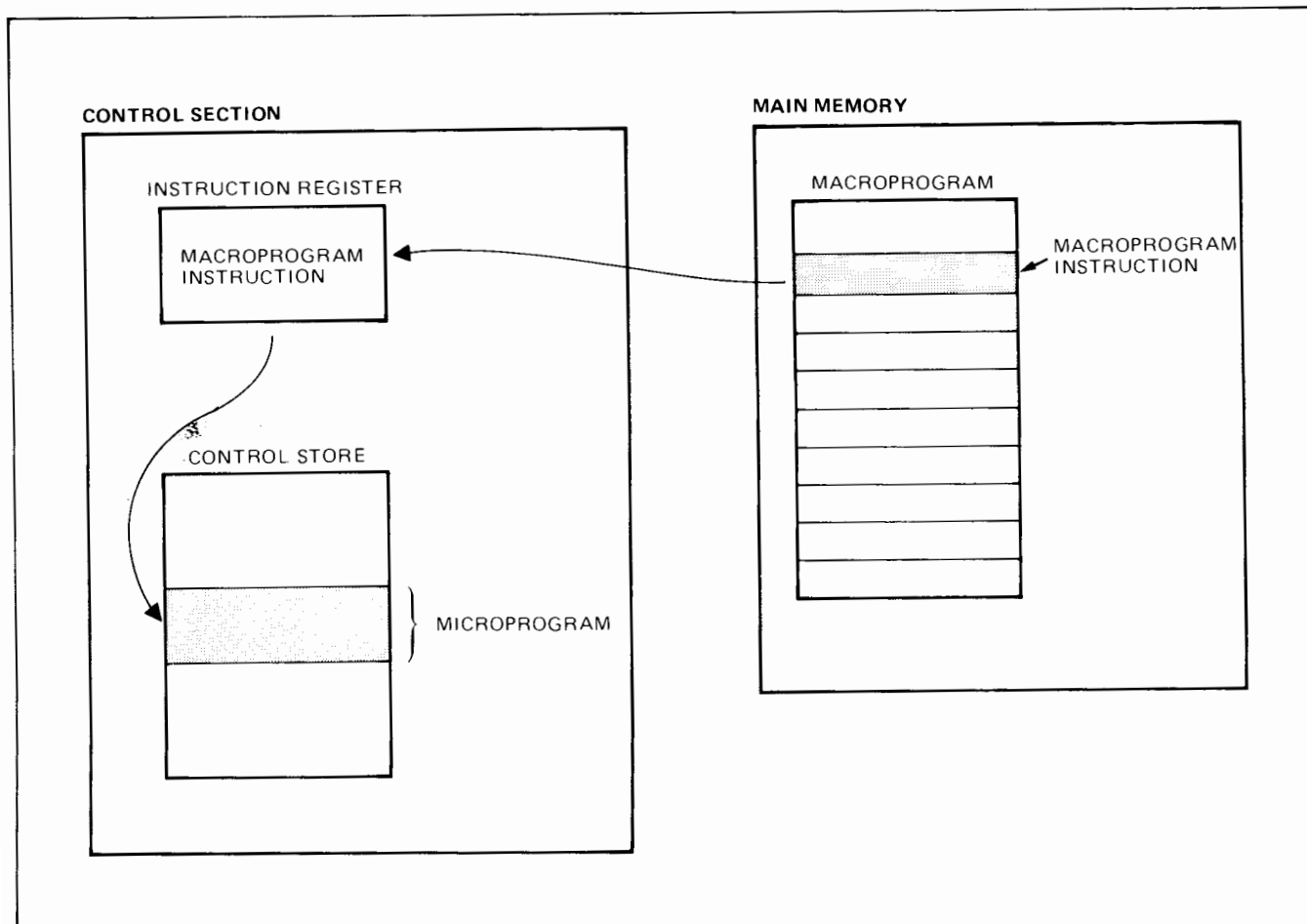


Figure 2-2. A Microprogram Implements One Macroprogram Instruction

Register (IR), as shown in figure 2-2. An appropriate microprogram is executed whose Control Store entry point address is determined by the IR. View, then, each program instruction in Main Memory as a jump to a microprogrammed routine, which resides in Control Store.

The storage area for these microprograms is Control Store which may be either a Read Only Memory (ROM) or Writable Control Store (WCS). In this manual, to distinguish programs in Main Memory from microprograms in ROM, Main Memory programs are called macroprograms. We refer to a Control Section that executes microprograms from ROM, as a Control Processor.

2-3. THE CONTROL PROCESSOR

A microprogram in the Control Processor is in command of the computer at all times. A microprogram which is part of the basic 21MX instruction set microprogram takes program instructions from Main Memory and stores them into the Instruction Register. The upper eight bits of the Instruction Register determine the microprogram address within one of the following instruction groups:

- Basic Instruction Set
- Extended Instruction Group
- Floating Point Instruction Group
- User Microprogram Group

Since the user is mainly interested in writing and executing his own microprograms, he can regard the Basic Instruction Set microprogram as a supervisor microprogram that determines when a user microprogram is called and then passes control to the user microprogram.

When the Instruction Register holds an octal 101rrr or 105rrr (see table 3-1 for possible values of rrr), a branch is made to the user microprogram area of Control Store.

When a microprogram has run to completion, it returns to location 0 in Control Store to take the next instruction from Main Memory and store it into the Instruction Register.

2-4. THE MICROPROGRAMMER'S ROADMAP

Appendix D holds the fundamental diagram of the computer required by the microprogrammer. This functional

block diagram is the “roadmap” that is used to determine possible data paths and to determine where logical decisions can be made. This diagram can be unfolded and referred to while reading other parts of the manual. Note that the four sections of the computer, illustrated in figure 2-1, are shown in more detail in the functional block diagram.

To read the functional block diagram, begin with a 101rrr or 105rrr instruction in the Instruction Register. The rrr specifies the octal Control Store entry point address according to the description in section 3-24, Mapping to a Module Address. This address is moved into the ROM Address Register (RAR). With a first address specified, the user microprogram begins execution. The contents of the Control Store location given in the ROM Address Register are moved into the ROM Instruction Register (RIR). The ROM Instruction Register now holds a 24 bit micro-instruction. The micro-instruction is decoded and the specified control functions are executed.

Successive micro-instruction addresses are determined in the following way. The ROM Address Register is incremented at the start of execution of each micro-instruction. When a jump is executed, the ROM Address Register is loaded with the jump target address. When a jump to subroutine is executed, the ROM Address Register is stored into the SAVE Register (save return address) and the jump target address is stored into the ROM Address Register. When a return from subroutine is executed (RTN), the SAVE Register contents are transferred into the ROM Address Register and the SAVE Register is cleared. Thus at the completion of execution of each micro-instruction, the ROM Address Register holds the address of the next micro-instruction.

2-5. DATA PATHS

The central data transfer path is the S-bus. The contents of all registers except the following can be directed onto the S-bus: L-register, RAR, SAVE Register, Extend Register, and the Overflow Register. The following registers can receive data from the S-bus:

- M-register
- T-register
- L-register
- Counter Register
- Display Register
- Display Indicator
- Instruction Register

The T-bus receives data only from the Rotate/Shifter (R/S) but can pass data to these registers:

- A-register
- B-register
- Scratch Pad Registers (S1 through S12)

- X-register
- Y-register
- P-register
- S-register (Front Panel Switch Register)

The I/O-bus serves to transfer data to and from external devices under programmed control.

Note in Appendix D, the functional block diagram, that the arrows are significant. For example, the B-register contents can be sent to the S-bus and thence to the M-register. However, the contents of the B-register cannot be sent to S12 (Scratch Pad 12) without passing through the ALU.



2-6. MAIN MEMORY

The M-register is a 15 bit register which holds memory addresses for reading from or writing into Main Memory. When storing from the M-register, bit 15 is clear (0). The T-register or Transfer register holds the data being transferred to or from memory. Contents of both these registers are transferred to and from the S-bus. Four loader ROMs, selectable by Instruction Register bits 15 and 14, each can contain a 64 word Main Memory program which may be loaded into Main Memory and used to load Main Memory from a peripheral device or to perform any other function desired by the user.

Two flags are associated with memory: the A-register Addressable Flag (AAF) and the B-register Addressable Flag (BAF). These flags are required to allow the A- and B-registers to be addressed as locations 0 and 1, respectively, of Main Memory.

2-7. I/O SECTION

The Central Interrupt Register (CIR) is a 6 bit register associated with the I/O interrupt circuitry. It is loaded with the Select Code of the interrupting device under program control and passed to the S-bus. Whenever the Central Interrupt Register is loaded, an Interrupt Acknowledge (IAK) signal is issued to the I/O device.

The I/O-bus transfers data to and from external devices.

Two flags are associated with I/O: the Interrupt Pending flag and the I/O Skip Condition Met (Main Memory instructions SFS and SFC) flag.

The Interrupt Enable Register is used to disable or enable the recognition of all interrupts, except Memory Protect, Parity, and Power Fail interrupts.

2-8. ARITHMETIC AND LOGIC SECTION

This section consists of the Arithmetic and Logic Unit (ALU), the Rotate/Shifter (R/S), 22 registers and six flags.

The ALU and R/S are the only units that execute functional modifications on the data. The ALU receives input from the S-bus and from the L-register (Latch Register). Output from the ALU goes to the R/S which places its output on the T-bus.

Output from the ALU and R/S can be stored in one of the following registers via the T-bus:

- A-register
- B-register
- Scratch Pad Registers (S1 through S12)
- X-register
- Y-register
- P-register
- S-register

Remember that the P-register holds the macroprogram (Main Memory) address. The P-register must be under control of the microprogram which must insure that it contains the proper address after the microprogram is complete. When the microprogram is complete, the resulting P-register value is the address of the next macro-instruction to be executed. Note that the Basic Instruction Set fetch routine (at Control Store address 0) automatically increments the P-register after the macro-instruction is fetched. Thus for one word user macro-instruction function codes, no further incrementing of the P-register is necessary in the user microprogram.

The S-register is reserved for internal storage of the Front Panel switch register. Note that all of these registers can also be sent along the S-bus for storage into memory, passage to an external device, or input to the ALU.

The Extend Register is a one bit register used in shift operations to link the A- and B-register or to indicate a "carry" arithmetic result out of the A- or B-registers. The Overflow is a one-bit register used to indicate an arithmetic overflow from the ALU. (See **21MX Computer Series Reference Manual**, where Overflow and Extend Register arithmetic results are fully explained.) These two registers can also be used as flags.

The 8 bit Counter Register, which passes to and from the S-bus, is used for repeat instructions, for Loader ROM addressing, and other general purposes, such as looping in a microprogram.

There are six flags dedicated to the Arithmetic and Logic Section. The CPU Flag is a general purpose flag. Four others signal output results from the ALU and one indicates the last T-bus value. ALU Ones is set when all ones are output from the ALU. ALU Carry Out is set when an ALU function produces a "carry" out of bit 15. ALU Bit 0 and ALU Bit 15 flags represent the last value of the specified bit in the ALU output. T-bus Zero flag is set if all bits of the T-bus are zero.

2-9. FRONT PANEL

Two registers and two flags are associated with the Front Panel Section. The Display Register holds the contents of the register A, B, M, T, P, or S, indicated by the Display Indicator. The Display Register and the Display Indicator are displayed on the Front Panel, as illustrated in figure 2-3.

The Run Mode flag indicates that the computer is in a Run or Halt condition. The Run Enable flag indicates whether the four position key-operated switch on the front panel is in Lock or Operate mode.

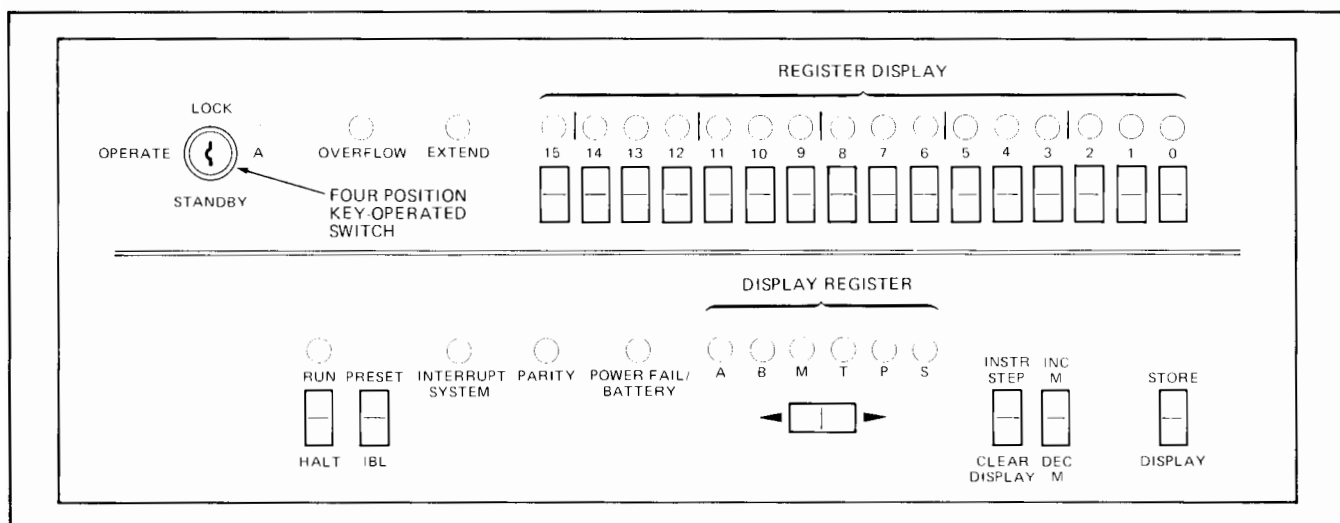


Figure 2-3. Front Panel Displays and Switches

This section introduces the basics of writing and debugging a microprogram in the micro-assembly language.

An assembly language programmer who codes programs for Main Memory may shun microprogramming because he regards it as too complex, mysterious, and the exclusive field of the computer designer.

However, Hewlett-Packard has especially designed the HP 21MX series computers to enable assembly language programmers to quickly get to the microcode level of computer logic so that they can attack the most time-consuming and least efficient parts of the software. Execution times can be cut with the proper application of microcode.

3-1. AN EXAMPLE

Figure 3-1 illustrates a segment of a microprogram. Ten micro-instructions are shown coded on the 21MX Micro-coding Form. The second micro-instruction shaded in figure 3-1 consists of the following four codes:

COV	PASS	M	P
-----	------	---	---

Each of the four codes are called micro-orders:

- P** takes the 16 bits in the P-register and puts them onto the S-bus.
- M** stores the 16 bits on the S-bus into the M-register (bit 15 of M-register is always 0).
- PASS** passes the 16 bits on the S-bus through the ALU without modification.
- COV** clears the Overflow Register.

Note in figure 3-1 that the various micro-orders of the micro-instruction begin in certain columns of the micro-coding form. These columns define the location of fields of the micro-instruction and each field holds a certain type of micro-order. In the case of the example micro-instruction, field 3 holds the special operation COV, field 4 holds the ALU operation PASS, field 5 holds the store operation M, and field 6 holds the data source P, that is, the data placed on the S-bus.

Section IV of this manual gives a full explanation of micro-instruction formats and micro-orders.

3-2. COMPARISON BETWEEN ASSEMBLY AND MICRO-ASSEMBLY LANGUAGE PROGRAMMING

The assembly language programmer is already familiar with the basic concepts of programming: the instruction, data source, data destination, data modification, data test, and branch. These concepts hold in microprogramming.

3-3. THE INSTRUCTION

The normal macro-instruction in Main Memory is 16 bits long. Most macro-instructions consist of one operation command (for example Add to A-register) and a data source or destination (for example Memory Location 1237). Thus there are usually two orders in a macro-instruction [Add to A-register] [Memory location 1237]. This is coded in Assembly Language as ADA VALU, where VALU is the label of memory location 1237.

The micro-instruction in Control Store is 24 bits long, which allows more control and flexibility to be coded into each instruction. A micro-instruction consists of up to five orders called micro-orders. Section 3-1 gives an example of four micro-orders coded into a micro-instruction.

There are four micro-instruction formats. Each format defines a micro-instruction Word Type (Word Type 1, Word Type 2, etc.) and determines a set of micro-orders which may be coded into the format. Micro-instruction Word Types and micro-orders are described in Section IV.

3-4. DATA SOURCE AND DATA DESTINATION

Both assembly and micro-assembly language instructions specify data source and data destination. In assembly language one of these is usually a Main Memory address and the other is a register, as in ADA VALU where the A-register is the destination of the data and VALU is the source of the data. With microprogramming both data source and data destination are usually registers, as more registers are available to the microprogram than to the assembly language program.

3-5. DATA MODIFICATION

Add, shift, set flag, and logical functions are performed similarly in both types of programming. In microprogramming, a wider range of basic operations, especially logical functions, is available. Complex operations, such as divide, multiply, and byte move, are performed by micro-programmed subroutines and are available in the Basic Instruction Set and Extended Instruction Group microprograms.

HEWLETT-PACKARD 21MX MICROCODING FORM									
<div> <div> <div>PROGRAMMER</div> <div>JOE CODER</div> </div> <div> <div>DATE</div> <div>6/10/74</div> </div> <div> <div>MICROPROGRAM</div> <div>Read Loader ROM</div> </div> <div> <div>PAGE</div> <div>1 OF 1</div> </div> </div>									
LABEL	OP	SPECIAL	ALU	STORE	S-BUS	COMMENTS	Word Type 1		
LABEL	"IMM"	SPECIAL	MODIFIER	STORE	OPERAND	COMMENTS	Word Type 2		
LABEL	"JMP" OR "JSB"	"COND"	JUMP SENSE	JUMP SENSE	OPERAND	COMMENTS	Word Type 3		
LABEL	"JMP" OR "JSB"	JUMP MODIFIER	JUMP MODIFIER	JUMP SENSE	OPERAND	COMMENTS	Word Type 4		
FIELD 1	FIELD 2	FIELD 3	FIELD 4	FIELD 5	FIELD 6	FIELD 7			
	IMM		LOW	CNTR	0B	CLEAR CNTR (ROM ADDR REG)			
		COV	PASS	M	P	PUT SA IN M: CLR OVF = NO OPER ERR			
LOOP1	L4		PASS	S1	LDR	PASS XXXXXXXXAAAAXXXX INTO S1:CNTR=X			
	ICNT		PASS	L	S1	CNTR=X01			
	L4		AND	S1	LDR	FORM XXXXAAAABBBBXXXX IN S1:CNTR=X01			
	ICNT		PASS	L	S1	CNTR=X10			
	L4		AND	S1	LDR	FORM AAAABBBBCCCCXXXX IN S1:CNTR=X10			
	ICNT		PASS	L	S1	CNTR=X11			
			NAND	S1	LDR	FORM AAAABBBBCCCCDDDD (CMPL FORM)			
	WRITE		PASS	T	S1	WRITE INTO MEMORY			

3-6. DATA TEST AND BRANCH

These operations are quite similar in the two languages. Many tests occur automatically in the course of transferring data in a microprogram. A test and branch out of a line of macro-instructions in normal assembly language, however, requires two instructions (4.6 μ s): a test instruction and a skip instruction.

For example:

```
SLA          skip if LSB of A=0
JMP OUT      branch out of code sequence
```

A test and branch out of a line of micro-instructions requires only one micro-instruction (.325 μ s).

For example:

```
JMP CNDX AL0 OUT  branch out of code sequence
                  if bit 0 of last ALU out-
                  put = 0
```

3-7. MICRO-INSTRUCTION FORMATS

Just as in normal assembly language coding, micro-assembly language source statements are coded in mnemonic form to define an instruction. Each source language statement defines a micro-instruction and consists of an optional label, five micro-order fields some of which may be left blank, and a comment field. The label is used when needed as a reference by other micro-instruction statements. The micro-orders consist of one to four mnemonic characters and specify functions to be performed by the Control Section. According to the type of micro-instruction being defined, one of the micro-orders is sometimes interpreted as an operand. When an operand is specified, it defines an integer or an address, depending on the type of micro-instruction being defined.

3-8. STATEMENT CHARACTERISTICS

Micro-assembly language source statements are divided into four formats, according to the function the micro-instruction is to perform. Each format is called a Word Type.

- Word Type 1 is the most commonly used micro-instruction format and specifies data transfer and modification. Word Type 1 source statement fields are:

```
Label
Op
Special
ALU
Store
S-bus
Comments
```

- Word Type 2 is used to send an 8 bit constant (immediate data) specified in the micro-instruction to a register. Word Type 2 source statement fields are:

```
Label
"IMM"
Special
Modifier
Store
Operand
Comments
```

- Word Type 3 is used to specify a conditional branch in the microprogram. Word Type 3 source statements fields are:

```
Label
"JMP"
"CNDX"
Condition
Jump Sense
Operand
Comments
```

- Word Type 4 is used to specify an unconditional branch in the microprogram. Word Type 4 source statement fields are:

```
Label
"JMP" or "JSB"
Jump Modifier
Operand
Comments
```

3-9. FIELDS

As shown in figure 3-1, the fields are fixed for micro-assembly language source statements. An entry in any field (except comments) must begin in the first column of that field.

- Field 1 begins in column 1 and holds a label that is no longer than eight characters.
- Field 2 begins in column 10 and contains a micro-order no longer than four characters. This field can also hold a Pseudo Instruction (refer to section 4-21 for the explanation of Pseudo Instruction mnemonic codes).
- Field 3 begins in column 15 and contains a micro-order no longer than four characters.
- Field 4 begins in column 20 and contains a micro-order no longer than four characters.

- Field 5 begins in column 25 and contains a micro-order no longer than four characters.
- Field 6 begins in column 30 and contains a micro-order no longer than four characters (Word Type 1) or an operand (Word Types 2, 3, and 4).
- Field 7 begins in column 40 and contains comments only; comments may begin and be placed anywhere from column 40 to column 80 (if column 39 contains the last character of the field 6 operand, field 7 must begin in column 41).

3-10. CHARACTER SET

The characters that may appear in a source statement are as follows:

A through Z
 0 through 9
 . (period)
 * (asterisk)
 + (plus)
 - (minus)
 (space)

Any ASCII character may appear in the comments field.

A space may only begin a field if no micro-order is specified in that field.

3-11. LABEL SYMBOL

A label may be one to eight characters consisting of A through Z, 0 through 9, and a period. The first character must be a letter.

Each label must be unique within the microprogram. Names which appear in \$EXTERNALS micro-assembler control input statements (refer to section 5-5) may not be used as statement labels in the same microprogram.

3-12. ASTERISK COMMENT

An asterisk in column one of the source statement indicates that the entire micro-assembler source statement is a comment.

3-13. MICRO-ORDERS: FIELDS 2 THROUGH 6

The micro-order fields define operations that are to be performed by the Control Section of the computer. The micro-orders applicable to each field are determined by the source statement Word Type. Section IV describes the micro-orders that apply to each Word Type and describes the operations that they specify.

3-14. OPERANDS IN FIELD 6

Word Types 2, 3, and 4 contain an operand in field 6.

In Word Type 2, the operand must be either a decimal or octal number. It cannot be an expression (refer to section 4-10 for definition of a Word type 2 operand).

In Word Types 3 and 4, the operand is a decimal number, octal number, or a number computed from an expression which can include a label (refer to section 4-16 for the definition of a Word Type 3 operand. Refer to section 4-20 for the definition of a Word Type 4 operand).

3-15. CODING THE FOUR WORD TYPES

The following sections describe how to code source statements in micro-assembly language. The reader should be familiar with Section IV of this manual before proceeding with these descriptions. Section IV describes the micro-orders that can be used with each Word Type. By referring to Section IV, the reader can see the options that are available to him as each Word Type is described. The reader will also need to refer to the functional block diagram in Appendix D.

3-16. CODING WITH WORD TYPE 1 — COMMON

This word type specifies data transfer and modification. The format of Word Type 1 is shown in section 4-1. As an example, a micro-instruction is developed that executes the following control functions:

- Store the A-register contents into the M-register
- Perform a memory protect check on the A-register contents
- Transfer the A-register contents to the ALU, increment this value in the ALU, and store the result into the P-register
 - a. Specify the register that is to be placed on the S-bus; the A-register is specified in the example:

OP	SPEC	ALU	STORE	S-BUS
				A

- b. Specify the function of the ALU; the increment function is specified in the example:

OP	SPEC	ALU	STORE	S-BUS
		INC		A

- c. Specify the Op field function; no Op field function is specified in the example. When no Op function is required, the standard operation is specified by either leaving the field blank or inserting NOP into the field:

OP	SPEC	ALU	STORE	S-BUS
NOP		INC		A

- d. Specify a Special function, if required; a memory protect check is specified in the example:

OP	SPEC	ALU	STORE	S-BUS
NOP	MPCK	INC		A

- e. Finally, specify where the resulting data is to be stored. Two store operations are required in the example. The unmodified A-register value on the S-bus must be stored into the M-register and the incremented A-register value on the T-bus must be stored into the P-register. The micro-order PNM performs both of these store operations and serves to illustrate that data stored from the S-bus is unmodified data and data stored from the T-bus can be modified by the ALU or R/S:

OP	SPEC	ALU	STORE	S-BUS
NOP	MPCK	INC	PNM	A

PNM is a unique micro-order. No other micro-order provides the ability to store into two registers in the same micro-instruction.

3-17. CODING WITH WORD TYPE 2 — IMMEDIATE DATA

This word type sends an 8 bit constant (immediate data) specified in the micro-instruction to a register. The format of Word Type 2 is shown in section 4-7. As an example, a micro-instruction is developed that specifies the following control function:

- Repeat the micro-instruction following this one ten times
 - Specify IMM in the Op Code field:

"IMM"	SPEC	MODIF	STORE	OPERAND
IMM				

- Specify the octal or decimal data to be placed on the S-bus; a decimal -10 is specified in the example:

"IMM"	SPEC	MODIF	STORE	OPERAND
IMM				-10

If an octal -10 is specified, it is written -10B.

- Specify one of the four possible data modifiers (refer to section 4-9); LOW (place the 8 bit operand in the lower half of the S-bus and ones in the upper half) is specified in the example:

"IMM"	SPEC	MODIF	STORE	OPERAND
IMM		LOW		-10

- Specify where the resulting data is to be stored; the Counter Register is specified in the example:

"IMM"	SPEC	MODIF	STORE	OPERAND
IMM		LOW	CNTR	-10

- Specify any special operations required; RPT (repeat the micro-instruction following this one the number of times specified in the Counter Register) is specified in the example:

"IMM"	SPEC	MODIF	STORE	OPERAND
IMM	RPT	LOW	CNTR	-10

3-18. CODING WITH WORD TYPE 3 — CONDITIONAL JUMP

This word type specifies a conditional branch in the microprogram. The format of Word Type 3 is shown in section 4-11. As an example, a micro-instruction is developed that specifies the following control function:

- Jump to the microprogram address labeled ERR2, if the last data on the T-bus was not zero.
 - Specify JMP and CNDX in the Op Code and Special fields:

"JMP"	"CNDX"	COND	JUMP	SENSE	OPERAND
JMP	CNDX				

- b. Specify the condition that must be tested for the jump to take place; T-bus equal to 0 is specified in the example:

"JMP"	"CNDX"	COND	JUMP	SENSE	OPERAND
JMP	CNDX	TBZ			

- c. Specify, if required, RJS (Reverse Jump Sense), which establishes whether the Condition code "true" means jump or "false" means jump. The TBZ used in the example means the test condition is T-bus equal to 0. If RJS is specified, T-bus not equal to 0 means perform the jump. If RJS is not specified (blank in the field), then T-bus equal to 0 means jump. RJS is specified in the example:

"JMP"	"CNDX"	COND	JUMP	SENSE	OPERAND
JMP	CNDX	TBZ	RJS		

- d. Specify the target address of the jump. The target address must have the same most significant three

bits as the address of this micro-instruction. The address label ERR2 (an address label in the current page) is specified in the example:

"JMP"	"CNDX"	COND	JUMP	SENSE	OPERAND
JMP	CNDX	TBZ	RJS		ERR2

3-19. CODING WITH WORD TYPE 4 — UNCONDITIONAL JUMP

This word type specifies an unconditional branch in the microprogram. The format of Word Type 4 is shown in section 4-17. As an example, a micro-instruction is developed that specifies the following control function:

- Jump to a microprogram subroutine whose address is derived by the following: the address labeled CLSUB supplies all bits of the subroutine address except bits 3-0; bits 3-0 are supplied by the Instruction Register.

- a. Specify JSB in the Op code field:

"JMP" OR "JSB"	JUMP MODIFIER	--	OPERAND
JSB			

- b. Specify a target address (to be modified) of the jump anywhere within the Control Store (0-7777); CLSUB is specified in the example:

"JMP" OR "JSB"	JUMP MODIFIER	--	OPERAND
JSB			CLSUB

- c. Specify any modification to the target address; J30 (replace bits 3 to 0 of the operand with bits 3 to 0 of the Instruction Register) is specified in the example:

"JMP" OR "JSB"	JUMP MODIFIER	--	OPERAND
JSB	J30		CLSUB

3-20. FROM CODE TO EXECUTION SUMMARY

Figure 3-2 helps to illustrate the process of implementing a microprogram. Writing a micro-assembly language program is essentially the same process as writing an

assembly language program. Micro-instructions are combined to form a microprogram. The microprogram is punched onto cards or paper tape and this source is read by the Micro-assembler. The Micro-assembler produces a listing and an object tape.

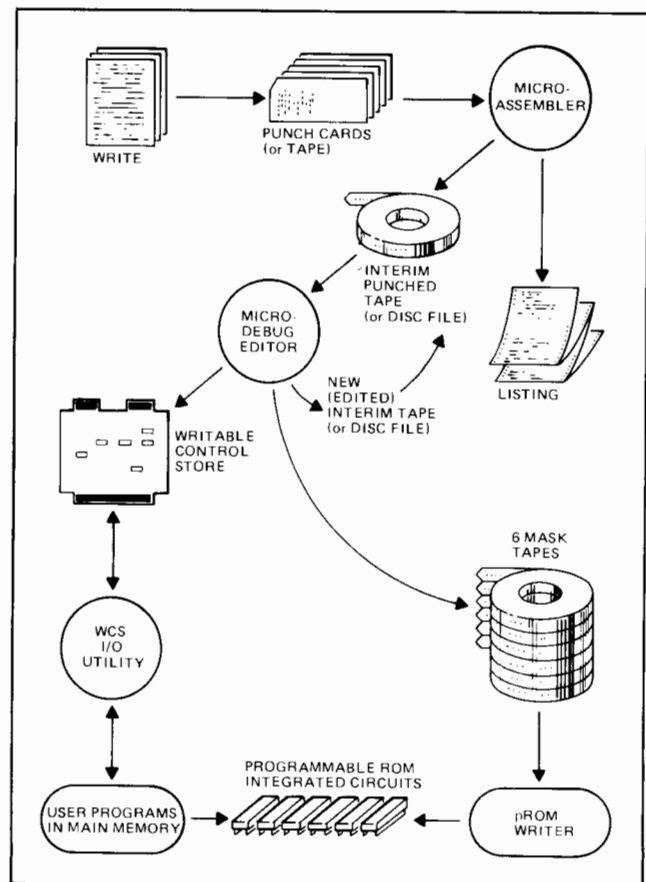


Figure 3-2. Microprogram Implementation Process

The object tape is loaded into Writable Control Store (WCS), executed, and debugged interactively using the Micro Debug Editor (MDE). When the microprogram is debugged, the source is corrected and the microprogram is reassembled. The microprogram can be loaded in two ways. It can be loaded into WCS by a call to the WCS I/O Utility subroutine from the user's Main Memory program or it can be burned into a programmable Read Only Memory. In the latter case, the object tape of the debugged microprogram is loaded into a buffer in Main Memory, using the Micro Debug Editor, and a set of six mask tapes are punched. These tapes are used by the HP 12909 pROM Writer to create ("burn") the programmed Read Only Memory (pROM) chip. The pROM chip is installed on an HP 12945A User Control Store board that is set by jumper wires to specify the proper Control Store module number.

3-21. ACCESS TO MICROPROGRAMS IN CONTROL STORE

The control processor microprograms are divided into three groups.

- a. The 21MX Instruction Set microprograms including the Basic Instruction Set, the Extended Instruction Group, and Floating Point.
- b. Hewlett-Packard supplied special microprograms (for example, the 12977A Fast FORTRAN Processor option) if installed.
- c. User microprograms, if installed.

The control processor reads a 16 bit instruction from Main Memory into the Instruction Register (IR), decodes it, and then determines which microprogram is called for by the instruction. This reading, decoding, and address determination is performed by microprograms that are an integral part of the Basic Instruction Set. The Basic Instruction Set microprogram is in some ways analogous to system software in a normal Main Memory operating system, since the Basic Instruction set performs the general control functions and passes control to the user microprogram area when the Instruction Register calls for a user microprogram. This enables the user-microprogrammer to concentrate effort on his special application.

For the purposes of decoding and implementing macro-instructions, the 21MX Instruction Set is divided into groups according to the general functions they perform. As shown in figure 3-3, there are five groups that encompass the 21MX Instruction Set. A sixth group called the User Instruction Group consists of the macro-instructions that allow the user to access the microprograms which he writes. Most instruction set enhancements or special microprograms will be accessed by the general classification of "user" macro-instructions.

Figure 3-3 summarizes the processing of the Instruction Register. A microprogram within the Basic Instruction Set reads an instruction from Main Memory into the Instruction Register and determines to which macro-instruction group (Alter/skip, Memory Reference, etc.) that instruction belongs. This is accomplished by a ROM table branch command (SPECIAL micro-order "JTAB") that uses the upper eight bits of the Instruction Register to jump, via the fixed ROM Main Look Up Table, to a Control Store microprogram address, according the value of those eight bits. Once the general instruction group is determined, the Instruction Register is further decoded and the logic implemented by the microprogram designed to implement that macro-instruction.

For example, if the instruction in the Instruction Register is in the Extended Arithmetic Unit (EAU) Group, the EAU Group microprogram address is found in the Main Look Up Table based on the Op Code of the instruction. Then the EAU Group microprogram executes the EAU instruction. Provided in the micro-instruction set are special jump parameters, such as "JEAU", to branch within the EAU Group microprogram according to which member of the group is being processed. Jump parameters are explained in Section IV of this manual.

3-22. USER FUNCTION CODE IN ASSEMBLY LANGUAGE

The assembly language program calls a microprogram using mnemonic codes that are assigned in the assembly language program. The pseudo op "MIC" is used to assign the mnemonic code. Refer to the HP Assembler Reference Manual (HP 24307-90014) for the use of the MIC pseudo op.

Using the MIC instruction, a binary function code is assigned to the mnemonic so that whenever the mnemonic appears, the function code is written into that location of the assembled program. The number of parameters is also specified.

The octal function code that calls the user microprogram is:

105rrr if bit 8 of the IR = 0

101rrr or 105rrr if bit 8 of the IR = 1

The value of rrr (bits 8-0) determines the Control Store module address. rrr is defined in table 3-1. Bit 11 in the third digit (5 or 1) is used by micro-instructions which test data in the Instruction Register, where the function code is interpreted. For example, see the "CAB" S-bus micro-order.

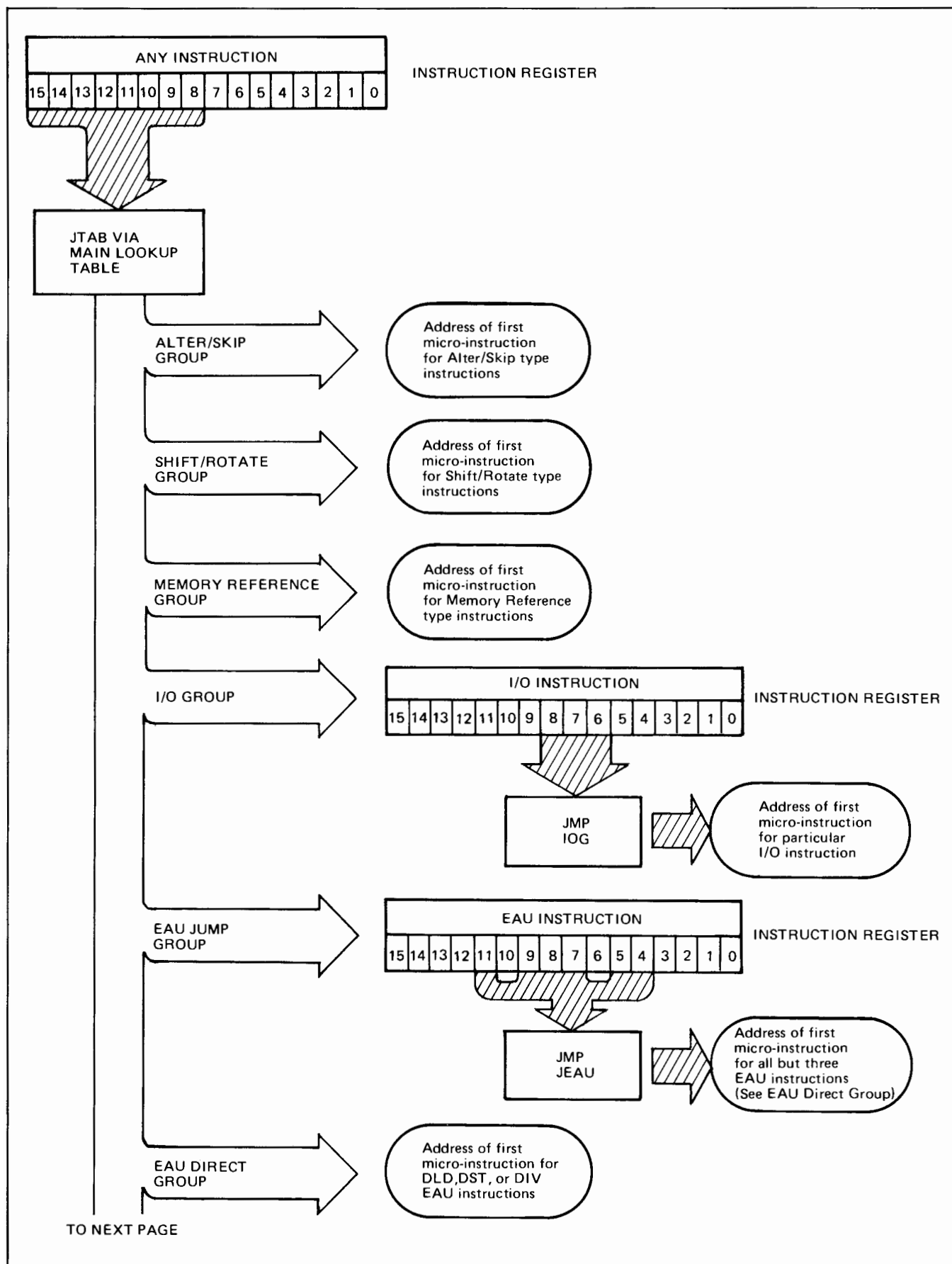


Figure 3-3. Processing the Instruction Register (Sheet 1 of 2)

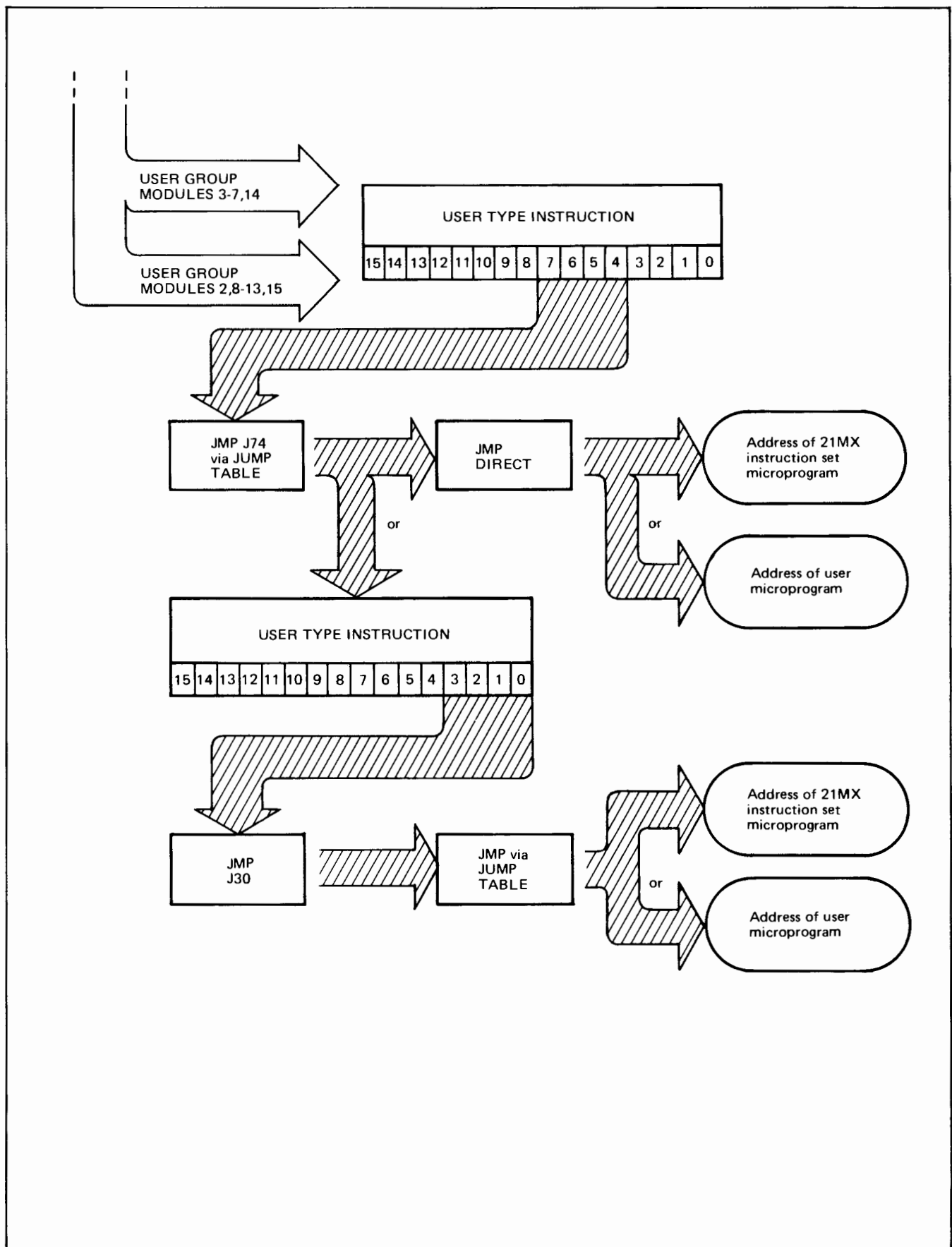


Figure 3-3. Processing the Instruction Register (Sheet 2 of 2)

3-23. CONTROL STORE MODULES AVAILABLE TO USER

The 4096 words of ROM are divided into sixteen 256-word modules, module 0 through module 15. Modules 0, 1, 14, and 15 hold the 21MX Instruction Set and are not available to the user microprogrammer. Modules 12 and 13 are reserved exclusively for user microprograms. Any other Control Store space, not filled by a microprogrammed option, is available to the user microprogrammer. Figure 3-4 summarizes the allocation of Control Store.

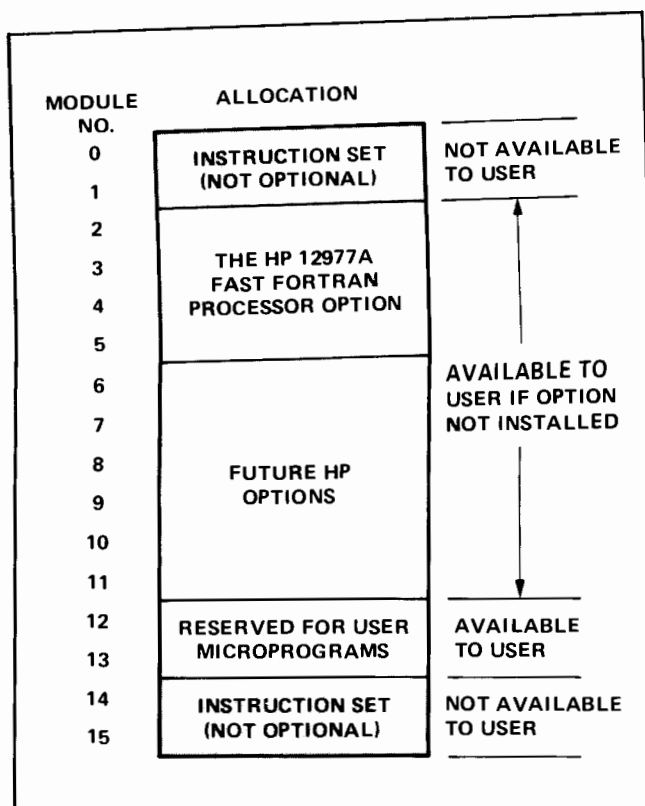


Figure 3-4. Allocation of Control Store by Modules

3-24. MAPPING TO A MODULE ADDRESS

Function codes available to the user are listed in table 3-1 together with the module address to which these function codes map. Some of these user function codes are assigned to the microprogrammed processors and options produced by Hewlett-Packard. The following function codes cannot be used:

105000 through 105137
 105740 through 105777
 101740 through 101777

If the HP 12977A Fast FORTRAN Processor is installed, the following function codes are not available to the user:

105140 through 105277
 105700 through 105737
 101700 through 101737

Note: If the function code maps to a Control Store module which is not present, the micro-instruction

JEAU PASS S S

is executed for each non-existent Control Store location. The ROM Address Register is incremented after each execution of the above micro-instruction until an installed module is encountered. No notification is given to the user or system that a non-existent module is being executed.

Table 3-1. User Function Code Mapping

Function codes 101rrr ₈ and 105rrr ₈ map to the module address given:			
	RANGE OF rrr VALUES	MODULE	RANGE OF OCTAL ADDRESSES
105rrr ₈ only	140 to 157	3	1400
	160 to 177	3	1400 to 1417
	200 to 217	4	2000
	220 to 237	4	2000 to 2017
	240 to 257	5	2400
	260 to 277	5	2400 to 2417
	300 to 317	6	3000
	320 to 337	6	3000 to 3017
	340 to 357	7	3400
	360 to 377	7	3400 to 3417
	400 to 417	8	4000
	420 to 437	8	4000 to 4017
	440 to 457	9	4400
	460 to 477	9	4400 to 4417
	500 to 517	10	5000
	520 to 537	10	5000 to 5017
101rrr ₈ or 105rrr ₈	540 to 557	11	5400
	560 to 577	11	5400 to 5417
	600 to 617	12	6000
	620 to 637	12	6000 to 6017
	640 to 657	13	6400
	660 to 677	13	6400 to 6417
	700 to 717	2	1000
	720 to 737	2	1000 to 1017

3-25. MICROPROGRAMMING INPUT AND OUTPUT FUNCTIONS

Microprogramming Input and Output (I/O) functions requires more care than any other type of microprogramming, because there are strict timing dependencies. The microprogram described in section 3-40 is an example of I/O microprogramming.

To maintain integrity of the I/O system, every control signal which goes to the I/O devices is generated in a specific time period (T-period). All micro-instructions, except those containing READ or WRTE micro-orders, are executed in one I/O T-period, where $T = 325$ ns. READ and WRTE each require two I/O T-periods. An I/O time cycle consists of five T-periods labelled T2, T3, T4, T5, and T6. Specific I/O activity is restricted to certain T-periods in order to synchronize setting of data flags, latching of data, and resolving of multiple interrupt requests.

The microprocessor must synchronize with T2 before initiating an I/O cycle. Thereafter, special consideration must be given to the order and timing of the I/O micro-instructions given.

3-26. SYNCHRONIZING WITH THE I/O SYSTEM

To initiate an I/O cycle, the IOG micro-order must be specified. When this occurs, the processor "freezes" (ceases executing micro-instructions) until time T2. The next micro-instruction is executed during time T3, the

next during T4, etc. IOG may occur with any micro-instruction which does not require some other Special or Jump Modifier (Field 3) micro-order.

Examples:

- a. READ IOG INC PNM P
- b. IOG PASS IR S3

3-27. I/O SIGNAL GENERATION

When IOG is specified, the I/O system generates control signals to the I/O devices starting at the next T2 time and according to the contents of the Instruction Register (IR).

IR bits 5-0 hold a Select Code (SC) signal (SC = The I/O slot number on the backplane or in I/O extenders) that determines which device will respond to the control signal. IR bits 11-6 determine which I/O signals are sent, as shown in table 3-2. The IR must be loaded prior to or during occurrence of the IOG to insure that the correct signals are generated to the proper SC.

Select Codes 0, 1, 2, 3, 4, and 5 have special functions concerning, respectively, the interrupt system, the Front Panel, the Dual Channel Port Controller (DCPC), Power Fail, and Memory Protect/parity. The "Interrupt and Control summary" table in the Appendix of the **HP 21MX Computer Series Reference** manual (HP 02108-90002) holds a description of the effect of these select codes (S.C. in the table).

Table 3-2. I/O Control Signal Generation Determined by IR Bits 11-6

IR*						I/O SIGNAL	TIME	GENERAL USE
11	10	9	8	7	6			
x	x	x	0	0	0	none	T3	Turns off the Run Flag on the CPU.
x	x	x	0	0	1	STF	T3	Set device flag.
x	x	1	x	x	x	CLF	T4	Clear device flag.
x	x	x	0	1	0	SFC	T3-T5	SKPF condition is true if and only if the device flag is clear.
x	x	x	0	1	1	SFS	T3-T5	SKPF condition is true if and only if the device flag is set.
—	—	—	—	—	—	IOI	T4-T5	Buffer the input data latch on the device onto the I/O-bus; this command must be stated explicitly in micro-code during these times.
x	x	x	1	1	0	IOO	T3-T4	Store the I/O-bus into the input data latch on the device.
0	x	x	1	1	1	STC	T4	Set device control flag.
1	x	x	1	1	1	CLC	T4	Clear device control flag.

*Bits marked with x are not significant for the I/O signal specified.

3-28. MEMORY PROTECTION IN RELATION TO I/O MICROPROGRAMMING

When the Instruction Register is loaded, the Memory Protect (MP) feature (12892A) records information on the instruction (from Main Memory) being stored in the IR. When an IOG micro-order is specified, MP checks the select code. If it is not equal to 1 (Front Panel) and MP control is set, MP will inhibit any I/O signals and prevent the CPU from altering memory or the P- or S-registers, and will generate an interrupt request. The microprogrammer cannot prevent this function, so the software operating system maintains security of I/O programming with MP in the microprogramming environment.

3-29. I/O CONTROL ROUTINE

This type of I/O function requires no data transfer. The IR must specify:

STF
CLF
SFS
SFC
STC
CLC

Note that CLF can be generated in conjunction with any other signal by merely letting bit 9 of the IR equal one. To simulate a CLF macro-instruction, specify CLF with STF. Once IOG has been given in an I/O control routine, there are no limitations in using micro-instructions because I/O signals are generated automatically.

For SFS and SFC, the state of the flag on the device may be tested with a "JMP CNDX SKPF" instruction. SKPF is true only when SFS is being executed and the flag is set, or when SFC is being executed and the flag is clear. The SKPF test should occur during T4 or T5 of a SFS or SFC routine. Any operation desired may be implemented as a result of this test. To cause a macroprogram skip, simply increment the P-register contents.

3-30. I/O OUTPUT ROUTINE

This routine is characterized by generation of the IOO micro-order. The IOO sends data from the I/O-bus into the input data latch on the device. The microprogram must put the proper data on the S-bus, then direct it onto the I/O-bus. The detailed timing requirements are:

- During T3, the S-bus must be driven by the register containing the output data to prepare for the transfer to the I/O bus.
- During T4 and T5, the S-bus must be driven by the same register and IOO must appear in the Store field. This insures valid data on the I/O bus.

For example, the sequence for a standard OTA macro-instruction is:

(Time T2)	IOG		
(Time T3)	PASS	CAB	
(Time T4)	IOO	CAB	
(Time T5)	RTN	IOO	CAB

3-31. I/O INPUT ROUTINE

This routine is characterized by use of IOI in the S-bus field. IOI is used in the I/O cycle during T4 or T5 to input data from the I/O device PCA onto the I/O-bus and then onto the S-bus. Any normal Word Type 1 instruction may be used to store the data input from the S-bus.

For example:

(Time T2)	IOG			
(Time T3)	NOP			
(Time T4)	NOP			
(Time T5)	RTN	PASS	CAB	IOI

It can be seen that during some parts of some I/O routines, there are instruction times which are unused. Caution is required when using these times. Do not use micro-instructions which may cause the processor to freeze (listed in section 3-36), until all I/O related code has been executed for that I/O cycle. In the above example, if the T3 and T4 NOPs were replaced by READ and T (S-bus field) micro-orders, the CPU would freeze in the middle of T4 and IOI would not be executed until T6 — too late to correctly handle the data transfer. On the other hand, during a control type routine which is not performing an SFS or SFC, many kinds of micro-instructions can be performed after the IOG. These include READ or even another IOG, since the I/O system requires no further assistance from the microprocessor.

3-32. INTERRUPT HANDLING

The presence of a pending interrupt or halt request may be detected by microcode in two ways:

- Performing a test with JMP CNDX on INT, NHOI, or RUN.
- Attempting to JMP or RTN to location 0 in Control Store; a pending interrupt or halt will cause Control Store address 4 to be loaded into the RAR.

The interrupt device select code (SC) can be read onto the S-bus (high order bits = 0) by specifying CIR in the S-bus field. This freezes the CPU until T6 and then sends IAK to the interrupting device. In the Basic Instruction Set microprogram, the select code from the CIR is loaded into the M-register and the Main Memory instruction at that address is executed. Note that the P-register is not altered during this process.



3-33. NORMAL USER INTERRUPT HANDLING APPLICATIONS

If a long microprogram is entered, the program itself has complete control over when it is terminated or suspended for a detected interrupt. It is not desirable to hold off interrupts very long. Magnetic tape, for example, might request an interrupt every 27 microseconds, if not transferring data by way of the Dual Channel Port Controller.

It is up to the microprogrammer to decide how long to wait before testing for an interrupt. When an interrupt is detected, a jump should be made to a routine to save whatever is necessary to allow the microprogram to continue after the interrupt is serviced or to provide for complete restart of the microprogram. The P-register must be reset to point to the Main Memory address of the macro-instruction interrupted. If parameters are saved, a test must be made at the beginning of the microprogram to determine if it was interrupted or if it executes from the beginning.

When the interrupt servicing is started, a JMP or RTN is made to Control Store location 4 where the Basic Set microcode takes the trap cell address from the Central Interrupt Register and then gives control to Main Memory programs which service the interrupt. After the interrupt routine is complete, the interrupted microprogram is restarted (assuming the P-register was reset upon interrupt detection).

3-34. MICRO-ORDERS AFFECTING MEMORY PROTECT

To fully use the level of protection afforded by the 12892A Memory Protect feature, some conventions must be followed in microprogramming to assure proper communication between the processor and the Memory Protect feature (MP).

Note that MP can only be enabled and disabled by the I/O system. There are no microcode commands for it. Refer to the Memory Protect Interrupt section in the **HP 21MX Reference Manual** for further discussion. The micro-orders which communicate with MP are listed below together with a description of their rules and functions:

- a. FTCH (Special field). This reads the M-register into the MP Violation register, clears out the MP Violation flag and resets the Indirect counter. It should be given when the address of the current instruction from Main Memory is being read (READ micro-order) or immediately after. FTCH occurs in the following places in the Basic Instruction Set Microprogram:

1. At location 0, the Fetch routine.
2. At the location MGOOD+1 in the Halt routine to reset the MP Violation flag and to enable alteration of P-register, S-register, and Main Memory from the Front Panel.

3. At location SCAN+12 as part of the single instruction fetch routine, where it serves the same purpose as at location 0.

- b. IR (Store field). Whenever the IR is specified in the Store field, MP records whether the instruction is a Halt, JMP, or neither, and whether or not IR bits 5-0 equal 01 or not. The IR must be loaded prior to initiating an I/O cycle with IOG to insure that the signal decoding logic will take effect.
- c. INCI (Special field). This micro-order should be used whenever another level of indirect addressing is detected by a microprogram. After 3 counts of the Indirect Counter, an ION (enable interrupts) micro-order is effectively performed by the Memory Protect option. A microprogrammed IOFF micro-order will have no effect after this occurs until after the next FTCH is executed.
- d. MPCK (Special field). There is no need to use this memory protect check micro-order if the Memory Protect feature (HP 12892A) is not installed. This micro-order should be used to insure that a microprogram will not alter protected memory. When this micro-order is used and a MP violation is detected:
 1. All future READ instructions put invalid data into the T-register.
 2. No WRTE instructions are performed.
 3. All attempts to alter the P- or S-registers fail.
 4. All I/O signals from the processor are inhibited until after the next FTCH or CIR is executed.
- e. IOG (Special and Jump Modifier). If Memory Protect has been enabled, this micro-order will set the Memory Protect Violation flag if the select code (IR bits 5-0) is not equal to one. If a MP violation is detected, the actions 1 through 4 described in d. MPCK take place.
- f. CIR (S-bus field). This micro-order causes a freeze until T6 and then issues an IAK to acknowledge the granting of an interrupt to the requesting device. If the select code is 5, the Parity indicator on the Front Panel is cleared and the Memory Protect Violation flag is cleared. Whenever CIR occurs, special logic on the Memory Protect PCA determines whether or not the MP should be disabled (Clear the Control bit). This determination is made six micro-instructions after the last CIR:

1. MP is not disabled if an I/O instruction (IOG) is executed that is not a halt.
2. MP is disabled if no I/O instruction (IOG) is executed or a halt is executed.

To re-enable Memory Protect, an STC 5 is required.

3-35. THE EFFECT OF THE DUAL CHANNEL PORT CONTROLLER ON MICROPROGRAMS

The Dual Channel Port Controller (optional hardware) steals full I/O cycles to perform direct transfers between external devices and Main Memory. This process is essentially transparent to the microprogram. The Dual Channel Port Controller (DCPC) is a hardware function that does not employ microcode. If the microprogram interferes with a DCPC cycle, the Control Processor freezes until DCPC completes its cycle. If DCPC takes a sequence of consecutive I/O cycles for input transfers, any attempted IOG, READ, or WRTE micro-orders will freeze the processor until DCPC is finished. If DCPC takes a sequence of consecutive I/O cycles for output transfers, the Memory Reference Group, the Alter/skip Group, and Shift Rotate Group macro-instructions can still proceed at between 40% and 60% normal execution rate; IOG will still freeze the Control Processor.

If DCPC takes as much as 50% of all I/O cycles, the overall efficiency of the basic instruction set execution is 60% to 70% for input or output transfers. Non-main Memory and I/O micro-instruction execution are only frozen 20% of each DCPC cycle. Thus arithmetic and logical micro-instructions execute at 80% efficiency, when DCPC takes every I/O cycle.

3-36. SUMMARY OF SPECIAL TIMING RULES

- a. Always load the M-register before specifying WRTE in the OP micro-order field.
- b. Load the M-register before or during micro-instructions containing READ in the OP field.
- c. Do not alter the T-register unless initiating a WRTE, since the T-register is internal to the Main Memory system and is used by DCPC and the CPU. The T-register is not intended to be a general purpose register, but to be used in referencing Main Memory.
- d. Load the T-register with data to be written in the same instruction as WRTE appears, or DCPC could alter it before WRTE is executed.
- e. The T-register must be placed on the S-bus no later than two micro-instructions after a READ is specified or the T-register will be disabled by the Memory system.
- f. When an I/O cycle (using IOG) is in progress, a READ or WRTE **must not be initiated** before T6 in the cycle under either of the following conditions:

1. An input or output routine (refer to sections 3-29 and 3-30) is in progress.
2. A skip flag test of the I/O system is taking place.

- g. Do not specify a READ or WRTE micro-order in the same micro-instruction that is transferring data from the T-register (T or TAB micro-order in the S-bus field). The reason is that if a freeze occurs as a result of such a READ or WRTE micro-order (see i. below) the data in the T-register will be invalid after the freeze.

For example, a sequence of micro-instruction similar to the following must not take place:

READ	—	INC	PNM	P
—	—	PASS	S4	L
READ	—	INC	M	TAB

- h. Do not start an I/O cycle (using IOG) before data is transferred from the T-register following a READ operation. The reason is that if the IOG results in a freeze (see i. below), the data in the T-register will be invalid.

For example, a sequence of micro-instructions similar to the following must not take place:

READ	—	INC	PNM	P
—		IOG PASS	S4	TAB

- i. The following conditions always cause a micro-processor freeze:
 1. The CIR micro-order is in the S-bus field and either the I/O cycle time is not T6 or the Dual Channel Port Controller is stealing a full I/O cycle.
 2. The IOG micro-order is in the Special field and either the I/O cycle time is not T2 or the Dual Channel Port Controller is stealing a full I/O cycle.
 3. A T or TAB micro-order is in the S-bus field and a READ or WRTE micro-order memory cycle is still in progress.
 4. A READ or WRTE micro-order is in the Op field and one of the following conditions is true:
 - (a) The semi-conductor Main Memory is being refreshed (two micro-instruction cycles are required every 32.5 microseconds for this purpose).
 - (b) The Dual Channel Port Controller is stealing an I/O cycle.
 - (c) A READ or WRTE memory cycle is still in progress.

3-37. SAMPLE MICROPROGRAMS

While reading the sample microprograms, the reader may find it useful to refer to the fold out functional block diagram in Appendix D. This diagram and the micro-order definitions in Section IV are the two basic sets of information used by the programmer in writing a microprogram.

3-38. SWAP MEMORY LOCATIONS

The sample microprogram illustrated in figure 3-5 swaps the contents of two Main Memory locations that are pointed to by the A- and B-registers (no indirect addresses).

Micro-instruction Commentary

READ	INC	M	A
------	-----	---	---

- Put the address in the A-register onto the S-bus.
- Store the S-bus into the M-register.
- Pass the S-bus through the ALU and increment data enabling the A- or B-register addressable test.
- Read the location in Main Memory pointed to by the M-register (this requires 2 micro-instruction cycles).

MPCK	PASS	M
------	------	---

- Put the M-register onto the S-bus.
- Pass the S-bus through the ALU (output not used).
- Since READ requires two cycles, an instruction cycle is available before data is available from memory. And since the M-register holds the address of the location that will eventually be written into, this cycle is used for the memory protect check.

PASS	S1	TAB
------	----	-----

- The read is complete and data from the memory location is in the T-register unless the AAF or BAF Flag is set. If AAF is set, the data is in the A-register. If BAF is set, the data is in the B-register.
- Put memory data on the S-bus.
- Pass S-bus through the ALU and R/S to the T-bus.
- Store data on T-bus into Scratch Pad Register 1 (S1).

READ	INC	M	B
------	-----	---	---

- Put the address in the B-register onto the S-bus.
- Store S-bus into the M-register.
- Pass the S-bus through the ALU and increment data enabling the A- or B-register addressable test.
- Read the Main Memory location pointed to by the M-register.

MPCK	PASS	M
------	------	---

- Put M-register (memory address) onto the S-bus.
- Pass the S-bus data through the ALU.
- Test the address for a Memory Protect violation.

PASS	S2	TAB
------	----	-----

- Put memory data (T-, A-, or B-register contents) onto the S-bus.

Op Code	Special	ALU	Store	S-bus	Comment
\$ORIGIN=2000B					
\$SYMTAB					
	READ	INC	M	A	READ WORD POINTED TO BY A
		MPCK	PASS	M	CHECK ADDRESS
			PASS S1	TAB	STORE DATA IN S1
	READ	INC	M	B	READ WORD POINTED TO BY B
		MPCK	PASS	M	CHECK ADDRESS
			PASS S2	TAB	STORE DATA IN S2
	WRTE		PASS TAB	S1	BEGIN WRITE
			INC	M	LOAD M WITH A
	WRTE	RTN	PASS TAB	S2	WRITE AND RETURN
\$END					

Figure 3-5. Swap Microprogram

- b. Pass S-bus through the ALU and R/S to the T-bus.
c. Store data on the T-bus into Scratch Pad Register 2 (S2).

WRTE	PASS	TAB	S1
------	------	-----	----

- a. The contents of the first memory location is in S1. Put S1 onto the S-bus.
b. Store the S-bus into T-register (or A- or B-register if AAF or BAF, respectively are set).
c. Pass S-bus data through the ALU.
d. Write T-register contents into Main Memory at address pointed to by the M-register. Note that the M-register still holds the second memory location address. It was loaded during last read operation.

INC	M	A
-----	---	---

- a. The A-register holds the first memory location. Put the A-register contents onto the S-bus.
b. Store the S-bus into the M-register.
c. Pass S-bus data through the ALU and increment data enabling the A- or B-register addressable test.

WRTE	RTN	PASS	TAB	S2
------	-----	------	-----	----

- a. The contents of the second memory location is in S2. Put S2 onto the S-bus.
b. Store the S-bus into the T-register (or A- or B-register, if AAF or BAF, respectively, are set).
c. Pass S-bus data through the ALU.
d. Write the T-register contents into Main Memory at the address pointed to by the M-register.
e. Exit (RTN micro-order).

3-39. BLOCK MOVE MICROPROGRAM

The sample program illustrated in figure 3-6 moves a group of words in Main Memory from one location to another. When the microprogram receives control, it is assumed that:

- The negative value of the number of words to be moved is in the A-register in two's complement form.
- The FROM address is in the B-register.
- The TO address is in the Main Memory location pointed to by the P-register and cannot be indirect.

	Op Code	Special	ALU	Store	S-bus	Comment
	\$ORIGIN=2000					
	\$SYMTAB					
	\$FILE=FILMOV					
	JMP				MOVE	
MOVE			PASS S1	A		WORD COUNT = 0 ?
	JMP	CNDX	TBZ		OUT	IF ZERO, THEN GO TO "OUT"
*						
	READ		INC M	P		GET "TO" ADDRESS
			PASS S2	TAB		PUT IT IN S2
*						
LOOP	READ		INC M	B		READ A DATA WORD
			PASS S3	TAB		STORE THE WORD IN S3 REG
	READ		INC M	S2		GET "TO" ADDRESS
			INC S2	S2		INCREMENT "TO" ADDRESS
	WRTE		PASS T	S3		WRITE A DATA WORD TO MEMORY
			INC B	B		INCREMENT "FROM" ADDRESS
			INC S1	S1		INCREMENT WORD COUNT
	JMP	CNDX	TBZ	RJS	LOOP	GO TO "LOOP" IF WORD
*						COUNT IS NOT ZERO
OUT		RTN	INC P	P		INCREMENT THE P REG AND EXIT
\$END						

Figure 3-6. Block Move Microprogram

The HP assembly language calling sequence is as follows:

```
LDA — (number-of-words)
LDB FROM-address
OCT 105200
DEF TO-address
```

Note: This microprogram is a translation of the Block Move microprogram shown in Section VI of the **HP 2100 Computer Microprogramming Software** manual (HP 02100-90133). Thus it can be used to compare HP 2100 microprogramming to HP 21MX microprogramming.

Micro-instruction Commentary

MOVE	—	—	PASS	S1	A
JMP		CNDX	TBZ	—	OUT

Store the contents of the A-register in Scratch Pad Register 1. If the contents of the A-register are zero, then go to OUT address and return to the calling program.

READ	—	INC	M	P
—	—	PASS	S2	TAB

Get the TO address and store it in Scratch Pad Register 2. The TO address cannot be indirect.

LOOP	READ	—	INC	M	B
—	—	—	PASS	S3	TAB

Read a data word from the Main Memory location pointed to by the FROM address and store the data word in Scratch Pad Register 3. Note that a Control Processor freeze will occur.

		INC	M	S2
		INC	S2	S2
WRTE	—	PASS	T	S3

Write the data (in Scratch Pad Register 3) into memory. Increment TO address pointer.

—	—	INC	B	B
—	—	INC	S1	S1
JMP	CNDX	TBZ	RJS	LOOP

Increment the FROM address pointer. Increment the word count. If the word count is not zero, go to LOOP.

OUT	—	RTN	INC	P	P
-----	---	-----	-----	---	---

Increment the P-register beyond the word containing the TO address and exit.

3-40. INPUT, SUM, AND SUM OF SQUARES MICROPROGRAM

The sample microprogram illustrated in figure 3-7 loads a 16 bit word from a device specified by its select code "SC". If the word is equal to 177777 (end of transmission word), the microprogram is finished and this is signalled by executing the next instruction in Main Memory; otherwise:

- The word is stored in memory location "DATA" indexed by the X-register.
- The word is added to a running total kept in memory location "SUM"
- The word is squared and added to a running total of squares in memory location "SQUAR".
- Another input is initiated from the specified device (STC SC,C).
- The next instruction in Main Memory is skipped to indicate that 177777 was not input from the specified device.

Conditions:

- All numbers are 16 bit positive integers.
- If SUM exceeds $2^{16}-1$, the Extend Register is set.
- If SQUAR exceeds $2^{16}-1$, the Overflow Register is set.
- If both SUM and SQUAR are less than $2^{16}-1$, the Extend and Overflow Registers are clear.
- Memory protect check is performed on addresses used for a write into Main Memory.

Microprogram storage:

The microprogram resides in module 12 starting at octal address 6017.

Microprogram initiation:

Entry into the microprogram is caused by the execution of the following 5 words in Main Memory:

```
105637  USER CALL TO CONTROL STORE
        ADDRESS 6017
0000nn  nn = SELECT CODE "SC"
0aaaaa  "DATA" STORAGE ADDRESS (a table
        holding all input data)
0bbbbbb "SUM" STORAGE ADDRESS
cccccc  "SQUAR" STORAGE ADDRESS
(end of transmission return) SUMMING TERMINATED BY EOT
(normal return)              SUMMING CONTINUES
```


	Op	Code	Special	ALU	Store	S-bus	Comments
\$ORIGIN=6017B							
	READ			INC	PNM	P	01/READ SC, INC P, SET UP TAB LOGIC
	IMM	L1		CMLO	S1	137B	02/000500 INTO S1=USE FOR INP COM LATER
				PASS	L	TAB	03/STORE SC INTO L
				IOR	S11	S1	04/CREATE INPUT COM 0005NN IN S11
	READ			INC	PNM	P	05/READ DATA ADR, INCR P, SET UP TAB LOGIC
	IMM	L4		CMLO	S1	303B	06/001700 INTO S1 FOR SET CONT COM LATER
				PASS	S3	TAB	07/STORE DATA ADR INTO S3
				PASS	IR	S11	08/LOAD IR WITH INPUT COMMAND
				I0G	I0R	S10 S1	09/
*09/	FREEZE TILL T2, START I/O, CREATE SET CONTROL COMMAND 0017NN IN S10						
				PASS	L	S3	10/T3 STORE DATA ADDRESS INTO L
				ADD	S3	X	11/T4 ADD INDEX TO L, STR INTO S3
	ASG			PASS	A	I0I	12/T5 GET DEV WRD FROM I/O BUS, ST INTO A
*							12.5/CLEAR E (IR6=1)
	JMP	CNDX	ONES			OUT	13/T6 JUMP OUT IF ALL ONES IN DEV WORD
	READ			INC	PNM	P	14/READ SUM ADR, INCR P, SETUP TAB LOGIC
				INC	X	X	15/INCR INDEX
				INC	M	TAB	16/ STORE SUM ADR IN M, PREPARE TAB LOGIC
	READ						16.5/ READ SUM
	IMM			LOW	CNTR	0B	17/CLEAR CNTR TO PREPARE FOR REPEAT
				PASS	L	TAB	18/STORE SUM INTO L
	ENVE			ADD	S7	A	19/ADD DEVICE WORD TO T, ENBL O&E, ST INS7
	MPCK			PASS		M	20/MEMORY PROTECT TEST ON SUM ADDRESS
	WRTE			PASS	TAB	S7	21/WRITE TOTAL INTO SUM ADDRESS
				COV	PASS	IR	22/CL OV, PUT SET CNTRL CL FLG COM INTO IR
				I0G	PASS	L	23/FRZ TILL T2, ST A INTO L, START I/O
	MPCK			INC	M	S3	24/T3:S3(DATA ADR&X) ST INTO M, MEM PROT
	WRTE			PASS	TAB	A	25/T4:WRITE DEV WORD INTO (DATA&X
	READ			INC	PNM	P	26/T5,T6:READ ADR OF SQUAR, SETUP TAB LOG
				INC	M	TAB	26.5/ PREPARE TAB LOGIC
	READ			INC	P	P	27/INCR P=NORMAL RETURN, READ SQUARE
	RPT			PASS	B	TAB	28/STORE SQUAR INTO B, SETUP REPEAT
	MPY	R1		ADD	B	B	29/
*29/	(A TIMES L)&B, STORE RESULT INTO B,A						
	JMP	CNDX	TRZ			NO,OVER	30/JMP IF MPY RESULTED IN B=0 (MSB IN B)
				SOV			31/SET OV BIT:RESULT GR TH ACCEPTABLE
NO,OVER	MPCK			PASS		M	32/MEM PROT CK ON SQUAR ADDRESS
	WRTE	RTN		PASS	TAB	A	33/WRITE RESULT INTO SQUAR LOCATION, RTN
OUT				INC	P	P	34/INCREMENT P
				RTN	INC	P	35/INCR P TO INDICATE EOT RETURN, RETURN
\$END							

Figure 3-7. Input, Sum, and Sum of Squares Microprogram

The above instruction is coded in assembly language by defining the mnemonic SSI, function code, and four parameters:

- a. Use the MIC pseudo op in the assembler to define the five word instruction by its mnemonic and number of parameters: MIC SSI,105637B,4

- b. Code the following when calling the SSI microprogram:

SSI SC DATA SUM SQUAR
(end of transmission return) SUMMING TERMINATED BY EOT
(normal return) SUMMING CONTINUES

DATA AREA **

SC EQU nnB SELECT CODE OF DEVICE
DATA BSS mm BUFFER ARE TO HOLD ALL INPUT DATA
SUM OCT 0 "SUM" STORAGE LOCATION
SQUAR OCT 0 "SQUAR" STORAGE LOCATION

Micro-instruction Commentary:

READ	—	INC	PNM	P
------	---	-----	-----	---

- Upon entry into the microprogram, P is the address in Main Memory that follows the instruction that calls microprogram. Hence P is the address of the address containing the select code.
- Place the P-register contents on the S-bus. Store the S-bus into the M-register. Pass the S-bus contents through the ALU incrementing the data in the ALU and store the result (from the T-bus) into the P-register. The address on the T-bus is tested by the T-or-A-or-B logic for use by the TAB micro-order.
- Read the contents of the location in Main Memory specified by the address in the M-register. The read requires two cycles.

IMM	L1	CMLO	S1	137B
-----	----	------	----	------

- While the read is still in progress, a memory cycle is used to construct an input command to be used later.
- Place an octal 137 in bits 7-0 of the S-bus. Bits 15-8 are automatically filled with ones.
- Pass the S-bus through the ALU complementing the data. Shift the data left one bit as it passes through the Rotate/Shifter inserting a zero into bit 0.
- Store the T-bus result into Scratch Pad Register 1. The result in S1 = 000500.

—	—	PASS	L	TAB
---	---	------	---	-----

- Store the result of the read from Main Memory (contents of T- or A- or B-register) onto the S-bus (the select code nn was read).
- Store the S-bus into the L-register and pass the S-bus contents through the ALU (the PASS is effectively a non-operation since the T-bus data is not stored).

—	—	IOR	S11	S1
---	---	-----	-----	----

- Place Scratch Pad Register 1 on the S-bus. Perform an "inclusive or" of L-register and S-bus in the ALU and store the result in S11.
- $$\left. \begin{array}{l} S1 = 00050 \\ L = nn \text{ (select code)} \end{array} \right\} \text{IOR} = 0005nn \text{ in } S11$$

The result in S11 is the complete input command for select code = nn.

READ	—	INC	PNM	P
------	---	-----	-----	---

- The P-register now points to the DATA address.
- Place the P-register on the S-bus. Store the S-bus into the M-register. Increment the S-bus contents as it passes through the ALU and store the resulting address into the P-register. The address on the T-bus is tested by the T-or-A-or-B logic for use by the TAB micro-order.
- Read the contents of the address in Main Memory specified by the M-register (read the DATA address).

IMM	L4	CMLO	S1	303B
-----	----	------	----	------

- While the read is still in progress, the memory cycle is used to construct a set control-clear flag I/O command.
- Place an octal 303 in bits 7-0 of the S-bus. Bits 15-8 are automatically filled with ones.
- Pass the S-bus through the ALU complementing the data. Rotate the data left four bits as it passes through the Rotate/Shifter.
- Store the T-bus result into Scratch Pad Register 1. The result in S1 = 001700.

—	—	PASS	S3	TAB
---	---	------	----	-----

- Place the result of the read from Main Memory (contents of T- or A- or B-register) onto the S-bus (the DATA address was read).
- Pass the S-bus data through the ALU and store it into Scratch Pad Register 3.

—	—	PASS	IR	S11
---	---	------	----	-----

- Place Scratch Pad Register 11 on the S-bus and store the S-bus into the Instruction Register (IR). IR now holds the input command 0005nn, where nn is the device select code.

—	IOG	IOR	S10	S1
---	-----	-----	-----	----

- IOG commands the microprocessor to freeze until time T2. At time T2 the input command in the Instruction Register is executed (transmitted to the device).
- The L-register still holds the select code of device.
- Place Scratch Pad Register 1 (holding 001700) on the S-bus. Perform an "inclusive or" with the L-register in the ALU. Store the result (0017nn) into Scratch Pad Register 10.

- d. The net result in S10 is the completed set control — clear flag command.

—	PASS	L	S3
---	------	---	----

- a. Place Scratch Pad Register 3 (holding DATA address) onto the S-bus and then store S-bus into the L-register.
- b. The PASS is essentially a non-operation.

—	ADD	S3	X
---	-----	----	---

- a. Place the X-register (index to the number of words so far input from the device) onto the S-bus.
- b. Add the S-bus to the L-register (now containing DATA address).
- c. Store the result in Scratch Pad Register 3.

ASG	—	PASS	A	IOI
-----	---	------	---	-----

- a. The time is T5. Take the word input from the Device from the I/O-bus and place it on the S-bus.
- b. Pass the S-bus data through the ALU and store it into the A-register.
- c. The IR = 0005nn, where nn is the device select code. Perform an Alter/Skip Group instruction (ASG) according to bits 7 and 6 in the IR. Since bits 7 and 6 = 01, perform a CLE (Clear Extend register bit).

JMP	CNDX	ONES	—	OUT
-----	------	------	---	-----

If the word last passed through the ALU (see previous micro-instruction) was all ones (end of transmission), jump to the location with the label OUT.

READ	—	INC	PNM	P
------	---	-----	-----	---

- a. The P-register now points to the SUM address.
- b. Place the P-register onto the S-bus. Store the S-bus into the M-register. Increment the S-bus contents as they pass through the ALU and store the resulting address into the P-register. The address on the T-bus is tested by the T-or-A-or-B logic for use by the TAB micro-order.
- c. Read the contents of the address in Main Memory specified by the M-register (read the SUM address).

—	—	INC	X	X
---	---	-----	---	---

Increment the X-register, which is an index to the number of words input from the device.

—	INC	M	TAB
---	-----	---	-----

- a. Place the result of the read from Main Memory (contents of T- or A- or B-register) onto the S-bus (the address of the SUM was read).
- b. Store the data on the S-bus into the M-register.
- c. Increment the data in the ALU and place it on the T-bus so that the data is tested by the T-or-A-or-B logic.

READ	—	—	—	—
------	---	---	---	---

Read the contents of the address in Main Memory specified by the M-register (the present SUM value).

IMM	—	LOW	CNTR	0B
-----	---	-----	------	----

- a. While the read is still in progress, the memory cycle is used to clear the Counter Register in preparation for the RPT used later in the microprogram.
- b. Place zero on the lower eight bits of the S-bus. All ones are automatically stored in the upper eight bits.
- c. Store the S-bus into the Counter Register.

—	—	PASS	L	TAB
---	---	------	---	-----

- a. Place the result of the read from Main Memory (contents of T- or A- or B-register) onto the S-bus (the present SUM value was read).
- b. Store the S-bus into the L-register.

ENVE	—	ADD	S7	A
------	---	-----	----	---

- a. The A-register still contains the word input from the device. Place the A-register onto the S-bus.
- b. Enable the Overflow test and Extend Register test in this micro-instruction only.
- c. Add the L-register (current SUM value) to the S-bus in the ALU.
- d. Store the result in Scratch Pad Register 7.

—	MPCK	PASS	—	M
---	------	------	---	---

- a. The M-register still holds the Main Memory address of SUM. Place the M-register onto the S-bus.
- b. Pass the S-bus through the ALU.

- c. Perform a memory protect check on the address since this address will be used for a write into Main Memory.

WRTE	—	PASS	TAB	S7
------	---	------	-----	----

- a. Place Scratch Pad Register 7 (holding the current DATA total) onto the S-bus.
- b. Store the S-bus into the T-register (or A- or B-register according to AAF or BAF flags).
- c. Initiate a write to Main Memory of the data in the T-register to the address in the M-register. This stores the new total of data words from the device back into the Main Memory address of SUM.

—	COV	PASS	IR	S10
---	-----	------	----	-----

- a. Scratch Pad Register 10 holds the set control-clear flag command, 0017nn, where nn = the select code. Place Scratch Pad Register 10 onto the S-bus.
- b. Store the S-bus into the Instruction Register.
- c. Clear the Overflow Register.

—	IOG	PASS	L	A
---	-----	------	---	---

- a. IOG commands the microprocessor to freeze until time T2. At time T2 the set control-clear flag command in the Instruction Register is executed (transmitted to the device).
- b. Place the A-register (which still holds the word input from the device) onto the S-bus.
- c. Store the S-bus into the L-register.

—	MPCK	INC	M	S3
---	------	-----	---	----

- a. Place Scratch Pad Register 3 (which holds DATA address + index X) onto the S-bus.
- b. Store the S-bus into the M-register.
- c. Increment the data as it passes through the ALU and place it onto the T-bus. The data is tested by the T-or-A-or-B logic.
- d. Perform a memory protect check on the S-bus data.

WRTE	—	PASS	TAB	A
------	---	------	-----	---

- a. Place the A-register (which still holds the word input from the device) onto the S-bus.
- b. Store the S-bus into the T-register (or A- or B-register if the AAF or BAF flag is set).

- c. Initiate a write to Main Memory of the data in the T-register to the address in the M-register. This stores the word input from the device into the Main Memory table of DATA values.

READ	—	INC	PNM	P
------	---	-----	-----	---

- a. The P-register now points to the SQUAR address. Place the P-register onto the S-bus.
- b. Store the S-bus into the M-register.
- c. Increment the S-bus data as it passes through the ALU and then store the T-bus into the P-register.
- d. Read the SQUAR address pointed to by the M-register.

—	—	INC	M	TAB
---	---	-----	---	-----

- a. Freeze until last READ is complete, then place SQUAR address just read from Main Memory onto the S-bus.
- b. Store the S-bus into the M-register.
- c. Increment the data as it passes through the ALU and place it onto the T-bus. The data is tested by the T-or-A-or-B logic.

READ	—	INC	P	P
------	---	-----	---	---

- a. Place the P-register onto the S-bus.
- b. Increment the data as it passes through the ALU and store it into the P-register. The P-register now contains the normal Main Memory return address.
- c. Read the SQUAR contents from Main Memory (contains the current total of data squares).

—	RPT	PASS	B	TAB
---	-----	------	---	-----

- a. Place the SQUAR contents (in the T- or A- or B-register) onto the S-bus.
- b. Pass the S-bus through the ALU onto the T-bus and then store the T-bus into the B-register. The B-register now holds the current total of device input word squares.
- c. Repeat the following micro-instruction incrementing the Counter Register after each repeat. When the Counter Register is equal to 377, execute the next micro-instruction.

MPY	R1	ADD	B	B
-----	----	-----	---	---



- a. Perform a multiply step where the multiplier is in the L-register and the multiplicand is in the A-register.
- b. Both the A- and L-registers hold the last word input from the device. The B-register holds the current total of word squares. Thus the result of 16 repeats of this multiply step is to square the word input from the device adding the result to the past total of squares $[(A \times L) + B]$.
- c. The 32 bit result is in the B- and A-registers with the most significant bits in the B-register.

JMP	CNDX	TBZ	—	NO.OVER
-----	------	-----	---	---------

- a. Jump to the location in the microprogram with the label NO.OVER if the last value that passed onto the T-bus was equal to zero.
- b. In a multiply step operation, the last data to go along the T-bus is the data that is stored into the B-register. Since the B-register holds the most significant bits of the multiplication result, if the result exceeds $2^{16}-1$, bits will be set in the B-register.

—	SOV	—	—	—
---	-----	---	---	---

Set the Overflow Register. The result of the multiplication operation (added to the B-register) exceeds $2^{16}-1$.

NO.OVER	—	MPCK	PASS	—	M
---------	---	------	------	---	---

- a. Place the M-register (the SQUAR address) onto the S-bus.

- b. Perform a memory protect check on the address on the S-bus. (To prepare to write the multiplication result back into the Main Memory data location (SQUAR.)

—	WRTE	RTN	PASS	TAB	A
---	------	-----	------	-----	---

- a. Place the A-register (the current total of squares) onto the S-bus.
- b. Store the S-bus into the T-register (or A- or B-register, if AAF or BAF flag is set).
- c. Write the contents of the T-register into Main Memory at the address given in the M-register (the address of SQUAR).
- d. Return to the Control Store address held in the SAVE Register. In general, this means return to 0 to read the next instruction from Main Memory at the address pointed to by the P-register.

OUT	—	—	INC	P	P
-----	---	---	-----	---	---

- a. This micro-instruction (label OUT) is branched to, if the end of transmission character (177777) has been received from the device.
- b. Increment the P-register.

—	—	RTN	INC	P	P
---	---	-----	-----	---	---

- a. Increment the P-register again to point to the end of transmission return address in Main Memory.
- b. Return to the Control Store address held in the SAVE Register. In general, this means return to 0 to read the next instruction from Main Memory at the address pointed to by the P-register.

3-41. READ A WORD FROM A LOADER ROM

The sample program segment illustrated in figure 3-8 reads four 4-bit bytes from a Loader ROM, constructs a 16 bit word, and then stores the word into Main Memory.

Conditions:

- The A-register holds the Main Memory address into which the 16 bits read from the Loader ROM are to be stored.
- The Loader ROM is selected by bits 15 and 14 of the Instruction Register. The particular Loader ROM selected does not affect the example.
- The Counter Register is set to address the first location in the Loader ROM at the beginning of the micro-program segment.

Micro-instruction Commentary:

—	IMM	—	LOW	CNTR	0B
---	-----	---	-----	------	----

- Place a 0 onto the S-bus in bits 7-0; bits 15-8 are automatically filled with ones.
- Store the S-bus into the Counter Register. Since the Counter Register is eight bits long, only bits 7-0 of the S-bus are stored into the Counter Register.
- The Counter Register is now zero.

—	—	—	PASS	M	A
---	---	---	------	---	---

- The P-register holds the Main Memory Address into which 16 bits are to be stored from the Loader ROM.
- Place the P-register contents onto the S-bus.
- Store the S-bus into the M-register for use later in the write to Main Memory of the word from the Loader ROM.

LOOP1	—	L4	PASS	S1	LDR
-------	---	----	------	----	-----

- The LOOP1 label is used to identify this microprogram segment in the Basic Instruction Set microprogram.
- Place a 4-bit-byte, addressed by the Counter Register, onto the S-bus. The Counter Register is equal 0; thus addressing byte 0 (there are 256 bytes addressed octal 0-377 in each Loader ROM). Note that each byte is stored on the S-bus in complemented form. Thus before a 16 bit word is stored into Main Memory, it must be complemented. This is taken care of by the next to last micro-instruction in this program segment.
- Pass the S-bus through the ALU to the Rotate/Shifter. Left shift the data four bits.
- Store the data on the T-bus into Scratch Pad Register 1 (S1). S1 now holds 16 bits of the form:

xxxxxxxxAAAAxxxx

where AAAA is the 4 bit byte just read.

—	—	INCT	PASS	L	S1
---	---	------	------	---	----

- Place Scratch Pad Register 1 onto the S-bus.
- Store the S-bus into the L-register.
- Increment the Counter Register to address Loader ROM byte 1.

—	—	L4	AND	S1	LDR
---	---	----	-----	----	-----

- Place byte 1 of the Loader ROM onto the S-bus.
- Perform a logical "and" of the S-bus and the L-register in the ALU.
- Left shift the data four bits in the Rotate/Shifter.

	Op Code	Special	ALU	Store	S-bus	Comments
	IMM		LOW	CNTR	B	CLEAR CNTR (ROM ADDR REG)
		COV	PASS	M	A	PUT SA IN M
LOOP1		L4	PASS	S1	LDR	PASS XXXXXXXXXXXXXXXX INTO S1; CNTR=X00
		INCT	PASS	L	S1	CNTR=X01
		L4	AND	S1	LDR	FORM XXXXAAAABBBBXXXX IN S1; CNTR=X01
		INCT	PASS	L	S1	CNTR=X10
		L4	AND	S1	LDR	FORM AAAABBBBCCCCXXXX IN S1; CNTR=X10
		INCT	PASS	L	S1	CNTR=X11
			NAND	S1	LDR	FORM AAAABBBBCCCCDDDD (CNPL FORM)
	WRITE		PASS	T	S1	WRITE INTO MEMORY

Figure 3-8. Reading From a Loader ROM

- d. Store the T-bus into Scratch Pad Register 1. S1 is now of the form:

xxxxAAAABBBBxxxx

where BBBB is the 4-bit-byte just read.

—	—	INCT	PASS	L	S1
---	---	------	------	---	----

- Place the contents of Scratch Pad Register 1 onto the S-bus.
- Store the S-bus into the L-register.
- Increment the Counter Register to address Loader ROM byte 2.

—	—	L4	AND	S1	LDR
---	---	----	-----	----	-----

- Place byte 2 of the Loader ROM onto the S-bus.
- Perform a logical “and” of the S-bus and the L-register in the ALU.
- Left shift the data four bits in the Rotate/Shifter.
- Store the T-bus into Scratch Pad Register 1. S1 is now of the form:

AAAABBBBCCCCxxxx

where CCCC is the 4-bit-byte just read.

—	—	INCT	PASS	L	S1
---	---	------	------	---	----

- Place S1 onto the S-bus.
- Store the S-bus into the L-register.
- Increment the Counter Register to address Loader ROM byte 3.

—	—	—	NAND	S1	LDR
---	---	---	------	----	-----

- Place byte 3 of the Loader ROM onto the S-bus.
- Perform a logical “nand” of the L-register and the S-bus (L “and” S, the result complemented) in the ALU.
- Store the T-bus in S1. S1 is now of the form:

AAABBBCCCCDDD

where DDD is the 4-bit-byte just read. S1 now holds the completed 16 bit macro-instruction.

—	WRTE	—	PASS	T	S1
---	------	---	------	---	----

- Place S1 onto the S-bus.
- Store the S-bus in the T-register (the Main Memory Data Register).
- Initiate a write to Main Memory (address in the M-register) of the data in the T-register.

This completes the reading of 4 bytes from the Loader ROM, constructing a 16 bit macro-instruction, and storing the macro-instruction in Main Memory.

MICROPROGRAMMING LANGUAGE

SECTION

IV

This section serves as a reference to micro-instruction word definitions and formats.

There are four micro-instructions word types. Their general uses are defined below:

- **Word Type 1** executes
 - a. Data transfers between Main Memory, I/O, and arithmetic and logic sections.
 - b. Logical and arithmetic functions on data.
- **Word Type 2** specifies octal data to be transferred to a specific register.
- **Word Type 3** executes a conditional jump based on flags or data values.
- **Word Type 4** executes an unconditional jump or sub-routine jump.

In addition, there are five Pseudo Instructions recognized by the micro-assembler.

Each word type has two formats. One format is the 24-bit **Binary Instruction Format**. This is the machine-language format; the format of the micro-instruction as it is stored in the ROM. The second format is the **Mnemonic Format**. This is the micro-assembler source format; the mnemonic-character representation of the micro-instruction.

Each micro-instruction consists of a number of **micro-orders**, which define the control steps to be executed within the system. The binary representation of the micro-orders falls within certain bits of the 24-bit Binary Instruction. The mnemonic representation of each micro-order falls within seven fields of the micro-instruction input record (e.g. a card). The binary and mnemonic formats are defined for word types in the following sections.

Common to all word types are the LABEL (Field 1), COMMENTS (Field 7), and "*" (column 1).

- **LABEL**

This optional field is a string containing any ASCII characters except +, -, or a space. The string of characters can be one through eight characters long and must always start in column one with a "." (period) or a letter. A maximum of 256 locations address labels are allowed in any microprogram.

- **COMMENT**

This optional field can be any string of up to 30 characters.

- *****

The asterisk indicates that the entire input record (card) is a comment field.

4-1. WORD TYPE 1 — COMMON

Character
Column:

1	10	15	20	25	30	40	80
Field 1	Field 2	Field 3	Field 4	Field 5	Field 6	Field 7	
* or LABEL	OP	SPECIAL	ALU	STORE	S-BUS	COMMENTS	

Figure 4-1. Word Type 1 Micro-assembler Mnemonic Format

Bit No.	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Fields	OP				ALU				S-BUS				STORE				SPECIAL							

Figure 4-2. Word Type 1 Binary Format

There are five micro-order classifications in Word Type 1:

- OP — 12 operations
- SPECIAL — 32 special operations
- ALU — 32 ALU functions
- STORE — 32 destinations of data generated by the micro-instruction
- S-BUS — 32 sources for data to be used by the micro-instruction.

Micro-orders for Word Type 1 are defined in the following paragraphs. The mnemonic code is defined first, followed by its binary equivalent, the meaning, and any special conventions in the use of the micro-order.

4-2. OP MICRO-ORDERS

Many operation codes require specific micro-orders in other fields of the micro-instruction. Those that do will be defined in terms of all required and optional micro-orders in the fields of the micro-instruction.

OP
ARS

Required micro-instruction mnemonic fields:

OP	SPECIAL	ALU	STORE	S-BUS
ARS	L1 or R1	PASS	B	B

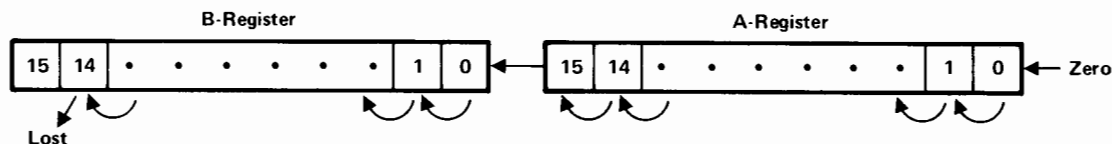
Equivalent micro-instruction binary fields:

23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OP				ALU				S-BUS				STORE				SPECIAL							
0	0	0	1	1	1	1	1	1	0	1	0	1	0	0	1	0	1	0	L1 or R1 Code				

Meaning: Perform a single bit Arithmetic shift of the A- and B-register combined, with the A-register forming the low-order 16 bits. The direction of the shift is specified in the SPECIAL field: L1 for left, R1 for right.

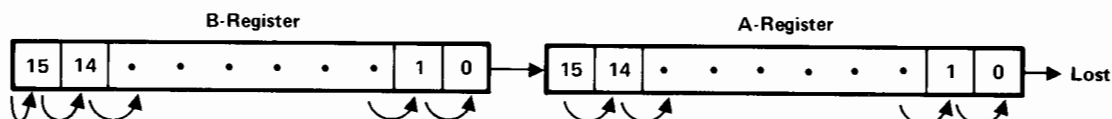
If L1, a 0 is shifted into bit 0 of the A-register; bit 14 of the B-register is lost, but the sign bit remains unchanged. The overflow register bit is set if bits 14 and 15 differ before the shift operation.

ARITHMETIC LEFT SHIFT: SPECIAL=L1



If R1, the sign is copied into bit 14 of the B-register and bit 0 of the A-register is lost.

ARITHMETIC RIGHT SHIFT: SPECIAL=R1



OP	BIT NO.	23	22	21	20
ASG	CONTENT	1	0	0	0

Conventions: This micro-order is used by the Basic Instruction Set microprograms which implement the Alter/skip Macro-instruction Group.

Meaning: Let bits 6 and 7 of the Instruction Register determine which of the following functions is to be performed; then clear the L-register.

IR Bit No.	7	6	
CLE	0	1	Clear Extend Register
CME	1	0	Compliment Extend Register
CCE	1	1	Set the Extend Register

} Alter/Skip instruction

OP
CRS

Required micro-instruction mnemonic fields:

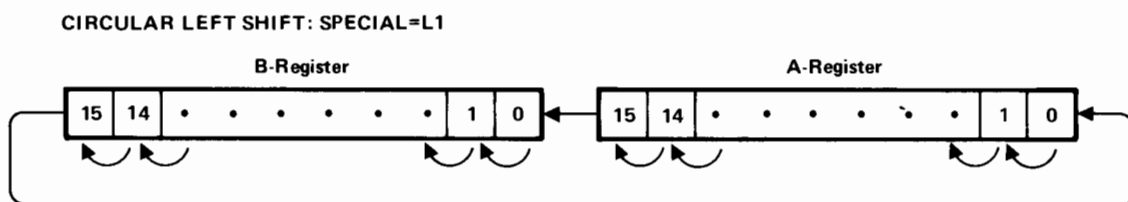
OP	SPECIAL	ALU	STORE	S-BUS
CRS	L1 or R1	PASS	B	B

Equivalent micro-instruction binary fields:

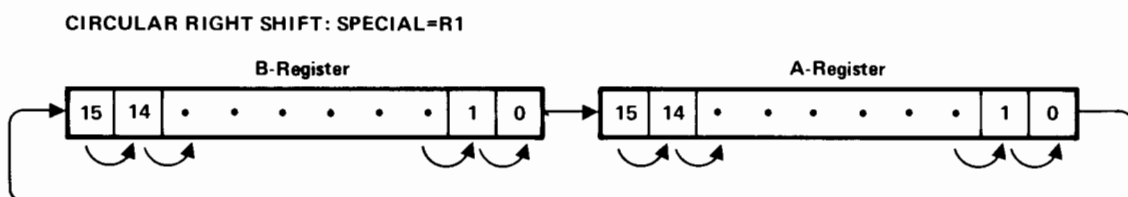
23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OP				ALU					S-BUS					STORE					SPECIAL				
0	0	1	0	1	1	1	1	1	0	1	0	1	0	0	1	0	1	0	L1 or R1 Code				

Meaning: Perform a single bit circular Rotate Shift of the A- and B-registers combined, with the A-register forming the low order 16 bits. The direction of the shift is specified in the SPECIAL field: L1 for left, R1 for right.

If L1, bit 15 of the B-register is transferred to bit 0 of the A-register.



If R1, bit 0 of the A-register is transferred to bit 15 of the B-register.



OP
DIV

Required micro-instruction mnemonic fields:

OP	SPECIAL	ALU	STORE	S-BUS
DIV	L1	SUB	B	B

Equivalent micro-instruction binary fields:

23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OP				ALU				S-BUS				STORE				SPECIAL							
0	1	0	1	0	0	1	1	0	0	1	0	1	0	0	1	0	1	0	0	0	0	1	0

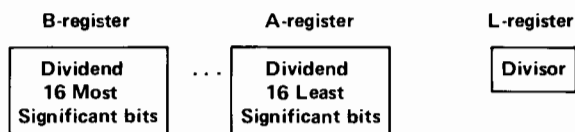
Meaning: Perform a divide step where the divisor is in the L-register and the 32 bit dividend is in the A- and B-registers (least significant bits in the A-register). This micro-order is repeated (16 times for a full word divisor) by specifying the Special micro-order RPT in the preceding micro-instruction. This performs the successive subtractions required in a divide algorithm.

The divide step is executed as follows:

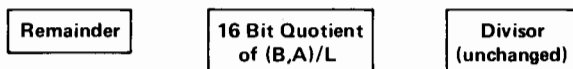
- Subtract the L-register from the B-register (ALU = B - L).
- If borrow is required to complete the subtraction, the ALU Carry Out Flag is clear (0). This Carry Out result means that the divisor (L-register) is too big. The ALU result is not stored. The A-register and B-register are left shifted one bit and the divide step is complete.
- If a borrow is not required to complete the subtraction, the ALU Carry Out Flag is set (1). This Carry Out result means that the divisor is small enough. The result of the subtraction is contained in the ALU and is left shifted one bit and stored back into the B-register. Bit 15 of the A-register shifts into bit 0 of the B-register and bit 0 of the A-register is set to 1 (the Carry Out result). The divide step is complete.

Usage: The base set divide operation is shown in the Basic Instruction Set microprogram in Appendix E at the label = DIV.

Initial Contents:



After Repeat 16 Times of Divide Step:



OP	BIT NO.	23	22	21	20
ENV		1	0	1	0

Meaning: Enable the overflow test for the current ALU operation.

Usage: To detect an overflow condition (that is, set the Overflow register bit), ENV or ENVE (see below) must be specified as the OP Code of the micro-instruction in which the condition is to be tested. Overflow is set if the S-bus and L-register bits 15 are the same and bit 15 output from the ALU is different.

Caution: Caution is advised in the use of DEC (decrement) or INC (increment) in conjunction with ENV. The L-register is always compared.

OP	BIT NO.	23	22	21	20
ENVE		1	0	1	1

Meaning: Enable the overflow test and the extend test for the current ALU operation.

Usage: To detect an Overflow condition (that is, set the Overflow register bit), ENV (see above) or ENVE must be specified as the OP Code of the micro-instruction. To set the Extend Register as a result of the ALU operation, the ENVE micro-order must be specified as the OP code of the micro-instruction. The Extend Register bit is set if there is a carry generated by the ALU (ALU Carry Out = 1).

OP
LGS

Required micro-instruction mnemonic fields:

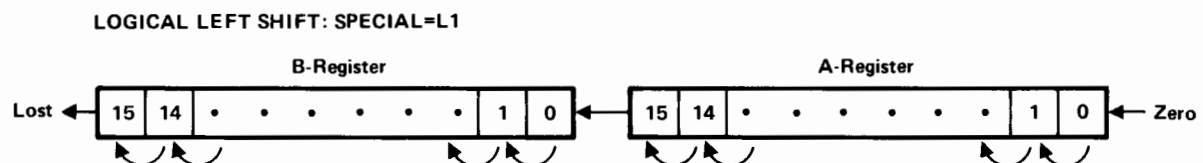
OP	SPECIAL	ALU	STORE	S-BUS
LGS	L1 or R1	PASS	B	B

Required micro-instruction binary fields:

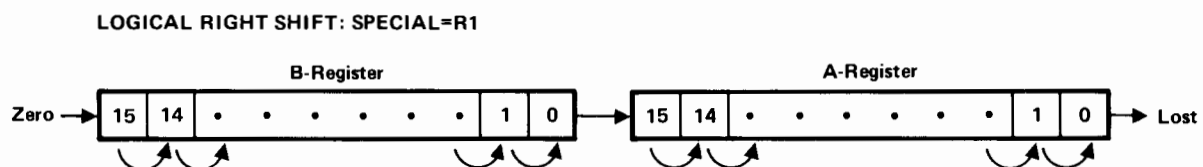
23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OP				ALU					S-BUS					STORE					SPECIAL				
0	0	1	1	1	1	1	1	1	0	1	0	1	0	0	1	0	1	0	L1 or R1 Code				

Meaning: Perform a single bit Logical Shift of the A- and B-registers combined, with the A-register forming the low order 16 bits. The direction of the shift is specified in the SPECIAL field: L1 for left, R1 for right.

If L1, a 0 is shifted into bit 0 of the A-register and bit 15 of the B-register is lost.

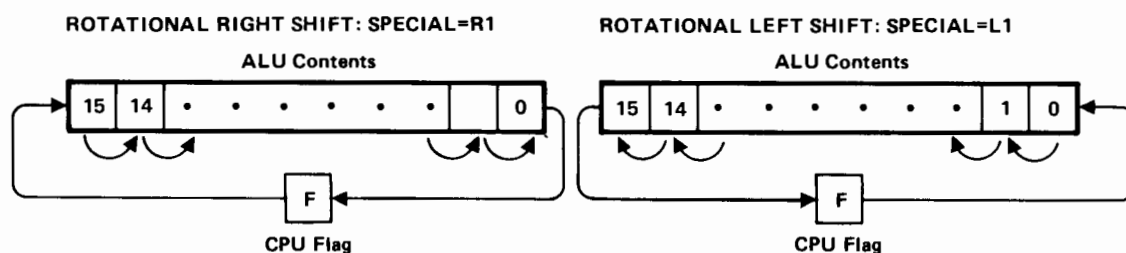


If R1, a 0 is shifted into bit 15 of the B-register and bit 0 of the A-register is lost.



OP	BIT NO.	23	22	21	20
LWF		0	1	1	0

Meaning: Perform a one bit rotational shift of a 17 bit operand in the Rotate/Shifter where bit 17 is formed by the CPU Flag. The rotate moves left one bit, if L1 is the SPECIAL code, or right one bit, if R1 is the SPECIAL code. If neither L1 or R1 are specified, LWF has no effect.



OP
MPY

Required micro-instruction mnemonic fields:

OP	SPECIAL	ALU	STORE	S-BUS
MPY	R1	ADD	B	B

Required micro-instruction binary fields:

23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OP				ALU				S-BUS				STORE				SPECIAL							
0	1	0	0	0	1	0	0	1	0	1	0	1	0	0	1	0	1	0	0	0	1	0	0

Meaning: Perform a multiply step where the multiplier is in the L-register and the multiplicand is in the A-register. The multiply step is executed as follows:

- Test bit 0 of the A-register.
- If the test bit is a one, the L-register is added to the S-bus (B-register value) in the ALU. The result is shifted right one bit and stored back into the B-register with the ALU Carry Out bit forming bit 15.
- If the test bit is a zero, the S-bus (B-register value) is shifted right one bit and stored back into the B-register with the ALU Carry Out bit forming bit 15.
- In either case, the A-register is shifted right and ALU bit 0 fills vacated bit position 15. Bit 0 of the A-register is lost. The multiply step is complete.

Usage: This micro-instruction, repeated 16 times by specifying the SPECIAL code RPT in the preceding micro-instruction, performs the successive additions required in a multiply algorithm. The base set multiply operation is shown in the Basic Instruction Set microprogram in Appendix E at the label =MPY.

Each step of the multiply algorithm effectively multiplies the L-register by the A-register bit that corresponds to the step; that is, step one multiplies the L-register by bit 0 of A-register, step two multiplies the L-register by bit 1 of the A-register, etc. Thus to multiply the L-register by all 16 bits of the A-register, MPY must be repeated 16 times.

Since the B-register goes through successive right shifts and additions as described under "Meaning", the initial contents of the B-register are added to the final result of the multiply algorithm. If the B-register is not zero before the multiply steps are begun, 16 multiply steps will yield the 32 bit result in the B- and A-registers (where the Least Significant Bits (LSB's) are in the A-register):

$$(B,A) = [(A \times L) + B]$$

This may be useful in some computational procedures. For example: $X(2) = X(1) + (Y \times Z)$.

Initial Contents:

B-register	A-register	L-register
Value to be added to the final result	Multiplicand	Multiplier

After Repeating the Multiply Step 16 Times:

(A×L)+B 16 Most Significant bits	(A×L)+B 16 Least Significant bits	Multiplier (unchanged)
-------------------------------------	--------------------------------------	------------------------

OP	BIT NO.	23	22	21	20
READ	CONTENT	1	0	0	1

Meaning: Read data into the T-register from the Main Memory address specified in the M-register. The CPU will freeze until Main Memory is not busy.

Usage: The data must be removed from the T-register two micro-instructions after the READ instruction. Note that the M-register must be loaded (M, PNM, or CM in the Store field) prior to or during the Read micro-instruction. The A- or B-register Addressable Flags (AAF or BAF, respectively) are set, according to data present on the T-bus when the M-register is loaded. Specify INC in the ALU field when the address being stored into the M-register could be a 0 or 1 (A- or B-register addressed). This assures that data is extracted from the proper register when TAB micro-order is used in the S-bus field.

T-bus when M Store is specified	Register Referenced By TAB in S-bus or Store Field	
	AAF	BAF
1	1	0
2	0	1
any other value	0	0

OP	BIT NO.	23	22	21	20
NOP	CONTENT	0	0	0	0

Meaning: Standard Operation. No operation is specified for the Op Code field.

Usage: This is the default micro-order when the OP Code Field is left blank.

OP	BIT NO.	23	22	21	20
WRTE	CONTENT	0	1	1	1

Meaning: Write data from the T-register into the Main Memory address specified in the M-register. The CPU will freeze until Main Memory is not busy. Two micro-instruction times are required to complete the write.

Usage: The T-register should be loaded during the write instruction and must not be altered by the next sequential micro-instruction; otherwise the Dual Channel Port Controller data-transfers could destroy the data.

4-3. SPECIAL MICRO-ORDERS

SPECIAL	BIT NO.	4	3	2	1	0
CLFL	CONTENT	0	1	0	0	1

Meaning: Clear the CPU Flag.

SPECIAL	BIT NO.	4	3	2	1	0
COV	CONTENT	0	1	1	0	0

Meaning: Clear the Overflow Register bit.

SPECIAL	BIT NO.	4	3	2	1	0
FTCH	CONTENT	0	1	0	1	0

Meaning: Move the Main Memory address contained in the M-register (usually the address of the next macro-instruction to be executed) to the Memory Protect Violation Register. Clear out the Memory Protect Violation flag and reset the Indirect Counter.

Usage: This micro-order must be used during, or one micro-instruction after, the initiation of a READ from the address of the next macro-instruction to be executed. This micro-order must be used if the Memory Protect feature is installed on the computer.

SPECIAL	BIT NO.	4	3	2	1	0
ICNT	CONTENT	1	0	0	1	1

Meaning: Increment the Counter Register by one.

SPECIAL	BIT NO.	4	3	2	1	0
INCI	CONTENT	1	0	1	0	1

Meaning: Increment the Indirect Counter in the Memory Protect Option (if installed) by one.

Usage: Used by microprograms that implement indirect addressing. If INCI is executed three times within the same microprogram, the Interrupt Enable Flag is set to allow the CPU to recognize interrupts. Used to prevent multiple indirect addressing levels from holding off recognition of I/O interrupt requests.

SPECIAL	BIT NO.	4	3	2	1	0
IOFF	CONTENT	0	0	0	0	0

Meaning: Turn off the Interrupt Enable flag to disable recognition of normal interrupts (does not disable memory protect, parity, or power fail interrupts).

Usage: After three occurrences of INCI (see INCI Usage) in the SPECIAL Field, interrupts are again recognized and cannot be disabled until a FTCH micro-order occurs. The ION micro-order is normally used to re-enable interrupt recognition.

IOFF should be used with caution, since holding off interrupts could cause the loss of input and output data.

SPECIAL	BIT NO.	4	3	2	1	0
IOG	CONTENT	1	0	0	1	0

Meaning: Freeze the CPU until time period T2. Then execute the base set I/O macro-instruction that is in the Instruction Register.

Usage: Microprogrammed input and output require cooperation between the I/O Section and microprogram control. Familiarity with the I/O system is mandatory. The user may execute input or output in micro-code by loading a standard base set I/O macro-instruction into the Instruction Register and then executing the following micro-instruction:

JSB IOG — — IOCNTL

where IOCNTL = octal 62

See section 3-25 and the following sections for a more detailed description of I/O microprogramming.

SPECIAL	BIT NO.	4	3	2	1	0
ION	CONTENT	0	0	1	0	1

Meaning: Turn the Interrupt Enable flag on to enable recognition of interrupts. Allow the CPU to recognize standard device interrupts until the micro-order IOFF is executed.

Usage: After ION has been executed, the CPU can detect an interrupt from any I/O device in two ways:

- If a JMP or RTN to location 0 of Control Store (the macro-instruction read and decode routine) is executed and an interrupt is pending or the Run flag is clear, execution is forced to location 4 in Control Store, which is the interrupt handler routine.
- A test for interrupt pending or Run flag clear can be performed by the executing microprogram by executing INT, NHOI, or RUN in the Jump Condition field.

ION allows interrupts to be recognized. However interrupts are not generated by the interrupt system until a STF 0 I/O control command is executed. Refer to the discussion of the interrupt system in the **HP 21MX Computer Series Reference Manual**.

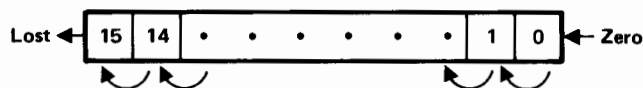
SPECIAL	BIT NO.	4	3	2	1	0
JTAB	CONTENT	1	1	0	1	1

Meaning: Perform a jump to a location within the Basic Instruction Set microprogram, based on the eight most significant bits (bits 15 through 8) of the Instruction Register. This is accomplished via a table look-up of the address in the main jump table for the basic instruction set (see figure 3-2).

The Save Register is cleared to 0. JTAB overrides the effects of JMP or JSB in the OP code field.

SPECIAL	BIT NO.	4	3	2	1	0
L1	CONTENT	0	0	0	1	0

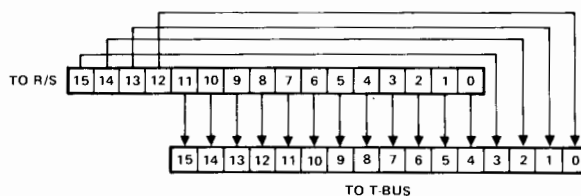
Meaning: Left one bit command to the Rotate/Shifter.



Usage: See MPY, DIV, CRS, LGS, ARS, LWF. Without one of the previous Op Codes, L1 performs a one bit logical left shift on data leaving the ALU.

SPECIAL	BIT NO.	4	3	2	1	0
L4	CONTENT	0	0	0	1	1

Meaning: Four bit circular left shift command to the Rotate/Shifter (R/S).



Usage: Used in conjunction with the shift and rotate operations.

SPECIAL	BIT NO.	4	3	2	1	0
MPCK	CONTENT	1	0	0	0	1

Meaning: Check the address placed on the S-bus for a memory protect violation.

Usage: An S-BUS micro-order must be used in conjunction with MPCK.

This check should be performed before any write to Main Memory (WRTE OP-code), if the memory protect feature is installed. Refer to section 3-27 for details on use of MPCK with the I/O system.

SPECIAL	BIT NO.	4	3	2	1	0
NOP	CONTENT	0	0	1	1	1

Meaning: No SPECIAL operation is performed.

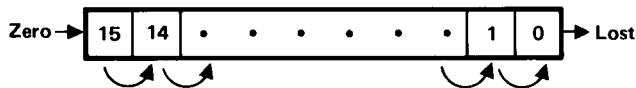
Usage: This is the default operation if none is specified in the SPECIAL field.

SPECIAL	BIT NO.	4	3	2	1	0
RPT	CONTENT	0	1	1	0	1

Meaning: Repeat the following micro-instruction incrementing the Counter Register after each time the repeat is executed. When the lower four bits of the Counter Register are set, execute the following micro-instruction once. The lower four bits of the Counter Register are set at the completion of the repeat sequence. Thus, the repeat is executed the number of times specified in the lower four bits of the Counter Register in two's complement form.

SPECIAL	BIT NO.	4	3	2	1	0
R1	CONTENT	0	0	1	0	0

Meaning: Right one bit command to the Rotate/Shifter.



Usage: Used in conjunction with the shift and rotate instructions. See MPY, DIV, ARS, CRS, LGS, LWF. Without one of the previous micro-orders, a single bit logical right shift is executed.

SPECIAL	BIT NO.	4	3	2	1	0
RTN	CONTENT	1	1	1	1	0

Meaning: Return from subroutine. Jump to the address held in the Save register and clear the Save register.

Usage: No more than one subroutine level is permissible. The second RTN encountered causes a jump to ROM address 0 (the address contained in the Save register) where the macro-instruction pointed to by the P-register is read. RTN overrides the effect of a JMP or JSB in the OP code field.

SPECIAL	BIT NO.	4	3	2	1	0
SHLT	CONTENT	1	0	1	0	0

Meaning: Clear the Run Flag (request a CPU halt).

Usage: The Run Flag is actually cleared at the completion of the micro-instruction following the one specifying SHLT. This micro-order should be used with caution by the microprogrammer. Once the Run Flag is clear, the halt request (SHLT) is detected:

- when a RTN or JMP to address 0 in Control Store (fetch routine) is executed
- when the Run Flag is tested by RUN or NHOI Jump Condition micro-order.

SPECIAL	BIT NO.	4	3	2	1	0
SOV	CONTENT	0	1	0	1	1

Meaning: Set the Overflow Register

SPECIAL	BIT NO.	4	3	2	1	0
SRGE	CONTENT	0	1	1	1	0

Meaning: If Instruction Register bit 5 is set, clear the Extend Register bit.

Conventions: This micro-order is used by the Basic Instruction Set that implements the Extend Register instructions.

SPECIAL	BIT NO.	4	3	2	1	0
SRG1	CONTENT	0	0	1	1	0

Meaning: Execute the Shift/Rotate function specified by bits 6 through 9 of the Instruction Register (Shift/Rotate instruction in the first position; see **HP 21MX Computer Series Reference Manual**.) The Shift/Rotate function is performed on the data that leaves the ALU. The function performed in the R/S is determined by IR bits 6 through 9 as follows:

Bits 9 8 7 6	Function Performed In R/S
1000	Arithmetic left shift one bit
1001	Arithmetic right shift one bit
1010	Rotational left shift one bit
1011	Rotational right shift one bit
1100	Arithmetic left shift one bit, clear sign bit 15
1101	Rotational right shift one bit with E-register forming bit 16 (the 17th bit)
1110	Rotational left shift one bit with E-register forming bit 16 (the 17th bit)
1111	Rotational left shift four bits
0xxx	No shift (bits 8, 7, and 6 can have any setting)

SPECIAL	BIT NO.	4	3	2	1	0
SRG2	CONTENT	0	0	0	0	1

Meaning: Execute the Shift/Rotate function specified by bits 0 1, 2 and 4 of the Instruction Register (Shift/Rotate instruction in the second position; see **HP 21MX Computer Series Reference Manual**.) The Shift/Rotate function is performed on the data that leaves the ALU. The function performed in the R/S is determined by IR bits 0, 1, 2 and 4.

Bits 4 2 1 0	Function Performed in R/S
1000	Arithmetic left shift one bit
1001	Arithmetic right shift one bit
1010	Rotational left shift one bit
1011	Rotational right shift one bit
1100	Arithmetic left shift one bit, clear sign bit 15
1101	Rotational right shift one bit with E-register forming bit 16 (the 17th bit)
1110	Rotational left shift one bit with E-register forming bit 16 (the 17th bit)
1111	Rotational left shift four bits
0xxx	No shift (bits 8, 7, and 6 can have any setting)

SPECIAL	BIT NO.	4	3	2	1	0
SRUN	CONTENT	1	0	1	1	1

Meaning: Set the Run Flag (remove the CPU halt request).

SPECIAL	BIT NO.	4	3	2	1	0
STFL	CONTENT	0	1	0	0	0

Meaning: Set the CPU flag.

4-4. ALU MICRO-ORDERS

ALU	BIT NO.	19	18	17	16	15
ADD	CONTENT	0	1	0	0	1

Meaning: Add the data placed on the S-bus to the contents of the L-register; the L-register contents are not disturbed; pass the result to R/S.

Usage: The L-register must be loaded in a previous micro-instruction.

ALU	BIT NO.	19	18	17	16	15
AND	CONTENT	1	1	0	1	1

Meaning: Logical and of L-register and S-bus ($L \cdot S$); the L-register contents are not disturbed; pass the result to R/S.

Usage: The L-register must be loaded in a previous micro-instruction.

ALU	BIT NO.	19	18	17	16	15
CMPL	CONTENT	1	0	1	0	1

Meaning: Ones complement the L-register; pass the result to Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
CMPS	CONTENT	1	0	0	0	0

Meaning: Ones complement the data on the S-bus; pass the result to Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
DEC	CONTENT	0	1	1	1	1

Meaning: Decrement the data on the S-bus by one; pass the result to the Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
INC	CONTENT	0	0	0	0	0

Meaning: Increment the data on the S-bus by one; pass the result to the Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
IOR	CONTENT	1	1	1	1	0

Meaning: Logical inclusive or of L-register and S-bus ($L + S$); L-register contents are not disturbed; pass result to Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
NAND	CONTENT	1	0	1	0	0

Meaning: Logical nand of L-register and S-bus ($\overline{L \cdot S}$); pass result to Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
NOR	CONTENT	1	0	0	0	1

Meaning: Logical nor of L-register and S-bus ($\overline{L + S}$); pass result to Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
NSAL	CONTENT	1	0	0	1	0

Meaning: Logical and of the complement of the S-bus and the L-register ($\overline{S} \cdot L$); pass result to Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
NSOL	CONTENT	1	1	0	0	0

Meaning: Logical or of the complement of the S-bus and the L-register ($\overline{S} + L$); pass result to Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
ONE	CONTENT	1	1	1	0	0

Meaning: Set all 16 bits (logical one) and pass them to the Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
OP1	CONTENT	0	0	0	0	1

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$(S+L) \text{ plus } 1$$

where "+" means logical function "or".

ALU	BIT NO.	19	18	17	16	15
OP2	CONTENT	0	0	0	1	0

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$(S+\bar{L}) \text{ plus } 1$$

where "+" means logical function "or" and \bar{L} means the ones complement of the L-register (not L).

ALU	BIT NO.	19	18	17	16	15
OP3	CONTENT	0	0	1	0	0

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$S \text{ plus } (\bar{S} \cdot \bar{L}) \text{ plus } 1$$

where " \cdot " means logical function "and" and \bar{L} means the ones complement of the L-register (not L).

ALU	BIT NO.	19	18	17	16	15
OP4	CONTENT	0	0	1	0	1

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$(S+L) \text{ plus } (S \cdot \bar{L}) \text{ plus } 1$$

where " \cdot " means logical function "and", "+" means logical function "or", and \bar{L} means the ones complement of the L-register (not L).

ALU	BIT NO.	19	18	17	16	15
OP5	CONTENT	0	0	1	1	1

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$(S \cdot \bar{L})$$

where " \cdot " means the logical function "and" and \bar{L} means the ones complement of the L-register (not L).

ALU	BIT NO.	19	18	17	16	15
OP6	CONTENT	0	1	0	0	0

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$S \text{ plus } (S \cdot L)$$

where " \cdot " means the logical function "and".

ALU	BIT NO.	19	18	17	16	15
OP7	CONTENT	0	1	0	1	0

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$(S+\bar{L}) \text{ plus } (S \cdot L)$$

where "+" means logical function "or", " \cdot " means logical function "and", and \bar{L} means the ones complement of the L-register (not L).

ALU	BIT NO.	19	18	17	16	15
OP8	CONTENT	0	1	0	1	1

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$(S \cdot L) \text{ minus } 1$$

where " \cdot " means the logical function "and".

ALU	BIT NO.	19	18	17	16	15
OP9	CONTENT	0	1	1	0	0

Meaning: Perform the following logical function in the ALU with the S-bus:

$$S \text{ plus } S$$

ALU	BIT NO.	19	18	17	16	15
OP10	CONTENT	0	1	1	0	1

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$(S+L) \text{ plus } S$$

where "+" means the logical function "or".



ALU	BIT NO.	19	18	17	16	15
OP11	CONTENT	0	1	1	1	0

Meaning: Perform the following logical function in the ALU with the L-register and the S-bus:

$$(S + \bar{L}) \text{ plus } S$$

where "+" means the logical function "or" and \bar{L} means the complement of the L-register (not L).

ALU	BIT NO.	19	18	17	16	15
PASL	CONTENT	1	1	0	1	0

Meaning: Pass the L-register to the Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
PASS	CONTENT	1	1	1	1	1

Meaning: Pass the S-bus data to the Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
SANL	CONTENT	1	0	1	1	1

Meaning: Logical and of the S-bus and the complement of the L-register ($S \cdot \bar{L}$); pass the result to the Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
SONL	CONTENT	1	1	1	0	1

Meaning: Logical or of the S-bus and the complement of the L-register ($S + \bar{L}$); pass the result to the Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
SUB	CONTENT	0	0	1	1	0

Meaning: Subtract the L-register from the S-bus and pass the result to Rotate/Shifter.

ALU	BIT NO.	19	18	17	16	15
XNOR	CONTENT	1	1	0	0	1

Meaning: Logical exclusive nor of the L-register and the S-bus; ($\bar{L} \oplus S$) and pass it to the Rotate/Shifter (\oplus means "exclusive or".)

ALU	BIT NO.	19	18	17	16	15
XOR	CONTENT	1	0	1	1	0

Meaning: Logical exclusive or of the L-register and the S-bus ($L \oplus S$); pass the result to the Rotate/Shifter (\oplus means "exclusive or".)

ALU	BIT NO.	19	18	17	16	15
ZERO	CONTENT	0	0	0	1	1

Meaning: Pass All zeros to the Rotate/Shifter.

4-5. STORE MICRO-ORDERS

STORE	BIT NO.	9	8	7	6	5
A	CONTENT	0	1	0	1	1

Meaning: Store the data on the T-bus in the A-register.

STORE	BIT NO.	9	8	7	6	5
B	CONTENT	0	1	0	1	0

Meaning: Store the data on the T-bus in the B-register.

STORE	BIT NO.	9	8	7	6	5
CAB	CONTENT	0	0	0	0	1

Meaning: Store the data on the T-bus in the A- or B-register according to the value of IR bit 11:

IR bit 11 set means B-register

IR bit 11 clear means A-register

STORE	BIT NO.	9	8	7	6	5
CM	CONTENT	0	1	1	0	1

Meaning: Store the data on the S-bus in the M-register, if the IR holds any Memory Reference instruction except a direct jump (JMP). Refer to the **HP 21MX Computer Series Reference Manual**, for a description of the Memory Reference instructions.

AAF or BAF is set as described under Usage for the M Store micro-order, whether or not the IR holds a Memory Reference instruction.

STORE	BIT NO.	9	8	7	6	5
CNTR	CONTENT	0	0	1	0	1

Meaning: Store the lower eight bits of the S-bus (bits 0-7) in the Counter Register.

STORE	BIT NO.	9	8	7	6	5
DSPI	CONTENT	0	0	1	1	1

Meaning: Store the lower six bits of the S-bus in the Display Indicator on the front panel.

Display Indicator Bit	5	4	3	2	1	0
Register Displayed	S	P	T	M	B	A

Usage: The six indicators on the front panel, labelled A, B, M, T, P and S are lit according to the bit(s) cleared in the Display Indicator. At power-up all bits are set until programmatically changed.

STORE	BIT NO.	9	8	7	6	5
DSPL	CONTENT	0	0	1	1	0

Meaning: Store the data on the S-bus in the Display Register on the Front Panel.

STORE	BIT NO.	9	8	7	6	5
IOO	CONTENT	0	0	1	0	0

Meaning: Direct the S-bus onto the I/O-bus.

Usage: This micro-order when used must be in the second and third instructions (T3 and T4) after IOG Special micro-order. See section 3-25 and the following sections for a description of I/O microprogramming.

STORE	BIT NO.	9	8	7	6	5
IR	CONTENT	0	1	0	0	0

Meaning: Store the data on the S-bus in the Instruction Register. Record the type of macro-instruction stored there in the Memory Protect hardware for use in determining error conditions during Instruction Register execution. See sections 3-28 and 3-34 for a description of Interfacing With Memory Protect feature.

STORE	BIT NO.	9	8	7	6	5
L	CONTENT	0	0	0	1	1

Meaning: Store the data on the S-bus in the L-register (Latch).

STORE	BIT NO.	9	8	7	6	5
M	CONTENT	0	1	0	0	1

Meaning: Store the data on the S-bus in the M-register.

Usage: An ALU micro-order (for example, INC) should also be specified in the micro-instruction. This will activate an A- or B-register addressable test. If bits 14 through 0 on the T-bus equal 1 or 2, the AAF or BAF, respectively, will be set. The M-register may be stored into immediately after a READ or WRTE Op micro-order.

STORE	BIT NO.	9	8	7	6	5
NOP	CONTENT	0	1	1	1	1

Meaning: No store operation is performed; this is the default micro-order when the Store field is left blank.

STORE	BIT NO.	9	8	7	6	5
P	CONTENT	1	1	1	1	0

Meaning: Store the data on the T-bus in the P-register (Program Address Register).

STORE	BIT NO.	9	8	7	6	5
PNM	CONTENT	0	1	1	1	0

Meaning: Store the data on the T-bus in the P-register (Program Address Register), and the data on the S-bus into the M-register (Memory Address Register).

Usage: Useful in microprograms which perform multiword READ operations from Main Memory, where the P-register points to the address in Main Memory to be read. In a single micro-instruction the microprogram can store P into the M-register via the S-bus and then increment P via the T-bus. An example of such an application is the following:

READ - - INC PNM P

The A- or B-register addressable test is activated. See Usage under M micro-order, above.

STORE	BIT NO.	9	8	7	6	5
S	CONTENT	1	1	1	1	1

Meaning: Store the data on the T-bus in the S-register (Switch Register).

STORE	THRU	STORE	BIT NO.				
S1		S12	CONTENT				
			9	8	7	6	5
			1	n	n	n	n

nnnn is binary representation of decimal number 0 + 11

nnnn is binary representation of decimal number 0 + 11

Meaning: Store the data on the T-bus in the indicated Scratch Pad Register S1 to S12.

STORE	BIT NO.				
T	9	8	7	6	5
CONTENT					
0 0 0 1 0					

Meaning: Store the data on the S-bus in the T-register (Memory Data Register).

Usage: This micro-order should occur concurrently when a WRTE micro-order is used. The T-register is internal to the Memory System. It must not be used as a working register.

STORE	BIT NO.				
TAB	9	8	7	6	5
CONTENT					
0 0 0 0 0					

Meaning: Store the data on the T-bus in the A-register if the AAF (A addressable Flag) is set; store the data on the T-bus in the B-register if the BAF (B addressable Flag) is set; store the data on the S-bus into the T-register (Memory Data Register) if neither AAF nor BAF is set.

Usage: Same as T micro-order.

STORE	BIT NO.				
X	9	8	7	6	5
CONTENT					
1 1 1 0 0					

Meaning: Store the data on the T-bus in the X-register.

STORE	BIT NO.				
Y	9	8	7	6	5
CONTENT					
1 1 1 0 1					

Meaning: Store the data on the T-bus in the Y-register.

4-6. S-BUS MICRO-ORDERS

S-BUS	BIT NO.				
A	14	13	12	11	10
CONTENT					
0 1 0 1 1					

Meaning: Direct the data in the A-register onto the S-bus.

S-BUS	BIT NO.				
ADR	14	13	12	11	10
CONTENT					
0 1 0 0 0					

Meaning: An address is formed on the S-bus using IR bits 0-9 and M-register bits 10-14; if IR bit 10 is clear, bits 10-14 of the address formed on the S-bus are clear. Bit 15 is always clear. IR bit 10 is the zero page/current page flag.

S-BUS	BIT NO.				
B	14	13	12	11	10
CONTENT					
0 1 0 1 0					

Meaning: Direct the contents of the B-register onto the S-bus.

S-BUS	BIT NO.				
CAB	14	13	12	11	10
CONTENT					
0 0 0 0 1					

Meaning: Direct the contents of the A- or B-register onto the S-bus according to the value of IR bit 11:

IR bit 11 set means B-register

IR bit 11 clear means A-register

S-BUS	BIT NO.				
CIR	14	13	12	11	10
CONTENT					
0 0 0 1 1					

Meaning: At I/O time T6 place the contents of the Central Interrupt Register onto the S-bus and generate an IAK (Interrupt Acknowledge) signal to the I/O device. (See section 3-33 for CIR description in relation to Interrupt Handling).

Usage: This micro-order must be used after detection of an I/O interrupt to determine the select code of the interrupting device and to acknowledge that the interrupt is being serviced.

S-BUS	BIT NO.				
CNTR	14	13	12	11	10
CONTENT					
0 0 1 0 1					

Meaning: Direct the contents of the Counter Register onto the S-bus. The 8 bit Counter Register is placed onto the low 8 bits of the S-bus; the upper 8 bits are set to ones.

S-BUS	BIT NO.				
DSPI	14	13	12	11	10
CONTENT					
0 0 1 1 1					

Meaning: Direct the six bits of the Display Indicator from the Front Panel to the S-bus.

Usage: See DSPI Store field definition for Display Indicator bit significance.

S-BUS	BIT NO.				
DSPL	14	13	12	11	10
CONTENT					
0 0 1 1 0					

Meaning: Direct the contents of the Front Panel Display Register onto the S-bus.

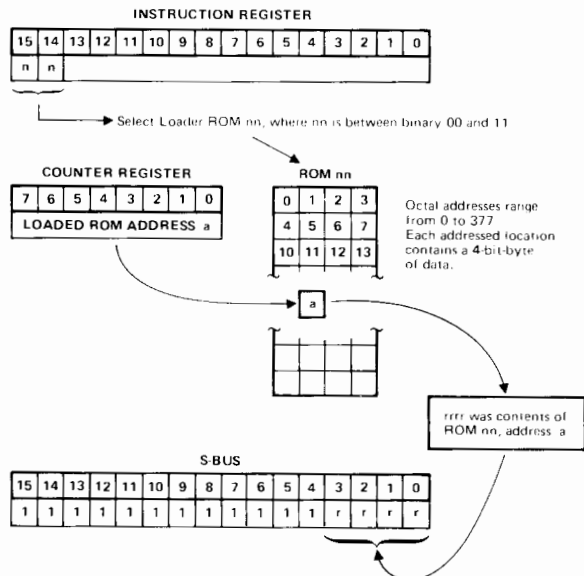
S-BUS	BIT NO.	14	13	12	11	10
IOI	CONTENT	0	0	1	0	0

Meaning: Direct the I/O bus onto the S-bus. (See section 3-25, Microprogramming Input and Output Functions.)

Usage: This is used to transfer data from an I/O device to the S-bus. When not in use, the I/O bus is all zeros. However, do not try to use the I/O bus for a source of zero data, since it is used by the Dual Channel Port Controller at indeterminate times.

S-BUS	BIT NO.	14	13	12	11	10
LDR	CONTENT	0	1	1	0	0

Meaning: Place one 4-bit-byte from a Loader ROM on the S-bus. The 4-bit-byte address is contained in the Counter Register. Determination of which Loader ROM, of the four Loader ROMs available, is specified by bits 15 and 14 in the Instruction Register.



Usage: See sample microprogram in section 3-41 for an illustration of the use of the LDR micro-order.

S-BUS	BIT NO.	14	13	12	11	10
M	CONTENT	0	1	0	0	1

Meaning: Direct the 15 bit contents of the M-register onto the S-bus. Bit 15 of the S-bus is cleared.

S-BUS	BIT NO.	14	13	12	11	10
NOP	CONTENT	0	1	1	1	1

Meaning: The S-bus holds all ones.

Usage: This is the default micro-order when the S-bus field is left blank.

S-BUS	BIT NO.	14	13	12	11	10
P	CONTENT	1	1	1	1	0

Meaning: Direct the contents of the P-register onto the S-bus.

S-BUS	BIT NO.	14	13	12	11	10
S	CONTENT	1	1	1	1	1

Meaning: Place the contents of the S-register (Front Panel Switch Register) onto the S-bus.

S-BUS	THRU	S-BUS	BIT NO.	14	13	12	11	10
S1		S12	CONTENT	1	n	n	n	n

nnnn is binary representation of decimal numbers 0 to 11

Meaning: Place the contents of the indicated Scratch Pad Register S1 to S12 onto the S-bus.

S-BUS	BIT NO.	14	13	12	11	10
T	CONTENT	0	0	0	1	0

Meaning: Direct the contents of the T-register (Memory Data Register) onto the S-bus.

Usage: Data in the T-register that resulted from a READ operation must be removed within two micro-instructions after the READ or the Dual Channel Port Controller could alter the data.

S-BUS	BIT NO.	14	13	12	11	10
TAB	CONTENT	0	0	0	0	0

Meaning: Direct the contents of the T-register (Memory Data Register) onto the S-bus if neither AAF (A addressable Flag) nor the BAF (B addressable Flag) is set; read the A-register onto the S-bus, if the AAF is set; read the B-register onto the S-bus if the BAF is set.

Usage: See T-register Usage description.

S-BUS	BIT NO.	14	13	12	11	10
X	CONTENT	1	1	1	0	0

Meaning: Direct the contents of the X-register onto the S-bus.

S-BUS	BIT NO.	14	13	12	11	10
Y	CONTENT	1	1	1	0	1

Meaning: Direct the contents of the Y-register onto the S-bus.

4.7. WORD TYPE 2 — IMMEDIATE DATA

Character
Column:

1	10	15	20	25	30	40	80
Field 1	Field 2	Field 3	Field 4	Field 5	Field 6	Field 7	
* or LABEL	" IMM "	SPECIAL	MODIFIER	STORE	OPERAND	COMMENTS	

Figure 4-3. Word Type 2 Micro-assembler Mnemonic Format

Bit No.	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Fields	"IMM" op CODE				OPERAND								STORE				SPECIAL							
					MODIFIER																			

Figure 4-4. Word Type 2 Binary Format

There are five micro-order classifications in Word Type 2:

- "IMM" — OP Code specifying Word Type 2.
- SPECIAL — Special operations and modifiers.
- MODIFIER — A special modifier for the Immediate Operation.
- STORE — Destination of the data.
- OPERAND — Binary data that is to be placed on the S-bus.

The STORE and SPECIAL micro-orders applicable to Word Type 2 are exactly the same as those defined for Word Type 1. Consequently, only the other three micro-order groups are defined in the following sections. The "IMM" and MODIFIER micro-order groups are defined by the mnemonic, by its binary equivalent, and finally, by the meaning.

4.8. "IMM" MICRO-ORDER

"IMM"	BIT NO.	23	22	21	20
IMM	CONTENT	1	1	1	0

Meaning: Place 16 bits onto the S-bus consisting of the 8 bit binary OPERAND and another 8 bits of all ones. Determination of which 8 bits of the S-bus receive the OPERAND and which 8 bits receive all ones is made by the MODIFIER.

4.9. MODIFIER MICRO-ORDERS (BITS 19 AND 18 OF THE MICRO-INSTRUCTION)

- Bit 19 Set: specifies complement the S-bus data in the ALU.
- Bit 19 Clear: specifies pass the S-bus data through the ALU.
- Bit 18 Set: specifies OPERAND goes in bits 7-0 of the S-bus.
- Bit 18 Clear: specifies OPERAND goes in bits 15-8 of the S-bus.

MODIFIER	BIT NO.	19	18
CMHI	CONTENT	1	0

Meaning: The 16 bits received by the S-bus consist of the following:

Bits 15-8 = OPERAND

Bits 7-0 = all ones

The S-bus is then complemented as it passes through the ALU.

S-Bus	{	BIT NO.	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			OPERAND								1	1	1	1	1	1	1	1
Out of ALU	{	BIT NO.	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			OPERAND Complemented								0	0	0	0	0	0	0	0

MODIFIER	BIT NO.	19	18
CMLO		1	1

Meaning: The 16 bits received by the S-bus consist of the following:

Bits 15-8 = all ones

Bits 7-0 = OPERAND

The S-bus is then complemented as it passes through the ALU.

S-Bus	{	BIT NO.	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			1	1	1	1	1	1	1	1	OPERAND							
Out of ALU	{	BIT NO.	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

MODIFIER	BIT NO.	19	18
HIGH		0	0

Meaning: The 16 bits received by the S-bus consist of the following:

Bits 15-8 = OPERAND

Bits 7-0 = all ones

The S-bus is then passed through the ALU without modification.

S-Bus and Out of ALU	{	BIT NO.	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			OPERAND								1	1	1	1	1	1	1	1

MODIFIER	BIT NO.	19	18
LOW	CONTENT	0	1

Meaning: The 16 bits received by the S-bus consist of the following:

Bits 15-8 = all ones

Bits 7-0 = OPERAND

The S-bus is then passed through the ALU without modification.

[illegible]

4-10. OPERAND MICRO-ORDER

OPERAND	BIT NO.	19	18	17	16	15	14	13	12	11	10
Integer	CONTENT	Binary Integer Equivalent									

The Integer can be an octal number or decimal number:

- Decimal number in range 0 to 255.
- Octal number in range 0 to 377, followed by “B”.

Examples:

117B, 117, 198, 5, 10B

4-11. WORD TYPE 3 – CONDITIONAL JUMP

Charactor Column :

1	10	15	20	25	30	40	80
Field 1	Field 2	Field 3	Field 4	Field 5	Field 6	Field 7	
* or LABEL	" JMP "	" CNDX "	CONDITION	JUMP SENSE	OPERAND	COMMENTS	

Figure 4-5. Word Type 3 Micro-assembler Mnemonic Format

Bit No.	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Fields	"JM op CODE				CONDITION				JUMP SENSE	OPERAND										"CNDX" SPECIAL CODE				

Figure 4-6. Word Type 3 Binary Format



There are five micro-order classifications in Word Type 3:

- “JMP” — Op Code used in conjunction with “CNDX” specifies Word Type 3, a conditional jump.
- “CNDX” — SPECIAL Code specifying Word Type 3.
- CONDITION — Condition that must be satisfied before jump is executed.
- JUMP SENSE — Optional code to invert the jump condition.
- OPERAND — Target address of jump.

All micro-order groups, except the OPERAND, are defined by the mnemonic, its binary equivalent, meaning, and, where necessary, by conventions in their use.

4-12. “JMP” MICRO-ORDER

“JMP”	BIT NO.	23	22	21	20
JMP	CONTENT	1	1	0	1

Meaning: Used in conjunction with the SPECIAL Code “CNDX”, the CONDITION code specifies the condition under which a jump to the address specified in the OPERAND will take place. If the JUMP SENSE code “RJS” is specified, the CONDITION code specifies the condition under which no jump will take place.

4-13. “CNDX” MICRO-ORDER

“CNDX”	BIT NO.	4	3	2	1	0
CNDX	CONTENT	1	1	0	0	1

Meaning: Used in conjunction with the Op code “JMP”, this micro-order specifies a conditional jump and Word Type 3.

4-14. CONDITION MICRO-ORDERS

The ALU and T-bus condition flags are set after each Word Type 1 or 2 micro-instruction. They are not changed during JMP or JSB micro-instructions (Word Types 3 and 4). Thus, several different jump tests can be made without losing the flag results.

CONDITION	BIT NO.	19	18	17	16	15
AL0	CONTENT	0	0	0	1	1

Meaning: Bit 0 of the last output from the ALU was set (tested before the Rotate/Shifter) by the last Word Type 1 or 2 micro-instruction.

CONDITION	BIT NO.	19	18	17	16	15
AL15	CONTENT	0	0	1	0	0

Meaning: Bit 15 of the last output from the ALU was set (tested before the Rotate/Shifter) by the last Word Type 1 or 2 micro-instruction.

CONDITION	BIT NO.	19	18	17	16	15
ASGN	CONTENT	0	1	1	1	0

Meaning: Alter/skip macro-instruction condition is not satisfied.

CONDITION	BIT NO.	19	18	17	16	15
CNT4	CONTENT	1	1	1	1	0

Meaning: The right (least significant) 4 bits of the Counter Register are all ones.

CONDITION	BIT NO.	19	18	17	16	15
CNT8	CONTENT	0	0	1	1	0

Meaning: All eight bits of the Counter Register are ones.

CONDITION	BIT NO.	19	18	17	16	15
COUT	CONTENT	0	0	0	1	0

Meaning: The ALU Carry Out Flag bit was set by the last ALU operation (tested before the Rotate/Shifter) of the last Word Type 1 or 2 micro-instruction.

CONDITION	BIT NO.	19	18	17	16	15
E	CONTENT	0	1	0	0	1

Meaning: The Extend Register bit is set.

CONDITION	BIT NO.	19	18	17	16	15
FLAG	CONTENT	0	1	0	0	0

Meaning: The CPU FLAG bit is set.

CONDITION	BIT NO.	19	18	17	16	15
FPSP	CONTENT	0	0	1	1	1

Meaning: A special signal is present issued by certain non-standard CPU Front Panels.

CONDITION	BIT NO.	19	18	17	16	15
INT	CONTENT	1	1	0	1	0

Meaning: An Interrupt is pending.

CONDITION	BIT NO.	19	18	17	16	15
IR2	CONTENT	0	1	1	1	1

Meaning: Instruction Register bit 2 is set.

CONDITION	BIT NO.	19	18	17	16	15
NDEC	CONTENT	1	0	0	1	1

Meaning: The "DEC M" (Decrement M-register) button on the Front Panel was **not** actuated.

CONDITION	BIT NO.	19	18	17	16	15
NHOI	CONTENT	0	1	1	0	0

Meaning: The RUN/HALT switch on the Front Panel is set to "Run" and there is no interrupt pending (i.e. no halt and no interrupt).

Usage: This micro-order is recommended for use in long microprograms. (75 microseconds or longer is the criterion used by Hewlett-Packard produced microprograms.) This is necessary since microprograms cannot be interrupted. A pending interrupt or halt condition is not detected unless a specific test is made for them.

CONDITION	BIT NO.	19	18	17	16	15
NINC	CONTENT	1	0	0	1	0

Meaning: The "INC M" (Increment M-register) button on the Front Panel was not actuated.

CONDITION	BIT NO.	19	18	17	16	15
NLDR	CONTENT	1	0	0	0	0

Meaning: The "IBL" (loader) button on the Front Panel was not actuated.

CONDITION	BIT NO.	19	18	17	16	15
NLT	CONTENT	1	0	1	0	1

Meaning: The "←" REGISTER SELECT LEFT button on the Front Panel was not actuated.

CONDITION	BIT NO.	19	18	17	16	15
NMLS	CONTENT	0	0	1	0	1

Meaning: Memory was not lost as a result of the last power down or power failure.

CONDITION	BIT NO.	19	18	17	16	15
NOP	CONTENT	1	1	1	0	1

Meaning: No condition test is made; no jump occurs.

Usage: This is the default micro-order if none is specified in the condition field.

CONDITION	BIT NO.	19	18	17	16	15
NRST	CONTENT	1	0	1	1	1

Meaning: The DISPLAY button on the Front Panel was not actuated.

CONDITION	BIT NO.	19	18	17	16	15
NRT	CONTENT	1	0	1	0	0

Meaning: The "→" REGISTER SELECT RIGHT button on Front Panel was not selected.

CONDITION	BIT NO.	19	18	17	16	15
NSFP	CONTENT	1	1	0	0	1

Meaning: A standard Front Panel is not installed on the CPU.

CONDITION	BIT NO.	19	18	17	16	15
NSNG	CONTENT	1	0	0	0	1

Meaning: The INSTR STEP (Instruction Step) button on the Front Panel was not actuated.

CONDITION	BIT NO.	19	18	17	16	15
NSTB	CONTENT	1	1	0	0	0

Meaning: None of the following Front Panel buttons were actuated:

INSTR STEP (Instruction Step)
 "→" REGISTER SELECT RIGHT
 "←" REGISTER SELECT LEFT
 DISPLAY
 IBL (Binary Loader)
 INC M (Increment M-register)
 DEC M (Decrement M-register)
 STORE
 RUN
 PRESET

CONDITION	BIT NO.	19	18	17	16	15
NSTR	CONTENT	1	0	1	1	0

Meaning: The STORE button on the Front Panel was not actuated.

CONDITION	BIT NO.	19	18	17	16	15
ONES	CONTENT	0	0	0	0	1

Meaning: All 16 bits of the last output from the ALU were set (tested before Rotate/Shifter) as a result of the last Word Type 1 or 2 micro-instruction.

CONDITION	BIT NO.	19	18	17	16	15
OVFL	CONTENT	0	1	0	1	0

Meaning: The Overflow Register bit is set.

CONDITION	BIT NO.	19	18	17	16	15
RUN	CONTENT	0	1	0	1	1

Meaning: The CPU is in RUN mode (the Front Panel RUN flag is set).

CONDITION	BIT NO.	19	18	17	16	15
RUNE	CONTENT	1	1	1	0	0

Meaning: The four position STANDBY/OPERATE/LOCK/R switch on the Front Panel is not in the LOCK position.

CONDITION	BIT NO.	19	18	17	16	15
SKPF	CONTENT	0	1	1	0	1

Meaning: The I/O signal SFS is present (I/O time is T3 to T5) and the addressed I/O device Flag is set or the I/O signal SFC is present (I/O time is T3 to T5) and the addressed I/O device Flag is clear.

Usage: See section 3-25, Microprogramming Input and Output Functions, for the use of the micro-order SKPF.

CONDITION	BIT NO.	19	18	17	16	15
SRGL	CONTENT	1	1	0	1	1

Meaning: Bit 3 of the Instruction Register is set and bit 0 of the last output from the ALU was cleared as a result of the last Word Type 1 or 2 micro-instruction.

Usage: This micro-order is used by the Basic Instruction Set microprogram which implements the SLA and SLB macro-instructions of the Shift/rotate Group.

CONDITION	BIT NO.	19	18	17	16	15
TBZ	CONTENT	0	0	0	0	0

Meaning: The last output from the Rotate/Shifter onto the T-bus was equal to zero as a result of the last Word Type 1 or 2 micro-instruction.

4-15. JUMP SENSE MICRO-ORDER

JUMP SENSE	BIT NO.	14
RJS	CONTENT	0

Meaning: Perform the jump, if the jump condition is not met. The CONDITION micro-order specifies the condition under which a jump can take place; the RJS micro-order in effect reverses the sense of the jump. For example, if a conditional jump is specified if the Flag bit is set (jump if Flag bit set), the RJS micro-order will reverse the condition so that the jump occurs if the Flag bit is not set.

4-16. OPERAND MICRO-ORDER

OPERAND
An Address

BIT NO.	13	12	11	10	9	8	7	6	5
CONTENT	Binary Address Equivalent								

The address can be an octal, decimal or computed number:

Decimal number, d, in the range 0 to 511

Octal number, kB, in the range 0B to 777B, where the B signifies octal.

Computed number, c, which is within the decimal or octal range, according to whether it is computed from octal or decimal values, of the form:

- | | |
|-------------|-------------|
| a. $*+kB$ | f. LABEL-kB |
| b. $*-kB$ | g. LABEL+d |
| c. $*+d$ | h. LABEL-d |
| d. $*-d$ | i. LABEL |
| e. LABEL+kB | |

where * means "this address" and LABEL means a micro-instruction or pseudo-instruction label that is defined elsewhere in the microprogram.

The target address of the jump is not relative and must be within the current 1000 octal locations (two modules). The complete absolute address must be specified. For example, if a conditional jump micro-instruction is within Control Store addresses 3000 and 3777, no target address may be outside the range 3000 to 3777. A target address of 3377B would initiate a jump to the octal address 3377.

Examples:

1005, 2632, 2632B, START, START-11B, END-11

4-17. WORD TYPE 4 — UNCONDITIONAL JUMP

Character Column:

	1	10	15	20	25	30	40	80
Fields:	Field 1	Field 2	Field 3	Field 4	Field 5	Field 6	Field 7	
Content:	* or LABEL	"JMP" or "JSB"	JUMP MODIFIER	(blank)	(blank)	OPERAND	COMMENTS	

Figure 4-7. Word Type 4 Micro-assembler Mnemonic Format

Bit No:	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Fields:	"JMP" or "JSB" Op Code				(zero)		binary OPERAND														JUMP MODIFIER			

Figure 4-8. Word Type 4 Binary Format

Word Type 4 consists of three micro-order classifications:

- "JMP" or "JSB" — Operation, code used in conjunction with the JUMP MODIFIER, specifies Word Type 4, an unconditional jump or subroutine jump.
- JUMP MODIFIER — Specifies modification to the OPERAND jump address.
- OPERAND — Target address of jump, prior to any modification.

Micro-orders, except the OPERAND, are defined by the mnemonic, binary equivalent, meaning, and, where necessary, by conventions in their use.

4-18. "JMP" AND "JSB" MICRO-ORDERS

"JMP" or "JSB"	BIT NO.	23	22	21	20
JMP	CONTENT	1	1	0	1

Meaning: Jump unconditionally to the address specified in the OPERAND, modified according to the JUMP MODIFIER micro-order.

"JMP" or "JSB"	BIT NO.	23	22	21	20
JSB	CONTENT	1	1	0	0

Meaning: Perform a subroutine jump unconditionally to the address specified in the OPERAND, modified according to the JUMP MODIFIER micro-order. The return address is stored in the Save register and recalled by the RTN micro-order (see section 4-3, SPECIAL Micro-orders for RTN definition).

4-19. JUMP MODIFIER MICRO-ORDERS

JUMP MODIFIER	BIT NO.	4	3	2	1	0
IOFF	CONTENT	0	0	0	0	0

Meaning: Disable recognition of normal interrupts (does not disable memory protect, parity, or power fail interrupts). Perform an unconditional jump. No modification is made to the jump OPERAND.

JUMP MODIFIER	BIT NO.	4	3	2	1	0
I0G	CONTENT	1	0	0	1	0

Meaning: Freeze the CPU until time period T2. Execute the I/O function according to the base set I/O macro-instruction that is in the Instruction Register. Perform the JMP or JSB modifying OPERAND bits 2 and 3 according to the I/O instruction jump table (bits 6, 7, and 8 of the I/O macro-instruction in the Instruction Register actually determine the OPERAND address modification):

IR Contains I/O Macro-instruction	IR Bits 8 7 6	OPERAND Bits 3 & 2 Replaced By:
MIA or MIB	1 0 0	1 1
LIA or LIB	1 0 1	1 0
OTA or OTB	1 1 0	0 1
HLT	0 0 0	0 0
CLO or CLF	0 0 1	0 0
STO or STF	0 0 1	0 0
SFC or SOC	0 1 0	0 0
SFS or SOS	0 1 1	0 0
STC or CLC	1 1 1	0 0

See section 3-25 and those following for a more complete description of the use of the IOG micro-order.

JUMP MODIFIER	BIT NO.	4	3	2	1	0
JEAU	CONTENT	1	1	1	1	1

Meaning: Enable the EAU jump table. According to the particular EAU macro-instruction held in the Instruction Register, the least significant three bits (0-2) of the OPERAND are replaced by EAU jump table bits (bits 4-9 and 11 of the Instruction Register actually determine the OPERAND address modification):

EAU Macro-instruction	Three LSB's of Address
RRR	000
ASR	001
LSR	010
(not used)	011
RRL	100
ASL	101
LSL	110
MPY	111

JUMP MODIFIER	BIT NO.	4	3	2	1	0
JIO	CONTENT	1	1	0	1	0

Meaning: Perform the JMP or JSB modifying OPERAND bits 2 and 3 according to the I/O instruction jump table (bits 6, 7, and 8 of the I/O macro-instruction in the Instruction Register actually determine the OPERAND address modification):

IR Contains I/O Macro-instruction	IR Bits 8 7 6	OPERAND Bits 3 & 2 Replaced By:
MIA or MIB	1 0 0	1 1
LIA or LIB	1 0 1	1 0
OTA or OTB	1 1 0	0 1
HLT	0 0 0	0 0
CLO or CLF	0 0 1	0 0
STO or STF	0 0 1	0 0
SFC or SOC	0 1 0	0 0
SFS or SOS	0 1 1	0 0
STC or CLC	1 1 1	0 0

JUMP MODIFIER	BIT NO.	4	3	2	1	0
JTAB	CONTENT	1	1	0	1	0

Meaning: Perform a jump to a location within the Basic Instruction Set microprogram based on the eight most significant bits of the Instruction Register. This is accomplished via a table look up of the address in the Main Jump Table for the basic instruction set. This micro-order is executed independently of word types; hence JMP or JSB need not be specified.

JUMP MODIFIER	BIT NO.	4	3	2	1	0
J30	CONTENT	1	1	1	0	1

Meaning: Replace the four Least Significant Bits of the OPERAND with bits 3 through 0 of the Instruction Register.

JUMP MODIFIER	BIT NO.	4	3	2	1	0
J74	CONTENT	1	1	1	0	0

Meaning: Replace the four Least Significant Bits of the OPERAND with bits 7 through 4 of the Instruction Register.

JUMP MODIFIER	BIT NO.	4	3	2	1	0
RTN	CONTENT	1	1	1	1	0

Meaning: Return to the address stored in the Save Register as a result of a subroutine jump (JSB); if the Save Register is equal to zero (no subroutine is active), return to address 0 of Control Store to initiate the reading of the next macro-instruction from Main Memory.

JUMP MODIFIER	BIT NO.	4	3	2	1	0
STFL	CONTENT	0	1	0	0	0

Meaning: Set the CPU Flag and then perform the JMP or JSB to the OPERAND address. No modification is made to the OPERAND address.

JUMP MODIFIER	BIT NO.	4	3	2	1	0
UNCD	CONTENT	1	1	0	0	0

Meaning: Perform the JMP or JSB to the OPERAND address. No modification is made to the OPERAND address.

Usage: This is the default micro-order if no JUMP MODIFIER is specified.

4-20. THE OPERAND MICRO-ORDER

OPERAND
An Address

BIT NO.	16	15	• • • • •	6	5
CONTENT	Binary Address Equivalent				

The ADDRESS can be a decimal, octal or computed number:

- Decimal number, d, in the range 0 to 4095
- Octal number, kB, in the range 0B to 7777B where B signifies octal
- Computed number, c, which is within the decimal or octal range, according to whether it is computed from octal or decimal values, of the form:
- a. $*+kB$
 - b. $*-kB$
 - c. $*+d$
 - d. $*-d$
 - e. LABEL+kB
 - f. LABEL-kB
 - g. LABEL+d
 - h. LABEL-d
 - i. LABEL

where * means "this address" and LABEL means a micro-instruction label that is defined elsewhere in the microprogram.

Examples:

$*+11B$, $*+9$, HERE+5, START

4-21. PSEUDO INSTRUCTIONS

There are five pseudo instructions recognized by the micro-assembler: DEF, EQU, ONES, SKP, and Zeroes.

4-22. DEF

The DEF statement creates a 24 bit micro-instruction word in ROM the contents of which is a 12 bit binary address defined by "ADDRESS" in the micro-assembler input record (Field 6). The binary address is associated in the microprogram with the optional LABEL, if defined.

The ADDRESS can be a decimal, octal or computed number:

- Decimal number, d, in the range 0 to 4095
- Octal number, kB, in the range 0B to 7777B, where B signifies octal
- Computed number, c, which is within the decimal or octal range, according to whether it is computed from octal or decimal values, of the form:
- a. $*+kB$
 - b. $*-kB$
 - c. $*+d$
 - d. $*-d$
 - e. LABEL+kB
 - f. LABEL-kB
 - g. LABEL+d
 - h. LABEL-d
 - i. LABEL

where * means "this address" and LABEL means a micro-instruction label that is defined elsewhere in the microprogram.

Examples of DEF statements:

Character Column:			
	1	10	30
Fields:	Field 1	Field 2	Field 6
Content:	AD1	DEF DEF DEF	SRF+150 ASGNOP 416B

Character Column:								
	1	10	15	20	25	30	40	80
Fields:	Field 1	Field 2	Fields 3-5			Field 6	Field 7	
Content:	LABEL (optional)	"DEF"	{blank}			ADDRESS	COMMENTS	

4-23. EQU

Character Column:		1	10	15	20	25	30	40	80
Fields:		Field 1	Field 2	Field 3	Field 4	Field 5	Field 6	Field 7	
Content:		LABEL	"EQU"	(blank)	(blank)	(blank)	ADDRESS	COMMENTS	

The EQU statement associates the stated LABEL with a 12 bit address. This statement does not result in an address being stored in ROM. The ADDRESS can be a decimal, octal or computed number:

- f. LABEL-kB
- g. LABEL+d
- h. LABEL-d
- i. LABEL

Decimal number, d, in the range 0 to 4095

Octal number, kB, in the range 0B to 7777B, where B signifies octal

Computed number, c, which is within the decimal or octal range, according to whether it is computed from octal or decimal values, of the form:

- a. *+kB
- b. *-kB
- c. *+d
- d. *-d
- e. LABEL+kB

where * means "this address" and LABEL means a micro-instruction label that is defined in the micro-program before this statement.

Examples of EQU statements:

Character Column:		1	10	30
Fields:		Field 1	Field 2	Field 6
Content:		HALT RELO START	EQU EQU EQU	400B 6000B RELO

4-24. ONES

Character Column:		1	10	15	20	25	30	40	80
Fields:		Field 1	Field 2	Field 3	Field 4	Field 5	Field 6	Field 7	
Content:		LABEL	"ONES"	(blank)	(blank)	(blank)	(blank)	COMMENTS	

The ONES statement creates a 24 bit micro-instruction word in ROM consisting of ones in all 24 bits.

Example of a ONES statement:

Character Column:		1	10
Fields:		Field 1	Field 2
Content:		NEG 1	ONES

4-25. SKP

Character Column:		1	10	15	20	25	30	40	80
Fields:	Field 1	Field 2	Field 3	Field 4	Field 5	Field 6	Field 7		
Content:	(blank)	"SKP"	(blank)	(blank)	(blank)	(blank)	(blank)	COMMENTS	

The SKP statement commands the micro-assembler to skip to the Top of the next page (TOP OF FORM command) during the listing of the microprogram. No locations in ROM are used, when this statement is specified.

Example of a SKP statement:

Character Column:		1	10
Fields:	Field 1	Field 2	
Content:		SKP	

4-26. ZEROES

Character Column:		1	10	15	20	25	30	40	80
Fields:	Field 1	Field 2	Field 3	Field 4	Field 5	Field 6	Field 7		
Content:	LABEL	"ZEROS"	(blank)	(blank)	(blank)	(blank)	(blank)	COMMENTS	

The ZEROES statement creates a 24 bit micro-instruction word in ROM consisting of zeroes in all 24 bits.

Example of a ZEROES statement:

Character Column:		1	10	40
Fields:	Field 1	Field 2	Field 7	
Content:	NULL	ZEROS	NO BITS	

MICROPROGRAMMING SOFTWARE

SECTION

V

Two sets of programs are provided to assemble, debug, and implement microprograms. One set operates in the BCS (Basic Control System) environment and the other operates in the DOS-III (Disc Operating System) environment.

5-1. MICROPROGRAMMING SOFTWARE SUMMARY

The following microprogramming software is provided:

- A two-pass micro-assembler, which converts the user's source microprogram record into an object tape and microcode listing.
- A Micro Debug Editor, which reads the object tape into Main Memory, outputs it to Writable Control Store (WCS), and allows the user to run the microprogram in WCS. The user can set breakpoints, change micro-instructions, change registers, etc. This program also provides the ability to punch the paper tapes that are used to create ("burn") programs into the ROM.
- A WCS I/O Utility subroutine, callable from FORTRAN and ALGOL libraries, that allows a microprogram, stored in a regular FORTRAN, ALGOL, or Assembler program buffer (in Main Memory), to be written into WCS.

5-2. MICRO-ASSEMBLER

The Micro-assembler accepts 80-character fixed-field card format records from a card reader, paper tape reader, or disc (using the DOS-III JFILE directive). Each record contains one micro-instruction coded in mnemonic format as described in Section IV of this manual. The micro-assembler processes input records and produces an object program paper tape which contains micro-instructions in binary format. Optionally output is a microprogram listing in both mnemonic and binary format, a symbol table, and error messages.

5-3. HARDWARE ENVIRONMENT

The BCS version requires the following as the minimum hardware:

- a. An HP 2105 or HP 2108 Processor with 8K of Main Memory.
- b. A Teleprinter.

This minimum system means that the assembly of the microprogram will be slow, since all input, listing, and punching must take place on the teleprinter.

The following additional hardware is supported:

- a. Paper Tape Reader for source microprogram input.
- b. Paper Tape Punch for binary object tape output.
- c. Card Reader for source microprogram input.
- d. Line Printer for microprogram assembly listing and symbol table listing.
- e. 7970 or 3030 Magnetic Tape Unit for temporary storage of source microprogram that is input to Pass 2 of the micro-assembler.

The DOS-III version of the micro-assembler requires the same hardware as the DOS-III system.

5-4. MICRO-INSTRUCTION SOURCE RECORD

A micro-instruction source record has the following characteristics:

- a. Length ≤ 80 characters.
- b. If not on a punched card, terminated by RETURN and LINE FEED.
- c. Seven fields with the starting column of each field as follows:

Field Number	Character Column
1	1
2	10
3	15
4	20
5	25
6	30
7	40

Figure 5-1 shows a card record.

Refer to Section IV, "Microprogramming Language," for a description of the micro-orders appropriate to the seven fields.

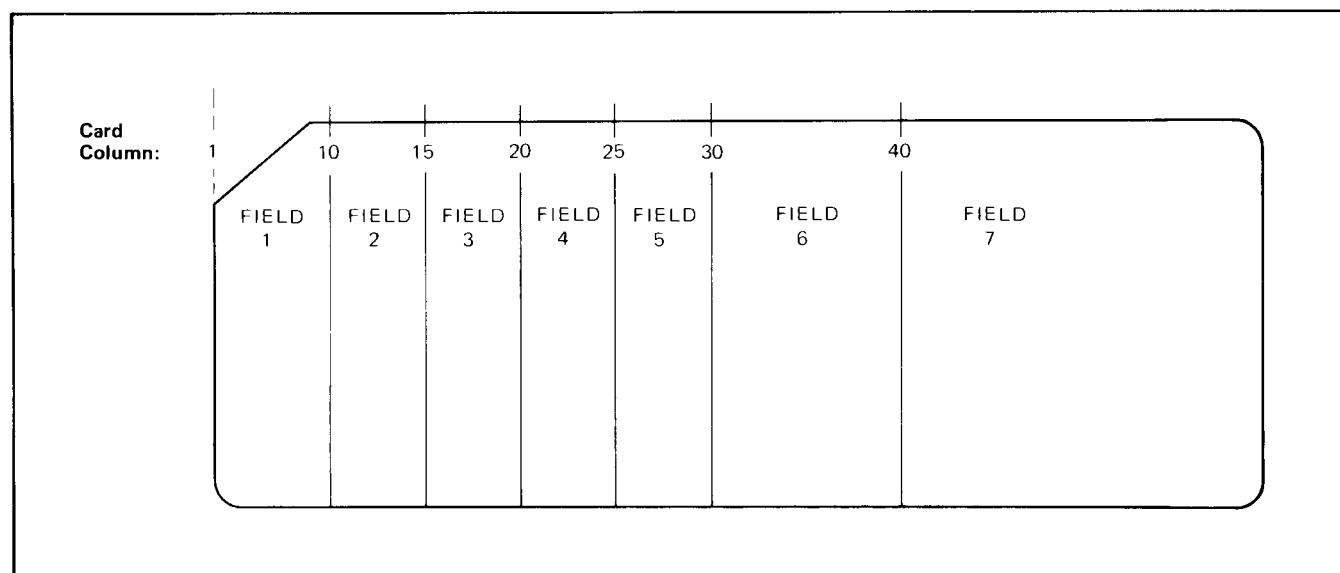


Figure 5-1. Micro-instruction Card Source Record

5-5. MICRO-ASSEMBLER CONTROL RECORD

Control statements are interspersed with micro-assembler language statements and specify control over the assembly process. For example, they may define the logical unit number of an input or output device or suppress listings.

There is one control statement per Control Record. If not on a card, it must be terminated by RETURN and LINE FEED.

Two control statements are required for every microprogram:

- a. \$ORIGIN statement
- b. \$END statement

All control statements start with a "\$" (Dollar character) in column 1. No intervening spaces are allowed in any control statement other than as specified. Details on each statement text and meaning are given below.

\$END

General Form: \$END

Meaning: End of microprogram

Purpose: Required as the last statement in every microprogram

Example: \$END

\$EXTERNALS

General Form: \$EXTERNALS = name1baddress1,
bnamebaddress2,
b. . .namenbaddressn

A comma and a space (*b*) separate each external name and address pair. Each "name" conforms to the Label definition in Section 4-1 and "address" means an octal address in the range 0 - 7777.

Meaning: Define the following label names:
name1 refers to address1
name2 refers to address2
:
:
:
namen refers to addressn

Purpose: Each \$EXTERNALS control statement provides for one or more branch (JMP or JSB) target addresses outside of the microprogram.

Example: \$EXTERNALS = OUTPUT 1012,
CHAR 736.

\$FILE

(Used by DOS-III systems only)

General Form: \$FILE = filename

The filename must be in accordance with DOS-III file name requirements.

Meaning: The object output file name for this microprogram is "filename."

Purpose: Provides the DOS-III micro-assembler with the name of the disc file into which the binary object code is to be stored.

Example: \$FILE=MOBJ

Note: Prior to assembling a microprogram with a \$FILE control statement, the user must have reserved a disc file using the DOS-III ":ST,B,..." directive.

\$INPUT

(Used by BCS systems only)

General Form: \$INPUT = lun

The logical unit number, lun, must be octal and in the range 1 - 74.

Meaning: The logical unit number of the device through which all subsequent input (to the next \$END statement) is to be read is "lun."

Purpose: When the assembly process is begun in BCS systems, the micro-assembler expects the first source statement to be entered through the system console device (logical unit number 5). The user may enter the whole source program through the system console device. Normally, however, the user enters a \$INPUT command specifying the logical unit number of the card reader or paper tape reader from which the rest of the source program is to be read.

Example: \$INPUT = 12

\$LIST**General Form:** \$LIST = lun

The logical unit number, lun, must be octal and in the range 1 - 74.

Meaning: The logical unit number of the listing device is "lun".

Purpose: To cause the assembly listing to be printed on the device having the specified unit number. If omitted, logical unit number is assumed to be 6 (standard list device).

Example: \$LIST = 16

\$ORIGIN**General Form:** \$ORIGIN = nnn

The origin, nnn, must be octal and in the range 0 - 7777.

Meaning: Set microprogram origin at octal address nnn in Control Store.

Purpose: Every microprogram must have its program address origin defined. New origins may be specified within the microprogram.

Example: \$ORIGIN = 427

\$RCASE**General Form:** \$RCASE

Meaning: Punch a special 32-micro-instructions/record object tape.

Purpose: This special object tape is reserved for system maintenance. Refer to Section 5-6 Micro-Assembler Output for a description of this special object tape.

Example: \$RCASE

\$OUTPUT**General Form:** \$OUTPUT = lun

The logical unit number, lun, must be octal and in the range 1 - 74. This statement may come anywhere before the \$END statement.

Meaning: lun is the logical unit number of the output device.

Purpose: To specify the device on which the micro-assembler object code is to be output. If this statement is omitted, logical unit of 4 is assumed.

Example: \$OUTPUT = 10

\$PASS 2

(Used by BCS systems only)

General Form: \$PASS2 = lun

The logical unit number, lun, must be octal and in the range 1 - 74. If present, this must be the first statement in the source deck or tape.

Meaning: lun is the logical unit number of the magnetic tape unit onto which all subsequent micro-assembler input is to be written.

Purpose: To cause all source input to be recorded on magnetic tape for use as input to Pass 2 of the micro-assembler. If this control statement is omitted, the computer halts at the end of Pass 1 to allow the operator to reload the microprogram source into the "\$INPUT" device.

Note: The only magnetic tape units supported by the micro-assembler are the HP 3030 and HP 7970.

Example: \$PASS2 = 23

\$SUPPRESS**General Form:** \$SUPPRESS**Meaning:** Suppress all warning error messages.**Purpose:** To cut down the volume of messages to the console device. Fatal error messages will still be printed.**Example:** \$SUPPRESS**\$SYMTAB****General Form:** \$SYMTAB**Meaning:** Print symbol table**Purpose:** To provide the user with label names and corresponding octal addresses used in his microprogram.**Example:** \$SYMTAB**5-6. MICRO-ASSEMBLER OUTPUT**

This section describes all forms of output from the micro-assembler. They are:

- Binary Object
- Symbol Table
- Source and Binary Microprogram Listing
- Error Messages

5-7. BINARY OBJECT OUTPUT

The Standard Object Tape output by the micro-assembler to paper tape or a disc file consists of one or more Instruction Records, the format of which is shown in Appendix A, Figure A-1. One Instruction Record holds up to 27 micro-instructions and five words of header information. Each micro-instruction requires 32 bits or two words in the format: an eight bit address and 24 bits for the micro-instruction. Hence the length of the record =

5 words of header

2n words for n micro-instructions (2 words for each micro-instruction)

5+2n words for one Instruction Record

No more than 27 micro-instructions are written into an Instruction Record. Hence the maximum length = $5+(2 \times 27)=59$ words. The last object record is a four word End Record. When the microprogram consists of more than 27 micro-instructions, a series of Instruction Records are produced with the last one holding 27 or less micro-instructions. For example, if 57 micro-instructions have been assembled, three Instruction Records and an End Record are required consisting of the following:

- Instruction Record 1 holds 27 micro-instructions and consists of
 - 5 words of header
 - 54 words for 27 micro-instructions
 - 59 words
- Instruction Record 2 holds 27 micro-instructions and consists of
 - 5 words of header
 - 54 words for 27 micro-instructions
 - 59 words
- Instruction Record 3 holds 3 micro-instructions and consists of
 - 5 words of header
 - 6 words for 3 micro-instructions
 - 11 words
- The End Record consists of
 - 4 words
 - 133 words for the entire microprogram Binary Object.

The Standard Object format is accepted by all programs which accept standard relocatable format. Thus a Standard Object tape can be stored in a DOS-III file using the “:STORE,R,...” directive. However, if the DOS-III user wants the Binary Object stored automatically in a disc file by the micro-assembler, the DOS-III directive “STORE,B,...” must have previously been used to reserve a disc file.

The Micro-assembler can also produce a non-standard object as the result of the inclusion of the \$RCASE control statement. This optional object is the HP ROM Simulator Object tape. The format of this tape is shown in Appendix A, Figure A-2.

5-8. SYMBOL TABLE LISTING

If the user has a \$SYMTAB control statement in his microprogram source input, then the micro-assembler will print a symbol table on the device with logical unit number 6 or on the device defined by the \$LIST control statement, if present.

An example of a symbol table is shown in Figure 5-2.

On the left are the symbols or labels in the microprogram. On the right is the value of the symbol; that is the six digit absolute octal address of the symbol. Where X follows the address, the symbol has been defined by a \$EXTERNAL control statement.

SYMBOL TABLE	
MOVE	002412X
GOTO	003421X
RET	002427X
LAST	002717X
OUT	002011
ERR1	002012

Figure 5-2. Symbol Table

5-9. MICROASSEMBLY LISTING

Unless suppressed by the \$NOLIST control statement, the micro-assembler provides a listing like the one shown in Figure 5-3. This listing is associated with the symbol table illustrated in Figure 5-2.

5-10. MICRO-ASSEMBLER ERROR MESSAGES

During the assembly process the micro-assembler checks each instruction for errors. If an error is detected, an error message of the following general form is printed in the Micro-assembly Listing.

****ERROR eeee IN LINE nnnn**

where

eeee

is an Error Code defined in Table 5-1 and

nnnn

is a line number in the Micro-assembly Listing.

Table 5-1 gives the meaning of each error code and the recovery procedure. Note that Figure 5-2 holds examples of two error messages in lines 9 and 11.

5-11. DOS-III OPERATION OF MICRO-ASSEMBLER

Before using the DOS-III version of the Micro-assembler, the following items must be available.

- A current DOS-III system.
- A source microprogram, on cards, paper tape, or in a source file on disc.

0001	\$ORIGIN=2000R FIRST ADDRESS OF MODULE 4									
0002	\$SYMTAB PRINT SYMBOL TABLE									
0003	\$EXTERNAL=MOVE 2412, GOTO 3421, RET 2427, LAST 2717									
0004	* P2=AKP1									
0005	2000	220	074457	READ	INC	M	P	READ ADDEND P		
0006	2001	017	126157		PASS	L	A	PUT AUGEND IN L AND ENABLE E & O		
0007	2002	264	101557	FNVE	ADD	S12	TAB	ADD MEMORY TO L AND STORE IN S12		
0008	2003	324	140531	JMP	CNDX	E	ERR1	IF E SET, GO TO ERR1		
**ERROR 0008 IN LINE 0009										
0009	2004	320	000030	JMP	CNDX	OVFL	ERR2	IF O SET, GO TO ERR2		
0010	2005	000	075717		INC	P	P	BUMP P FOR NEXT PARAMETER		
**ERROR 0003 IN LINE 0011										
0011	2006	017	136757	READ	INC	M	P	READ NEXT PARAMETER P2 ADDRESS		
0012	2007	000	000461		MPCK	INC	M	TAB		
0013	2010	177	166017	WRTE	PASS	TAB	S12	PUT IN M AND CHECK FOR M P ERR		
0014	2011	017	136776	OUT	RTN			PUT ADD RESULT INTO MEM ADD P2		
0015	2012	344	001757	ERR1	IMM	LOW	S	0		
0016	2013	320	100470		JMP			THE RETURN		
0017	2014	340	001757	ERR2	IMM	HIGH	S	0		
0018	2015	320	100470		JMP			SET UPPER BYTE FOR E ERR		
0019				\$END				RETURN		
** 0002 ERRORS**										
Line	ROM	Bits	Bits	Field	Field	Field	Field	Field	Field	
Number	Address	23-16	15-0	1	2	3	4	5	6	
Binary				Field						
Micro-instruction				7						

Figure 5-3. Micro-Assembly Listing

- c. The Micro-assembler program named MICRO stored in the DOS-III user library. If MICRO still is on relocatable object paper tape (HP 12978-160001), it can be loaded in the same way as any other relocatable object program.

- a. If there is a \$FILE control statement in the microprogram source, a binary file must be reserved on the disc before beginning the micro-assembly process to hold the relocatable object. The name of the reserved disc file must be the same as the one specified in the \$FILE control statement.

For the detailed description of DOS-III operation, see **HP 24307B DOS-III Reference Manual** (HP 24307-90006).

- b. Place the microprogram source in the input device; turn the device on; turn on the paper tape punch and the list device.

Table 5-1. Micro-assembly Error Messages

Error Code	Meaning/Recovery
1	Duplicate Label. The statement label of the micro-instruction in line nnnn is the same as another statement in the microprogram or the same as a declared \$EXTERNAL symbol. Assign a new statement label and reassemble.
2	Illegal Control Statement. Correct control statement in line nnnn and reassemble.
3	Illegal Field 2 Micro-order. A NOP is inserted in field 2 and assembly continues. Correct line nnnn and reassemble.
4	Illegal Field 3 Micro-order. A NOP is inserted in field 3 and assembly continues. Correct line nnnn and reassemble.
5	Illegal Field 4 Micro-order. A NOP is inserted in field 4 and assembly continues. Correct line nnnn and reassemble.
6	Illegal Field 5 Micro-order. A NOP is inserted in field 5 and assembly continues. Correct line nnnn and reassemble.
7	Illegal Field 6 Micro-order. A NOP is inserted in field 6 and assembly continues. Correct line nnnn and reassemble.
8	Illegal JMP or JSB Address. Address is outside permitted range, or target label address is undefined. A value of 0 will be inserted into address field of line nnnn and assembly continues. Redefine address and reassemble.
9	Microprogram Too Large. The last relative address in the program is 400 or greater. A \$ORIGIN statement must be changed or the program broken up into smaller parts before reassembly.
10	Missing \$ORIGIN Control Statement. At least one \$ORIGIN control statement is required. Insert \$ORIGIN statement and reassemble.
11	Illegal Word Type 2 Operand. Operand of the IMM micro-instruction is outside the permitted range. A value of 0 is inserted into the operand and assembly continues. Correct line nnnn and reassemble.
OR aaaa	Insufficient DOS-III File Space Reserved. Reserve a binary file with more sectors for storage of the file named in the \$FILE control statement (aaaa is an address in the micro-assembler and can be disregarded). See DOS-III manual section 15 under Error Conditions.
ABORT!	An irrecoverable error has occurred; correct error and reassemble.

- c. Summon the Micro-assembler with statement

```
:PR,MICRO,[p1,p2,p3,p4,99]
```

where

p1 = the input device logical unit number

p2 = list device logical unit number

p3 = paper tape punch device logical unit number

p4 = maximum number of lines-per-page on the list device.

If 99 is entered for any of the above parameters, that parameter and all those that follow are defaulted to "standard" values.

- d. The program title

MICRO-ASSEMBLER

is printed and Pass 1 begins. If a \$SYMTAB control statement is in the source microprogram, the symbol table is printed at the conclusion of Pass 1. Pass 2 begins immediately (from disc) and the listing and relocatable object tape are output. Micro-assembly is complete.

Note: If Pass 2 fails to begin, check that the paper tape punch is turned on. The micro-assembler will cycle in a loop until the punch is turned on.

5-12. BCS OPERATION OF MICRO-ASSEMBLER

Before proceeding, the following items must be available:

- An absolute BCS binary tape.
- A relocatable object tape of the Micro-assembler program MICRO (HP 12978-160003).
- A source microprogram either on cards or paper tape.

For a detailed description of BCS usage, see the **Basic Control System** manual (HP 02116-9017).

The following procedure need be performed only once. When an absolute binary tape of the Micro-assembler is punched, it is used as described in the procedure "Executing the Micro-assembler."

Making an Absolute Micro-assembler tape:

- a. Load the absolute BCS binary tape using the Basic Binary Loader.
- b. Set the P-register to 2. Set bit 14 of the Switch Register and clear all other Switch Register bits.

- c. Place the MICRO relocatable object tape in the paper tape reader. Check that the paper tape reader and the console device are on. Turn on the paper tape punch. Press PRESET and RUN on the CPU front panel. MICRO reads in and absolute binary tape is punched.

- d. The message

*LOAD

is printed and the computer waits. Set Switch Register bits 2 and 14 leaving all others clear. Load BCS Library tape into the paper tape reader. Press RUN.

- e. The BCS Library tape reads in and the rest of the absolute binary tape is punched. Linkage information is printed on the console device.

This is the absolute binary tape of MICRO, used for input to the next step.

Executing the Micro-assembler:

- a. Load the MICRO absolute binary tape using the Basic Binary Loader.
- b. When loading is complete, set P-register to 2. Press PRESET and RUN. The message

MICRO-ASSEMBLER

is printed followed by a request for the logical unit number of the source input device.

- c. Enter the logical unit number followed by carriage return/line feed. Pass 1 now begins. If a \$SYMTAB control statement is in the microprogram source, the symbol table is printed at the conclusion of Pass 1. (See Section 5-5 for a description of the \$SYMTAB control statement.)

- d. Turn on the paper tape punch.

- e. Pass 2 begins immediately. If no \$PASS2 control statement was included in the source, the message

RELOAD SOURCE, PRESS RUN

is printed. Reload the source microprogram into the input device and then press RUN on the front panel of the computer.

Note: If Pass 2 fails to begin, check that the paper tape punch is turned on. The micro-assembler will cycle in a loop until the punch is turned on.

If a teletype is used for both listing and punching, the computer halts (M-register = 102052) so that the operator can press the paper tape punch ON button to punch the microprogram object tape. The operator then presses RUN on the computer front panel.

When the paper tape is punched, another halt (M-register = 102053) occurs, so that the paper tape punch button can be set to OFF. Press RUN on the computer front panel.

- f. Pass 2 completes micro-assembly. The microprogram object tape is complete. To assemble another microprogram proceed from step b.

5-13. MICRO DEBUG EDITOR

The Micro Debug Editor (MDE) makes it possible to load the object microprograms output from the Micro-assembler into a Writable Control Store module. It also provides the ability to debug microcode stored in the WCS and to "burn" microprograms into ROM chips.

Before using the Micro Debug Editor to debug microprograms, the Writable Control Store PCAs must be set to the required control store module numbers. This is accomplished by the installation of a module selection Jumper Assembly (HP Part Number 5060-8342). Refer to Section 6 of this manual for installation of the module selection Jumper Assembly and the WCS PCAs.

5-14. HARDWARE ENVIRONMENT

The BCS version requires the following minimum hardware:

- a. HP 21MX Series Computer with 8K of Main Memory
- b. A console device
- c. A paper tape reader
- d. One or more WCS PCA's, depending on the size of the microprogram to be debugged.
- e. If a ROM program tape is to be punched, a paper tape punch is also required.

The DOS-III version of the MDE requires the same minimum hardware as the DOS-III system.

5-15. INITIALIZATION PROGRAM

When the Micro Debug Editor is to be run for debugging purposes (as opposed to being run merely to punch ROM program tapes), the user must supply an initialization program. The initialization program is an assembly language program that prepares the necessary parameters in

ASMB,R,B,L,T	Assembly parameters
NAM TEXT,6	Program name (DOS-III)
ENT TEST,MACRO	Entry points
TEST NOP	
.	Any initialization procedure re-
.	quired by the microprogram
.	
MACRO OCT 105xxx	(or 101xxx) Instruction that calls
	the user microprogram
DEF P1	
DEF P2	Parameter addresses required by
.	the microprogram
.	
.	
DEF Px	
JMP TEST,I	Return to calling program (MDE)
P1 (parameter 1 value)	
P2 (parameter 2 value)	Parameter values
.	
.	
.	
Px (parameter x value)	
END	

Figure 5-4. General Format of the Initialization Program

Main Memory and then executes a 101xxx or 105xxx macro-instruction.

The name of the initialization program must be TEST (required in BCS systems, is a NAM TEST statement; in DOS-III systems a NAM TEST, 6 statement). The program must also have the symbol "MACRO" declared as an entry point where MACRO is the symbolic address (label) of the macro-instruction (101xxx or 105xxx) which calls the microprogram under test. Note that there must only be one such macro-instruction in the TEST initialization program.

Figure 5-4 holds the general structure of the initialization program.

This initialization program is called as a relocatable sub-routine by MDE. Thus, its name is one of the references that must be satisfied when loading MDE.

A note of caution: a microprogram cannot be debugged using MDE unless the microprogram has:

- An entry point which is a "JMP" micro-instruction of Word Type 4 (described in Section 4-17).
- The micro-instruction jumped to by the JMP at the entry point must not contain a "READ" micro-order.

An example of a short initialization program is shown in Figure 5-5.

5-16. USING THE MICRO DEBUG EDITOR

Section 5-37 describes how to execute MDE using the DOS-III operating system. Section 5-38 describes how to execute MDE using the BCS operating system.

Before using the Micro Debug Editor to debug a microprogram, the Writable Control Store PCAs must have the correct terminal board plugged in, to establish the Control Store module number. Refer to Section VI of this manual for a description of setting module numbers in a Writable Control Store PCA.

When the module number has been set in the Writable Control Store PCA and it is plugged into the correct I/O slot, the user loads the microprogram object tape (produced by the Micro-assembler) using the Micro Debug Editor LOAD command. The microprogram is then output to the Writable Control Store using the WRITE command.

When the user is ready to execute his microprogram, the EXECUTE command is used. For the microprogram to execute properly, the following conditions must hold:

- The module that the microprogram was written into matches the range of addresses used by the microprogram. For example, a microprogram whose addresses are in the octal range 2400 to 2777 must be stored in a Writable Control Store PCA which has been set to module 5.

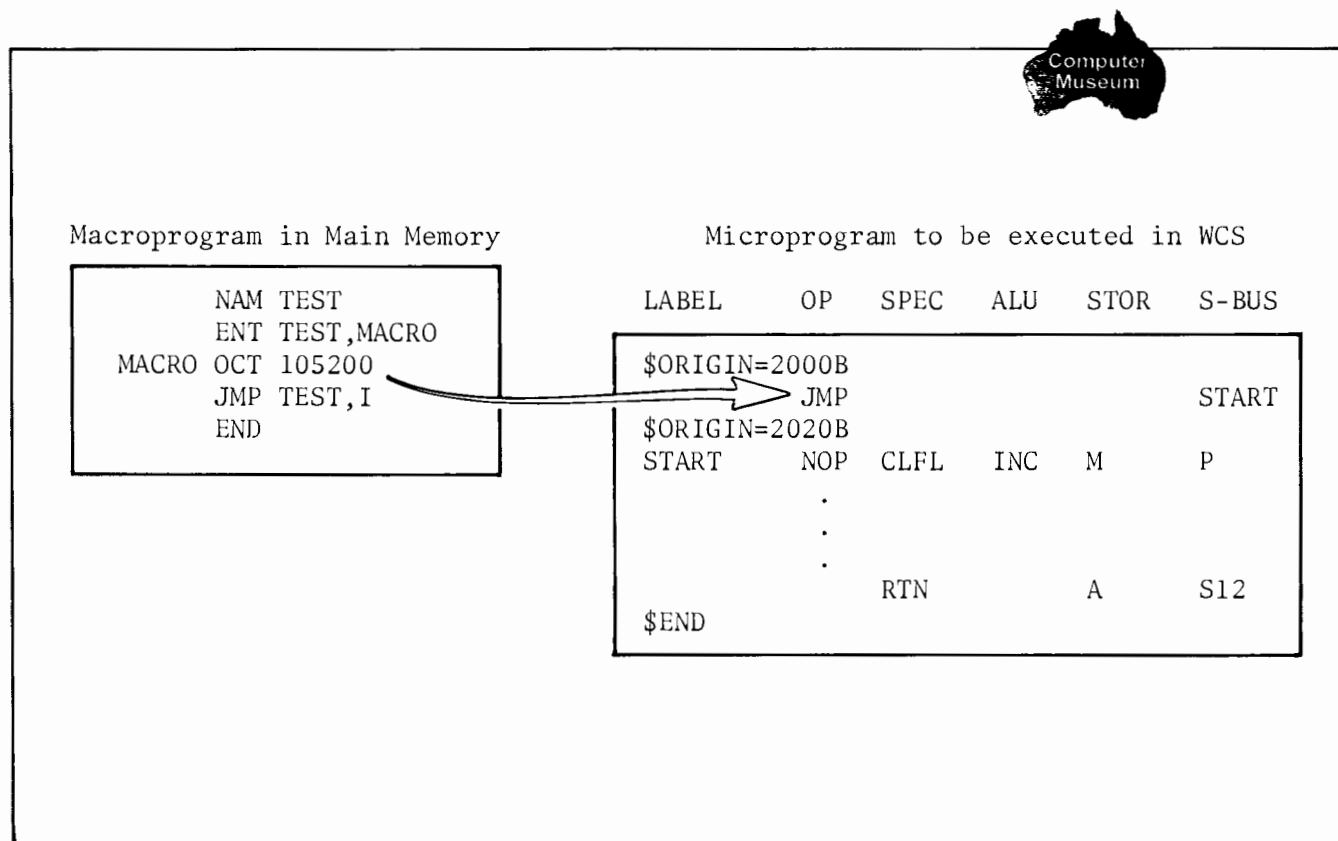


Figure 5-5. Test Program Call to Microprogram

- b. The macro-instruction in the TEST program must initiate entry into Control Store at the proper address of the microprogram to be tested.

Micro Debug Editor results are unpredictable if either of the above conditions are not met.

When MDE is executed, it prints the input prompt

COMMAND?

on the system teleprinter.

Respond by entering one of the input, edit, output, or debug commands described in Table 5-2 and the following pages. In most cases, the first letter of the command is sufficient to specify it to MDE. The two commands, "MOVE" and "MODIFY", require at least three letters to identify the command. After MDE has performed the specified operation, it again prints COMMAND? to repeat the cycle.

Terminate an MDE run by entering the FINISH command.

There are 13 MDE commands which are summarized in Table 5-2. A detailed description of each command follows. Whenever a logical unit number (lun) is called for, it must be entered in octal.

Note that the last octal 45 words of the lowest numbered WCS module loaded with a microprogram are used by Micro Debug Editor for its own resident microcode. If these locations are required by the user microprogram under test, use the MOVE command to relocate the MDE microcode before loading the user microprogram.

The Micro Debug Editor uses a Main Memory buffer to hold the microprogram object code. When the microprogram is loaded from an object tape, it is stored into this buffer. Most MDE commands make modifications or transfers to and from this buffer.

Use of the PREPARE command to punch the six ROM microprogram mask tapes has the following restriction. This buffer must have been loaded using an object tape produced by the micro-assembler and the buffer must not have been modified.

5-17. INPUT COMMANDS

5-18. LOAD[,X]

Meaning: Load the object microprogram produced by the Micro-assembler from disc or paper tape into the MDE buffer. The logical unit number (lun) of the input device is X.

Usage: The Micro-assembler control statement \$FILE can be used to specify (during assembly) the name of the DOS-III file into which the object code is to be stored. In the DOS-III version of MDE, if the logical unit number

entered is that of the disc, MDE will respond with a request for the name of the file in which the object code is to be stored:

FILENAME?

Enter the file name given to the object code by the \$FILE control statement.

Note: When loading the object microprogram for output to WCS (instead of punching ROM tapes), the LOAD command must be followed immediately by a WRITE command to the appropriate WCS PCA. No intervening commands are allowed. This allows the Micro Debug Editor to build a table relating microprogram addresses to WCS logical unit numbers.

Table 5-2. Micro Debug Editor Commands

INPUT

Commands: LOAD[,X]
READ,X

EDIT

Commands: SHOW,xxxx[,yyyy]
MODIFY,xxxx[,yyyy]

OUTPUT

Commands: DUMP[,X]
WRITE,X

PREPARE[,X]
VERIFY[,X]

TERMINATION

Command: FINISH

DEBUG

Commands: BREAK,yyyy
CHANGE[,mnemonic]
EXECUTE[,0 or yyyy]

RELOCATE MDE WCS-RESIDENT MICROCODE

Command: MOVE,yyyy

Note

The brackets indicate that the parameter may be omitted.

5-19. READ,X

Meaning: Read the contents of a WCS into the Micro Debug Editor buffer. **X** is the logical unit number of the WCS.

Usage: If no WCS is on the specified logical unit, the MDE buffer is unchanged. No notification is made to the user that the buffer is unchanged or that no WCS is on the logical unit specified. Thus, if READ or SHOW is being used to insure that a previous WRITE executed properly to the same (non-WCS) logical unit, the MDE buffer will still hold the data that was assumed to be written to that logical unit. The user could incorrectly assume that the non-existent WCS holds the proper data.

5-20. EDIT COMMANDS**5-21. SHOW,xxxx[,yyyy]**

Meaning: Display the WCS contents on the console device, where **xxxx** is the beginning address and **yyyy** is the ending address. Only the contents of the address **xxxx** are displayed, if **yyyy** is omitted.

Usage: See Usage under 5-19, READ,X.

The display format of each 24-bit word is:

```
aaa mmm nnnnnn
```

where **aaa** is the control store address of the location being displayed, **mmm** is the octal representation of bits 23-16 of the location, and **nnnnnn** is the octal representation of bits 15-0 of the location.

5-22. MODIFY,xxxx[,yyyy]

Meaning: Change the contents of the MDE buffer and the WCS where **xxxx** is the beginning WCS address and **yyyy** is the ending WCS address. Change WCS address **xxxx** if **yyyy** is omitted.

Usage: See Usage under 5-25, WRITE X.

“MOD” is the minimum input required to initiate the modify command. **xxxx** and **yyyy** must be absolute WCS addresses in a single WCS module. One at a time, the contents of each location are printed on the console device in the same format as the SHOW command above. Following the location contents, the operator enters the new location contents followed by a CARRIAGE RETURN and LINE FEED.

If fewer than 3 digits are entered for **mmm** or fewer than 6 digits are entered for **nnnnnn**, the number entered is right justified with zeros automatically filled to the left. To specify that no change is to be made, enter an asterisk (*), instead of **mmm** or **nnnnnn**.

Example:

```
MOD,4000,4003
4000 123 456777 *,123456
```

leaves bits 23-16 unchanged and sets bits 15-0 to 123456 in WCS location 4000.

```
4001 123 456777 6,123
```

is equivalent to entering 006,000123; bits 23-16 are set to 006 and bits 15-0 are set to 000123 in location 4001.

```
4002 123 456777 123, *
```

sets bits 23-16 to 123 and leaves bits 15-0 unchanged in location 4002.

```
4003 123 456777 *, *
```

makes no change to location 4003.

5-23. OUTPUT COMMANDS**5-24. DUMP[,X]**

Meaning: Punch the entire contents of the MDE buffer on the paper tape punch. **X** is the logical unit number of the paper tape punch. If **X** is omitted, it is assumed to be 4.

Usage: The DUMP command must be preceded by a READ or LOAD command to fill the MDE buffer. The tape produced is in the same format as the object tape produced by the Micro-assembler. If the tape is reloaded into the MDE buffer, the buffer cannot be used to punch (PREPARE command) a set of six pROM mask tapes. The primary use of this tape is to enable the user to save the results of a microprogram debug session for resumption later.

5-25. WRITE,X

Meaning: Write the contents of the MDE buffer into the WCS. **X** is the logical unit number of the WCS.

Usage: Since the Micro Debug Editor addresses the WCS by logical unit number, it is the responsibility of the user to insure that a WCS is installed with logical unit number **X** and that it is set to the proper module for the micro-code to be stored. If no WCS is on the specified logical unit, no notification is given to the user that a WRITE or MODIFY command failed to transmit data to the non-existent WCS.

5-26. PREPARE[,X]

Meaning: Punch a set of six pROM mask tapes each headed by three lines of I.D. and a checksum on the paper tape punch. **X** is the logical unit number of the device. If **X** is omitted, it is assumed to be 4.

Usage: Following entry of the PREPARE command, a cycle of dialogue is initiated between the operator and the console device. In the following procedure, the underlined characters indicate operator input is required at the console device. Each entry must be followed by a CARRIAGE RETURN and LINE FEED.

- a. Turn on the paper tape punch. The message cycle starts with:

GENERATION OF MASK BITS 23-20

where 23-20 represents the 4 bit range of bits to be punched into the first mask tape.

ENTER 3 LINES OF I.D. INFORMATION

LINE 1 — key in first line of tape I.D.

LINE 2 — key in second line of tape I.D.

LINE 3 — key in third line of tape I.D.

Enter up to 72 characters of identification information in each line.

- b. Following entry of the third I.D. line, the mask tape is punched for mask bits 23 to 20. This is for ROM chip number 6. The following cycle of dialogue is repeated for each of the remaining five mask tapes:

GENERATION OF MASK BITS UU-LL

UU - LL is the range of bits to be punched.

ANY CHANGE OF I.D. INFO IN LINE 1? key in N (no) or Y(yes) and new line 1 I.D.

LINE 2? key in N or Y and new line 2 I.D.

LINE 3? key in N or Y and new line 3 I.D.

- c. The next mask tape is punched. When all six mask tapes have been punched, the following message is output:

GENERATION OF TAPES COMPLETED

The six mask tapes have the following characteristics:

UU-LL	Punch Sequence	For Module ROM Chip No.
23-20	First tape	6
19-16	Second tape	5
15-12	Third tape	4
11-08	Fourth tape	3
07-04	Fifth tape	2
03-00	Sixth tape	1

Conventions: Line 1 I.D. holds module number, ROM chip number, number of bits (4), ROM size, and other I.D. information.

For example:

LINE 1-1,005, 4, 1025 REENTRY FACTOR

Line 2 I.D. holds part number or other central reference number. For example:

LINE 2-MT 38-0226 REVISION C

Line 3 I.D. holds date and any other I.D. information. For example:

LINE 3-04/01/75 PVT. D.M. BULMAN

5-27. VERIFY[,X]

Meaning: Compare the contents of the pROM mask tapes to the contents of the MDE buffer. The logical unit number of the paper tape reader is X.

Usage: Following entry of the command, the console device requests the range of bits in the pROM mask tape to be compared to the MDE buffer (underlined characters indicate operator entry).

TAPE NUMBER: uull

Enter CARRIAGE RETURN and LINE FEED after the bit range uu (upperlimit) and ll (lowerlimit). Refer to 5-26 PREPARE[,X] for valid bit ranges.

For example, the entry "2320" specifies verification of bits 23 to 20. The paper tape then reads the mask tape and compares its contents to the specified bits in the MDE buffer. As the tape is being read, the three lines of I.D. (see PREPARE command) and checksum are printed on the console device.

Note: If the DOS-III operating system is being used, and no errors were encountered, an I/O "error" message is printed at the console device:

I/O ERR ET EQT #n

Where n is the EQT number of the paper tape reader. This message notes a characteristic of the mask tape that DOS-III normally interprets as an error condition, but the message in fact, connotes no error.

If no errors were detected, the message

TAPE VERIFIED

is printed. Enter another bit range as before. The VERIFY command completes only after the bit range 03 to 00 has been entered and verified.

Errors: If errors are detected, dialogue between the console device and the operator is initiated. Follow each operator entry with CARRIAGE RETURN and LINE FEED.

- a. The message CHECKSUM ERROR OR BAD MASK TAPE is printed followed by a tape repunch request:

DO YOU WANT TO REPUNCH THIS TAPE?
enter Y or N

- b. If N is entered, another bit range request with the message

TAPE NUMBER?

Enter another bit range as before. The VERIFY command completes only after the bit range 03 to 00 has been entered and verified.

- c. If Y is entered, the following request is made:

ENTER PUNCH LOGICAL UNIT # enter octal
logical unit number of paper tape punch

The message

ENTER THREE LINES OF I.D. INFORMATION

is printed.

Enter up to 3 lines of tape I.D. information according to the procedure given in 5-26, PREPARE[,X]. The new mask tape is punched, headed by the I.D. information.

Special DOS-III operation: When a series of bit ranges are being verified, specification of each successive range at the console device (as a result of the message TAPE NUMBER?) will bring about the prompt character "@". To verify the specified bit range on paper tape:

- a. Enter the following command

:UP,n

where n is the EQT number of the paper tape reader.

- b. Then enter:

:GO

The next tape to be verified will read in as above.

Verify sequence: The mask tapes may be verified in any order with exception that the last tape verified must have the bit range 03 to 00.

5-28. TERMINATION COMMAND

5-29. FINISH

Meaning: Terminate the current MDE run.

5-30. DEBUG COMMANDS

5-31. BREAK,yyyy

Meaning: Set a Breakpoint at location yyyy and clear the previous one. If yyyy = 0, no breakpoint is set and the previous one is cleared.

Usage: Microcode execution is initiated by an EXECUTE command. When the Breakpoint address yyyy is reached,

REG'S?

is printed and microprogram execution ceases (breaks). Enter the mnemonics of the flags or registers that are to be displayed, separated by commas. The mnemonics are described under the CHANGE command. The entry is of the form

REG'S? m1,m2,m3, . . . mn

where m1 through mn are register and flag mnemonics. The resulting display is of the form

m1 = c1, m2 = c2, m3 = c3,, mn = cn

when c1 through cn are octal contents of the requested registers and flags.

Example of a display request:

REG'S A,B,1,2,3,4,14

The resulting display:

A = 00004, B = 103005, 1 = 000447,
2 = 00012, 3 = 00000, 4 = 00000,
14 = 034716

Enter "!" to display all registers and flags. Enter "/" to return to command entry mode.

Restrictions: Do not set a breakpoint

- in the WCS entry point address of the microprogram
- in a microprogram subroutine (within the JSB . . . RTN code limits)
- in an address where the micro-instruction passes information to or from the T-register immediately following a WRITE or READ micro-order.

5-32. CHANGE[,m]

Meaning: Alter the contents of one or more registers and flags. If the mnemonic **m** is specified, alter the contents of the register or flag which it specifies. If not specified, all registers and flags are displayed in sequence to prompt the user to make required changes.

Mnemonics: The list of register and flag mnemonics follows:

Mnemonic	Stands For	Mnemonic	Stands For
A	A-register	9	S9-register
B	B-register	10	S10-register
S	S-register	11	S11-register
P	P-register	12	S12-register
1	*S1-register	X	X-register
2	S2-register	Y	Y-register
3	S3-register	O	Overflow Register bit
4	S4-register	E	Extend Register bit
5	S5-register	F	CPU Flag bit
6	S6-register	CN	Counter Register
7	S7-register	L	L-register
8	S8-register		

*Scratch Pad Register 1; similarly for S2, S3, etc.

Usage: Upon entry of the command, the message

m xxxxxx =

is printed, where **m** is the register or flag mnemonic and **xxxxxx** is the octal representation of the contents. Enter the new contents or an asterisk (*) if no change is to be made.

Example of a CHANGE request:

CHANGE,6
6 173777 = 173770

This is a request for a change to S6-register (Scratch Pad Register 6). The original contents were octal 173777. The new contents are octal 173770.

5-33. EXECUTE[,yyyy]

Meaning: Execute microprogram.

If **yyyy = 0**, the TEST initialization program is run, which carries execution to the microcode in WCS. This is the normal mode of initiating microcode.

Note: If the entire system goes dead after entering an EXECUTE,0, the reason may be that the WCS with the correct module number is not plugged into the correct slot.

If **yyyy =** an absolute WCS address, execution of microcode begins at that address.

If **yyyy** is omitted, execution resumes from the last breakpoint with registers and flags set

- according to their setting when the breakpoint was encountered, or
- modified by the CHANGE command.

Usage: Execution will continue until a breakpoint is encountered or until the microprogram is completed. When complete, the command entry mode is repeated.

Before initiating a microprogram execute (other than EXECUTE,0), make sure that all registers and flags are preset using the CHANGE command, if necessary.

5-34. RELOCATE MDE WCS-RESIDENT MICROCODE

5-35. MOVE,yyyy

Meaning: Move the octal 45 word WCS-resident microprogram portion of MDE from the usually resident locations to locations beginning with **yyyy**.

Usage: "MOV" is the minimum input required to initiate the move operation. MDE requires a portion of WCS for register dump and register restore microprograms. These MDE microprograms are initially stored in relative octal locations 333 to 377 of the first WCS loaded. If the user requires these locations in Writable Control Store, he can move this resident MDE microcode elsewhere.

No check is made to see if a portion of the user microcode has been overlayed. The reason is that the user may actually want to situate the dump and restore microprograms on top of his own microcode as he debugs another portion of his code.

The actual relocation of the MDE microcode does not occur until the EXECUTE command is given.

5-36. MDE ERROR MESSAGES

During the use of MDE, commands, parameters, and processing functions are monitored. If an error condition is detected, an appropriate message is printed. Table 5-3 holds the list of MDE error messages plus their meaning and the recovery procedure.

5-37. DOS-III OPERATION OF MDE

Before using the DOS-III version of the Micro Debug Editor (MDE), the following items must be available.

- A current DOS-III system
- A relocatable object tape of MDE (HP 12978-16002).

- c. A relocatable object tape of the TEST initialization program if a debug run is to be made.
- d. A microprogram object tape output by the Micro-assembler.

The following is an example of how the user can proceed. For details on additional DOS-III options, see DOS-III manual (HP 24307-90006).

- a. Store the two tapes, MDE and TEST, on the disc using the DOS-III store command

:ST,R,filename, lun

where **filename** is any suitable label and **lun** is the logical unit number of the paper tape reader from which the tapes are entered.

- b. Make sure the list device is on. At the console device enter

:PR,LOADR,2

DOS-III responds with

ENTER FILE NAMES OR /E

- c. Respond as follows:

MDE filename, TEST filename, /E

where **MDE filename** and **TEST filename** are the chosen file names used with the "ST" store command (step A), and /E specifies end of entry.

If MDE is being used only to load WCS with a microprogram, the **TEST filename** may be omitted. The loader then reads the two files into main memory. If the TEST initialization program has been omitted, the message

UNDEFINED EXTS

is printed indicating TEST is an undefined external to the MDE program.

To proceed, enter

:GO,1

When loading is finished, the message

LOADER COMPLETE

is printed.

Table 5-3. Alphabetical List of MDE Error Messages

Message	Meaning/Recovery
CAN'T FILL MORE THAN 16 MODULES!	User has tried to write microprograms to more than the maximum of 16 WCS modules. The user can debug no more than 16 WCS modules at a time.
ILLEGAL COMMAND	Command just entered is not an MDE command; re-enter command.
ILLEGAL DIGIT	An "8" or "9" was entered in the previous command that called for an octal digit; re-issue the entire command.
ILLEGAL PARAMETER	An unacceptable parameter was entered in the previous command; re-issue command.
ILLEGAL REG. MNEMONIC	Register or flag mnemonic just entered is not one of those listed under the CHANGE command (section 5-32); enter correct mnemonic.
ILLEGAL TAPE #	Bit range entered is not one of those listed under PREPARE command (section 5-26).
MISSING PARAMETER	A required parameter was omitted from the previous command; re-issue command.
NO BREAKPOINT HAS BEEN SET!	An EXECUTE-from-breakpoint command was given without having set a breakpoint logically beyond the execute address.
WCS NOT LOADED	The Writable Control Store PCA corresponding to the logical unit specified in the command just entered, has not been loaded with a microprogram during this MDE session; load the WCS.

- d. Save the loaded MDE program with

:ST,P

To summon MDE from now on, enter

:PR,MDE

- e. The program title is then printed followed by command request:

MICRO-DEBUG EDITOR
COMMAND?

Now enter the MDE commands required as described beginning in Section 5-16.

5-38. BCS OPERATION OF MDE

Before proceeding, the following items must be available:

- An absolute BCS binary tape.
- A relocatable object tape of MDE (HP 12978-16004).
- A relocatable object tape of the TEST initialization program, if a debug run is to be made.
- A microprogram object tape.
- A BCS Library tape (HP 24145-60001), Revision B.

The following is an example of how the user can proceed. For details on additional BCS options, see the **Basic Control System** manual (HP 02116-9017).

- Load the absolute BCS binary tape using the Basic Binary Loader.
- Set the P-register to 2. Set bit 14 of the Switch Register and clear all other Switch Register bits.
- Place MDE relocatable object tape in the paper tape reader and insure that the paper tape reader and the console device are on. Turn on paper tape punch. Press PRESET and RUN on the CPU Front Panel.

The MDE tape is read and an absolute binary tape is punched.

- d. The message

*LOAD

is printed on the console device and the program halts.

If required, load the relocatable TEST Initialization Program tape into the paper tape reader. Press RUN.

The TEST tape is read and another absolute binary tape is punched.

- e. The message

*LOAD

is printed on the teleprinter and the program halts.

Set Switch Register bits 2 and 14 leaving all others clear. Load BCS Library tape into the paper tape reader. Press RUN.

- f. Library tape is read and more absolute binary tape is punched.

Linkage information is printed on the Teleprinter. Remove paper tape from punch. This is the complete absolute binary tape of the Micro Debug Editor including the TEST Initialization Program.

- g. Load this tape using the Basic Binary Loader.

- h. When loading is complete, set P-register to 2. Press PRESET and RUN. The message

MICRO-DEBUG EDITOR
COMMAND?

is printed.

- i. Now enter the required MDE commands as described beginning in Section 5-16.

5-39. WCS I/O UTILITY SUBROUTINE

This library subroutine provides the capability of writing a microprogram into and reading a microprogram from a WCS using a buffer in an Assembly Language, FORTRAN, or ALGOL program and operating in a BCS or DOS-III environment. This avoids the necessity of running MDE every time it is necessary to access a WCS. This subroutine is in the standard BCS and DOS-III libraries for 21MX Series Computers.

Unlike a ROM chip, whenever the computer power is turned off, the WCS contents are lost. Thus the WCS must be loaded before access can be made to microprograms. This WCS I/O utility has been provided to serve that purpose.

Besides the calling sequence, a buffer is required in the calling program large enough to hold the number of microinstructions being transferred in or out.

Initially, the microprogram is stored on an object paper tape, in an object file on disc, or as octal data stored in the Main Memory program. In the case where the microprogram is in the form of octal data in the Main Memory program, the octal data area serves as the buffer when the WCS I/O Utility is used to write the microprogram into the WCS.

In the case where the microprogram resides on disc or paper tape, the control system (BCS or DOS-III) must be used to read the tape or disc file into a buffer in the Main Memory program. It must be remembered that the microprogram object contains header and end record information that must be deleted before storing the microprogram in the buffer. (Header and end record information must not be written into the WCS.)

Refer to Section 5-7 for a description of the Binary object output by the micro-assembler. Appendix A illustrates the binary object format.

When the microprogram has been stored in the Main Memory program buffer, a WCS I/O Utility WRITE calling sequence is used to write the microprogram into the WCS.

To read the contents of the WCS, a WCS I/O Utility READ calling sequence is used.

The assembly language calling sequences are the following:

READ

JSB WREAD	Branch to WCS read subroutine
DEF *+5	Return address
DEF lun	Logical unit number of WCS
DEF BUFF	Address of microprogram buffer
DEF LENGTH	Number of words of transfer
DEF ADRS	WCS relative address

WRITE

JSB WWRITE	Branch to the WCS write subroutine
DEF *+4	Return address
DEF lun	Logical unit number of WCS
DEF BUFF	Address of microprogram buffer
DEF LENGTH	Number of words of transfer

Where lun contains the logical unit number of the WCS being accessed and BUFF contains the first word of a word pair that holds a micro-instruction. LENGTH contains the octal number of words in the transfer; if LENGTH is positive, the number of 24 bit words is specified; if LENGTH is negative, the number of 16 bit words is specified. ADRS contains the WCS relative address (between octal addresses 0 and 377) of where to start reading.

WRITABLE CONTROL STORE

SECTION

VI

This section covers general information, installation, programming, and general theory of operation for the HP 12978A Writable Control Store Interface Kit. Options 001 and 002 for the interface kit are also covered in this section.

6-1. GENERAL INFORMATION

The Hewlett-Packard 12978A Writable Control Store Interface Kit provides the HP 21MX Computers with the necessary logic to dynamically change the instruction set of the computer. The printed-circuit assembly and flat cable assembly contained in the interface kits are shown in figure 6-1 and listed in table 6-1.

6-2. IDENTIFICATION

Hewlett-Packard uses five digits and a letter (12978A) for standard kit designations. If the designation of your kit does not agree with this number, there are differences between your kit and the kit described in this manual.

6-3. INTERFACE KIT CONTENTS

Table 6-1. Interface Kit Contents

INTERFACE KIT	CONTENTS	HP PART NO.
12978A	Writable Control Store PCA	12908-60006
	Flat Cable Assembly 5 Connectors	5060-8393
	Microprogramming 21MX Computers	02108-90008
	Diagnostic Paper Tape	24359-16001
	Diagnostic Manual	12908-90009

6-4. CONTENTS OF INTERFACE KIT OPTIONS

There are two 12978A Interface Kit Options. They contain material in addition to that contained in the basic interface kit. Option 001 provides all the software required for use of

the writable control store in the DOS-III system. Option 002 provides all the software required for the use in the BCS system.

Table 6-2. Additional Material for Interface Options

OPTION	ADDITIONAL MATERIAL	HP PART NO.
12978A-001	DOS-III WCS Driver	24278-60001
	DOS-III WCS I/O Utility	24333-60001
	DOS-III Micro-assembler	12978-16001
	Dos-III Micro Debug Editor	12978-16002
12978A-002	BCS WCS Driver	24277-60001
	BCS WCS I/O Utility	24283-60001
	BCS Micro-assembler	12978-16003
	BCS Micro Debug Editor	12978-16004

6-5. SPECIFICATIONS

Table 6-3 lists the characteristics and specifications of the writable control store PCA.

6-6. INSTALLATION

6-7. UNPACKING AND INSPECTION

If the shipping carton is damaged upon receipt, request that the carrier's agent be present when the kit is unpacked. Inspect the kit for damage (cracked, broken parts, etc.). If the kit is damaged and fails to meet specifications, notify the carrier and the nearest HP Sales and Service Office immediately. (Sales and Service Offices are listed at the back of this manual.) Retain the shipping container and the packing material for the carrier's inspection. The HP Sales and Service Office will arrange for the repair or replacement of the damaged item without waiting for any claims against the carrier to be settled.

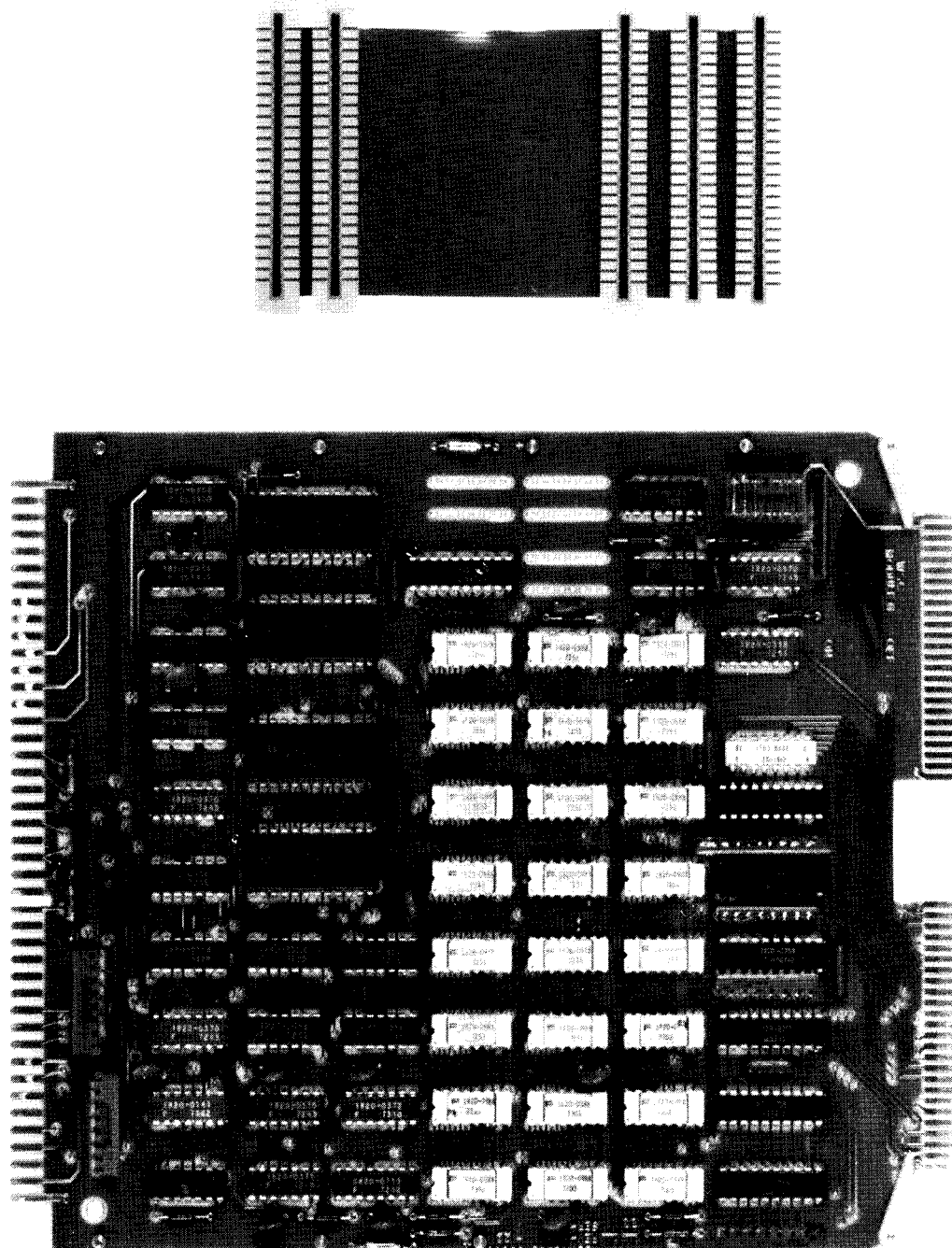


Figure 6-1. Writable Control Store Interface Kit

Table 6-3. Writable Control Store PCA Specifications

CAPACITY

Words Available: 256 per module

Maximum WCS Modules: one per HP 2105; two per HP 2108

Word Size: 24 bits

MICRO-INSTRUCTION TIME

Access: 162 ns.

Full Micro-instruction Cycle: 325 ns.

INSTALLATION

One writable control store PCA requires the use of one Input/Output slot (slot 10). Writable control store may be used as any module, except module 0.

DATA STORAGE

Input/Output Group instructions or an HP 21MX Dual Channel Port Controller are used to load the WCS.

DATA READBACK

Input/Output Group instructions only are used to read data from the WCS.

INTERFACE CURRENT SUPPLIED BY COMPUTER

0.15A (-2V supply); 4.6A (+5V supply)

PCA DIMENSIONS

Width: 7-3/4 inches (196.8 mm)

Height: 8-11/16 inches (220.7 mm)

PCA WEIGHT

Net Weight: 18 oz (511.2 gm) (card and cable only)

Shipping Weight: 4 lb (2.27 kg)

PCA INPUT LEVELS

"1" state: 1.9 volts minimum

"0" state: 1.1 volts maximum

PCA OUTPUT LEVELS

"1" state: 2.4 volts minimum

"0" state: 0.7 volts maximum

6-8. INSTALLATION

Install the writable control store kit as follows:

- a. Ensure that the computer operates properly prior to installing the writable control store interface kit.
- b. Turn off power at the computer.
- c. Remove the bottom and back access covers from the computer.
- d. On the writable control store remove the appropriate jumper wires from TB1 to select the desired module number (see figure 6-2 for pin number configuration). Refer to table 6-4 for the desired module number and jumper removal.
- e. On the writable control store PCA place the WCS module 0 enable switch S1 in the OFF position.
- f. Place the first writable control store PCA in slot number 10 (select code 10) of the I/O section of the computer. Any additional writable control store PCAs should be placed in slot 11.

Note: When WCS PCAs are installed, computer software must be reconfigured because of the changed I/O slot usage. If adding WCS PCA(s) will overburden the Power Supply of the computer, it may be necessary to move some I/O PCAs to an I/O Extender, HP 2156A.

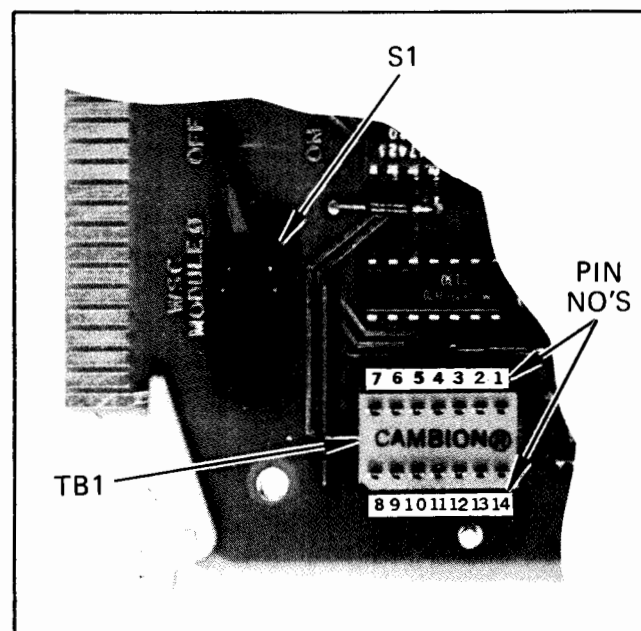


Figure 6-2. WCS Terminal Board for Selecting Module Number Position

Table 6-4. WCS PCA Jumper Removal on Terminal Board for Various Module Selections

MODULE	JUMPERS TO BE REMOVED
0	None
1	Pins 6,9
2	Pins 5,10
3	Pins 6,9; 5,10
4	Pins 4,11
5	Pins 6,9; 4,11
6	Pins 5,10; 4,11
7	Pins 6,9; 5,10; 4,11
8	Pins 3,12
9	Pins 6,9; 3,12
10	Pins 5,10; 3,12
11	Pins 6,9; 5,10; 3,12
12	Pins 4,11; 3,12
13	Pins 6,9; 4,11; 3,12
14	Pins 5,10; 4,11; 3,12
15	Pins 6,9; 5,10; 4,11; 3,12

g. Remove the ROM-CPU Interconnect assembly, part no. 5060-8344. Install the connectors of the flat cable assembly, part no. 5060-8393

1. on J1 of the ROM Control PCA 1, A7
2. on J2 of the CPU A1
3. on J1 of each WCS PCA

as shown in sideview on figure 6-3.

Note: If an I/O PCA is installed immediately above the WCS (refer to figure 6-3) that requires a cable (hood) connector on the back, then it may be necessary to double the flat cable assembly back or cut it to make room for the I/O cable connector.

h. Replace the bottom and back access covers on the computer.

i. Turn on power at the computer and perform the diagnostic test as outlined in the Diagnostic Program Procedures (part no. 12908-90009) shipped with the 12978A Interface Kit. If the diagnostic program is completed without error, the PCA is installed and operating properly. If the diagnostic program indicates errors, halt the computer, turn off power, and recheck all of the above installation procedures. Correct where necessary, then recheck and repeat the operating procedures of the diagnostic.

6-9. RESHIPMENT

If an item of the kit is to be shipped to Hewlett-Packard for service or repair, attach a tag to the item identifying the owner and indicating the service or repair to be accomplished. Include the model number of the kit. Package the item in the original factory packaging material, if available. If the original material is not available, standard factory packaging material can be obtained from a local Hewlett-Packard Sales and Service Office. If standard factory packaging material is not used, wrap the item in Air Cap TH-240 Cushioning (or equivalent) manufactured by Sealed Air Corp., Hawthorne, N.J. and place in a corrugated carton (200 pound test material). Seal the shipping carton securely and mark it "FRAGILE" to ensure careful handling.

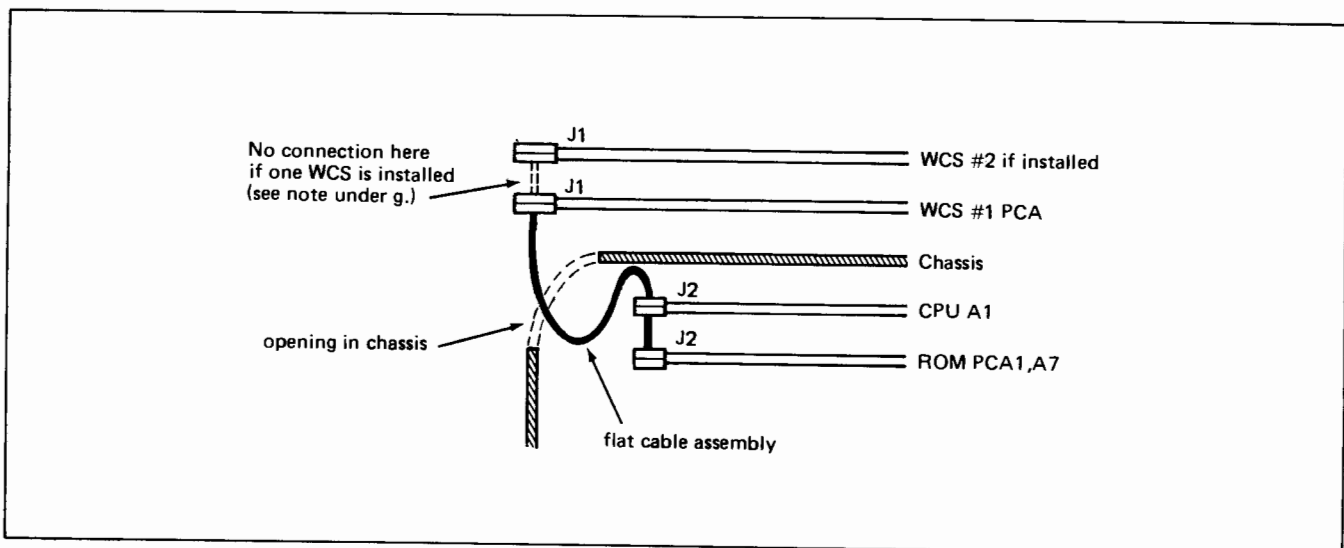


Figure 6-3. Installation of Flat Cable Assembly

Note: In any correspondence identify the kit by model number. Refer any questions to the nearest Hewlett-Packard Sales and Service Office.

6-10. PROGRAMMING

Two methods exist for writing data into (loading) a WCS module: under program control and under control of the Dual Channel Port Controller (DCPC). Under program control, prior to initiating the load routine, the data to be loaded must be stored in the computer memory. This requires a block of up to 512 words per module. The load routine will send two words from memory (32 bits which are mapped into an 8 bit address and a 24 bit micro-instruction) to the WCS module, issue a write command to that module and cause the data to be stored there. The load routine will repeat this process until the desired number of words have been stored in the WCS module.

Once loaded, the contents of the WCS module may be read back under program control via the I/O bus and compared with their counterpart in memory.

Timing sequences for flags used in the following examples are shown in figure 6-4.

6-11. PROGRAM EXAMPLE: LOADING WCS

The following is an example of the program sequence necessary for loading a WCS under program control. This example does not include block pointers, counters, etc., which are necessary for proper control.

Note: "SC" indicates select code of the WCS PCA.

STF	SC	Initializes the Direction FF (flip-flop or flag)
OTA	SC	Loads the first computer word into first WCS buffer and toggles the Direction FF. This word comprises the 8 bit address and the 8 most significant bits of the micro-instruction.
OTB	SC	Loads the second computer word into the second WCS buffer and toggles the Direction FF. This word comprises the 16 least significant bits of the micro-instruction.
STC	SC	Provides the write pulse to load the WCS buffers into the RAM.

The OTA, OTB, and STC instructions are normally in a loop that is repeated until the desired number of micro-

instructions have been stored. OTA/OTB was chosen as an example; any combination of these instructions is allowable.

6-12. PROGRAM EXAMPLE: READING WCS

An example of reading from WCS under program control via the I/O bus is shown below. This example is shown without regard to the block pointers, counters, etc., which are necessary for proper control.

STF	SC	Initializes the Direction FF.
OTA	SC	Sends the 8 bit address to the WCS module from the 8 most significant bits of the A-register. (B-register could be used, as well).
STF	SC	Re-initializes the Direction FF.
LIA	SC	Places eight zeros into the 8 most significant bit positions of the A-register and places the eight most significant bits of the micro-instruction into the eight least significant bit positions of the A-register.
LIB	SC	Places the 16 least significant bits of the micro-instruction into the B-register.

The STF, OTA, STF, LIA, and LIB sequence is normally in a loop that is repeated until the desired number of micro-instructions have been read in from WCS. LIA/LIB was chosen as an example; any combination of these instructions is allowable.

6-13. PROGRAM EXAMPLE: LOADING WCS BY DUAL CHANNEL PORT CONTROLLER

Under Dual Channel Port Controller (DCPC) control, the load routine must send only the three DCPC control words to the selected channel. When the channel is turned on, DCPC will utilize every I/O cycle until the entire block of data is sent to the WCS module (maximum of 512 cycles). DCPC will transfer these words at a rate of $1.62 \mu\text{s}/\text{word}$ (512 words will take $830 \mu\text{s}$ to transfer).

The following is an example of the program sequence necessary for loading WCS via DCPC. This example does not include block pointers, counters, etc., which are necessary for proper control.

Writable Control Store

LDA	CW1	Get the first DCPC control word.
OTA	6	Send the first DCPC control word to the selected DCPC channel (DCPC channel 1 has been selected here for demonstration purposes only).
CLC	2	Prepare the selected DCPC channel to receive the second DCPC control word.
LDA	CW2	Get the second DCPC control word.
OTA	2	Send the second DCPC control word to the selected DCPC channel.
STC	2	Prepare the selected DCPC channel to receive the third DCPC word.
LDA	CW3	Get the third DCPC control word.
OTA	2	Send the third DCPC control word to the selected DCPC channel.
STC	6,C	Turn on the selected DCPC channel.
STF	SC	Initialize the Direction FF.
CLF	SC	Start DCPC transfer.
SFS	6	Test for the completion of the transfer.
JMP	*-1	Loop until done.
CW1	OCT	12000SC
CW2	OCT	(Starting address of the block to be transferred)
CW3	OCT	(Two's complement of the number of computer words to be transferred.)

6-14. GENERAL THEORY OF OPERATION

Writable Control Store (WCS) consists of a bipolar semiconductor Random Access Memory (RAM) containing 24 integrated circuit (IC) packages mounted on a 2100-size printed-circuit assembly (PCA). Also included is the flat jumper cable assembly necessary for complete mechanization within the HP 21MX Computer. The WCS PCA should be installed only in slots 10 (standard) and 11 of the computer I/O slots. Each IC package is configured in 256 bits and organized as one bit per word. Thus one module of WCS is capable of storing 256 words of 24 bits each.

For the purpose of execution of WCS instructions, WCS can be configured to be addressed as any one of the computer's ROM modules except module 0. One WCS module can be installed on an HP 2105 Computer. Two WCS modules can be installed on an HP 2108 Computer.

6-15. WCS MODULE IDENTIFICATION

For proper addressing of WCS, an integrated circuit comparator and terminal board (with jumpers) on the WCS PCA is used to identify the PCA as a particular module of Control Store. For example, if the terminal board is configured for module 2, the PCA will be enabled when the ROM Address Register (RAR) contains the pattern "0010" in its four most significant bits (11-8), and disabled otherwise. When enabled, the word in WCS addressed by RAR bits 0 through 7 will be sent to the ROM Instruction Register (RIR) as signals ROM0 through ROM23. The computer will then execute this word (micro-instruction) as though it came from a standard ROM PCA. The access time of data from WCS (162 ns.) allows the computer to operate at its normal clock rate. If it is desired to replace any module already existing in ROM with a WCS module, that ROM module must be removed in order to prevent unwanted "or" conditions on the data lines.

Note: The ON position of switch S1 (figure 6-2) is not intended for use in the 21MX computers. All Control Store is disabled, if S1 is set to ON.

6-16. WCS CONNECTION

WCS is connected to the computer central processor through the I/O structure (for loading and checking), and also through a 50 wire flat cable connector. It is this connector that enables WCS to be used as an extension of the computer's basic control store. The cable connects one or two WCS PCAs to ROM control PCA 1,A7 and to CPU A1. The ROM address register on the CPU sends a 12 bit address to the WCS PCA or PCAs through this cable, and the addressed PCA then sends its data (micro-instruction) from that address back through this cable, where it is merged with the outputs of ROM. From there the data is sent to the ROM instruction register as though it was from ROM.

6-17. WCS ADDRESSING

To load the WCS RAM circuits, the WCS PCA must be addressed through the I/O interface structure of the computer. A 32 bit format is necessary and requires that a 2 word transfer be used in the loading procedure through the computer A- and/or B-registers. Two computer words and thus two transfer operations are required for one WCS word. The eight most significant bits of the first computer word transferred is the WCS RAM circuit address. The remaining eight bits of the first computer word and all 16 bits of the second computer word (total of 24 bits) are stored in WCS at the address specified.

Once loaded, WCS becomes an extension of the ROM. Thus the WCS may be used to alter the computer instruc-

tion set while the computer is in an operating condition. This feature permits dynamic expansion of the computer instruction set.

6-18. WCS Loading Timing diagram

Figure 6-4 illustrates the WCS timing.

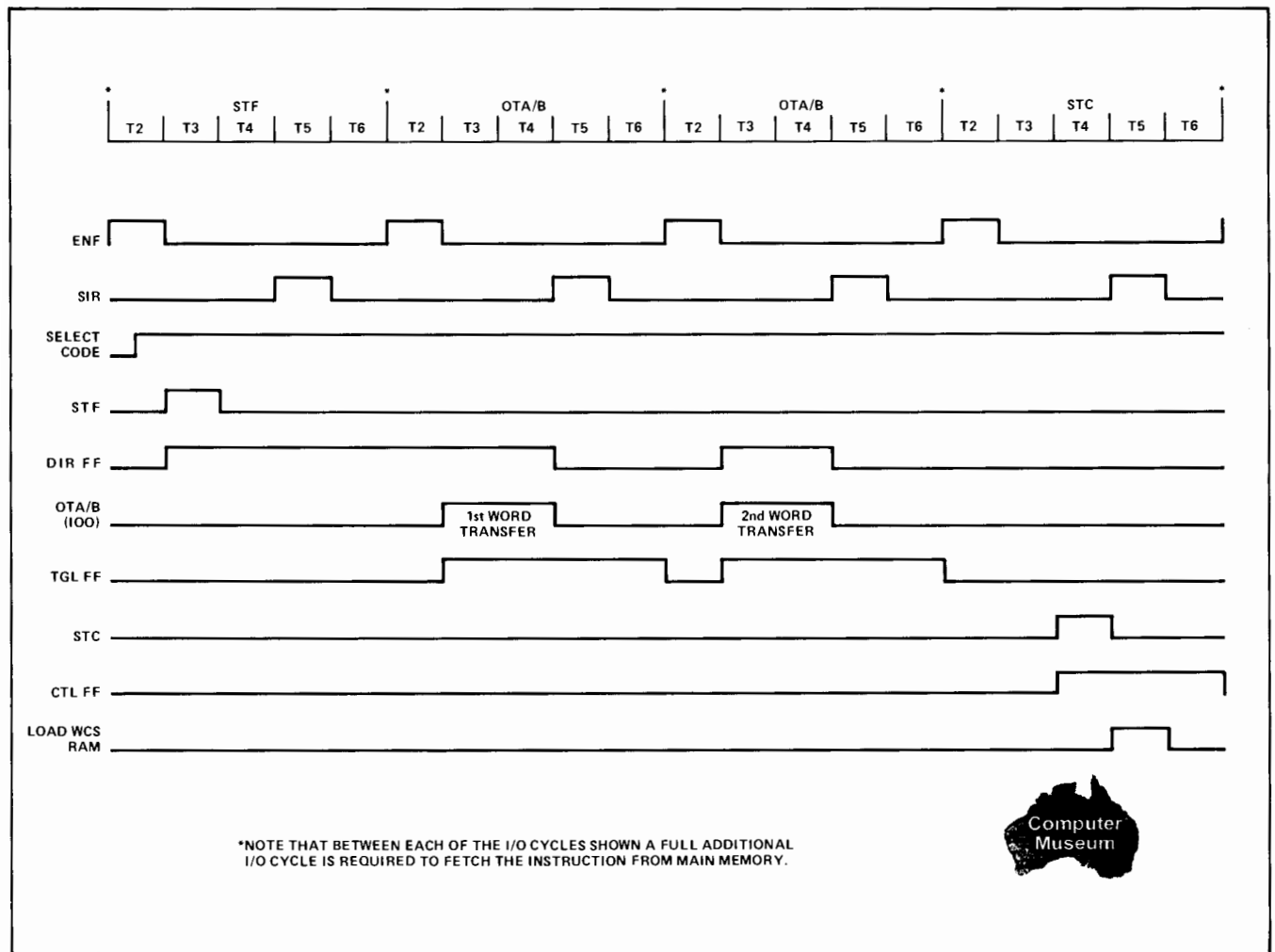


Figure 6-4. WCS Loading Timing Diagram

OBJECT TAPE FORMATS

APPENDIX

A

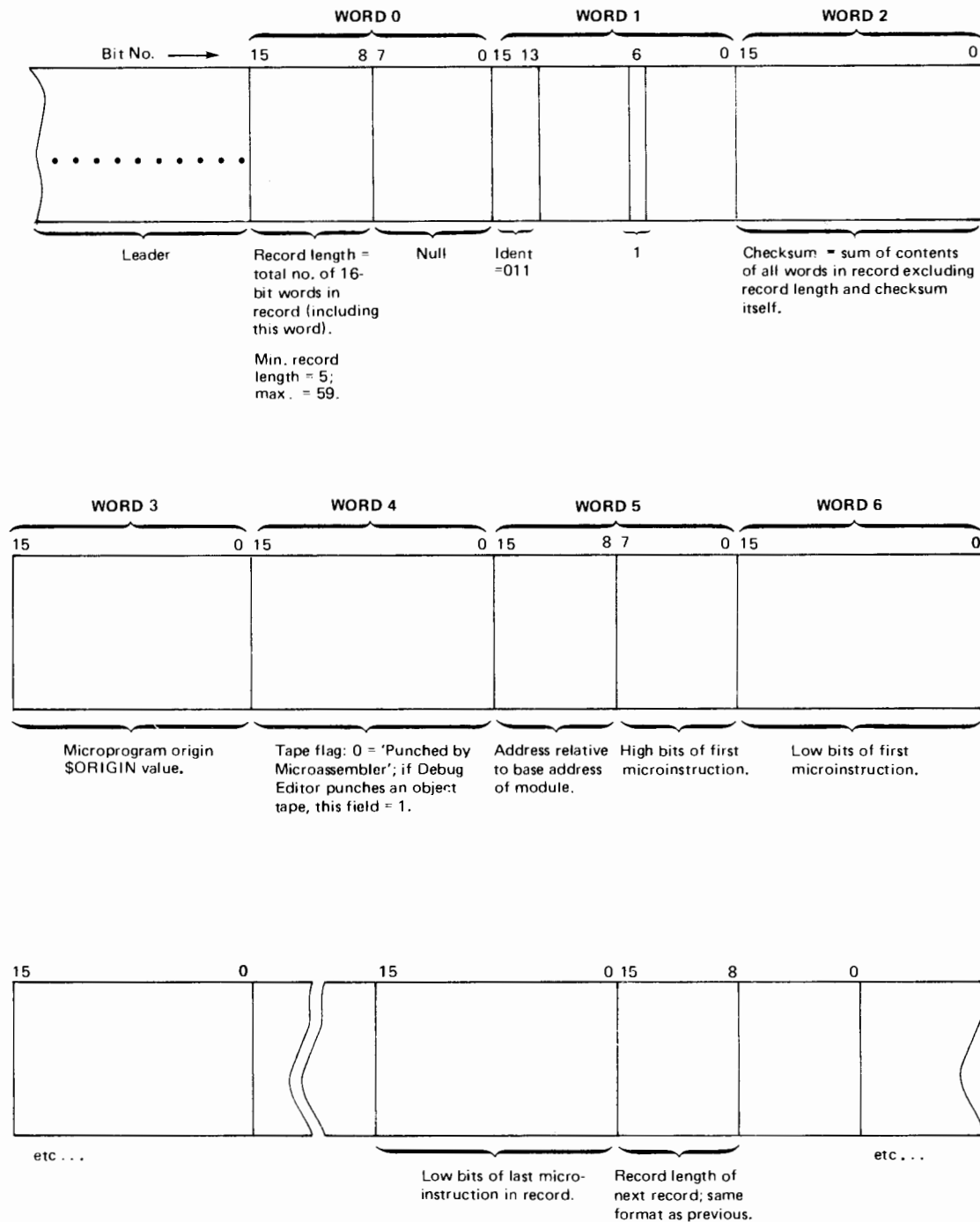


Figure A-1. Format of Standard Object Tape (Sheet 1 of 2)

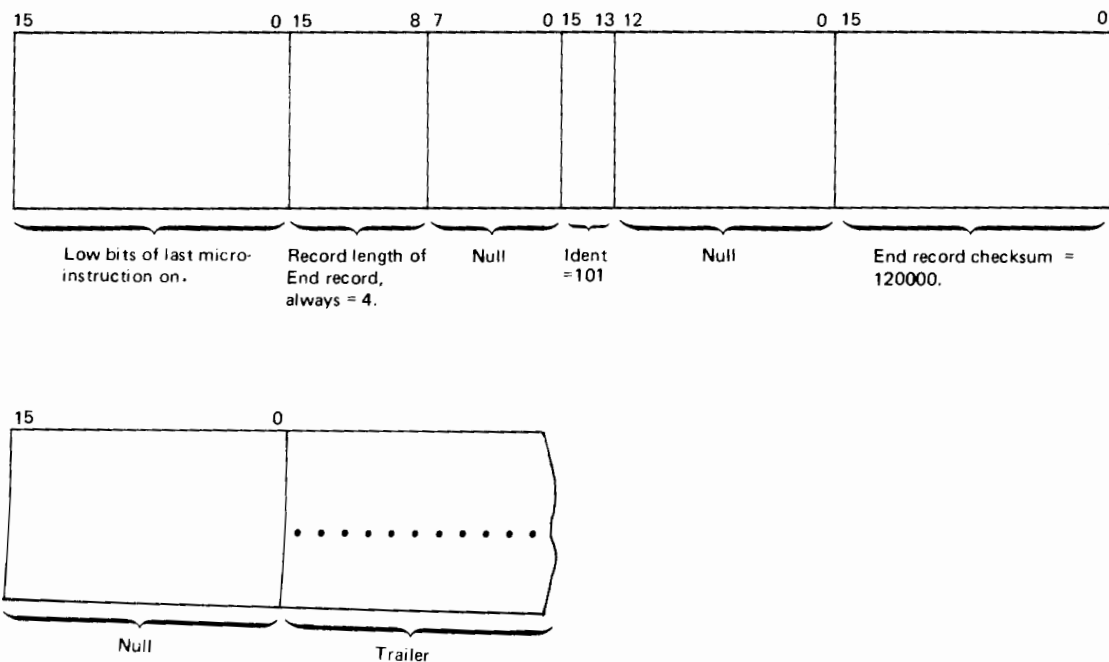


Figure A-1. Format of Standard Object Tape (Sheet 2 of 2)

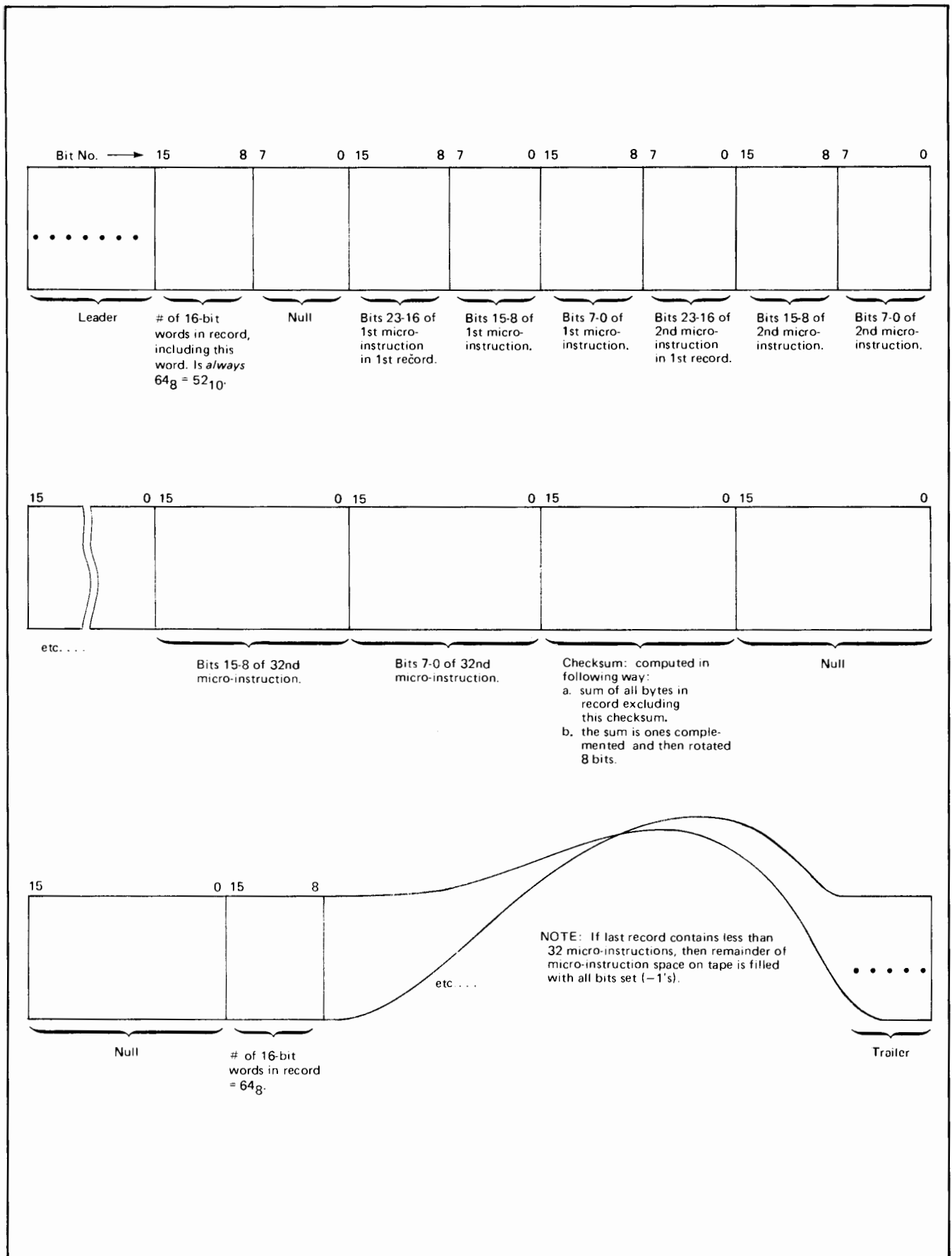


Figure A-2. Format of the \$RCASE Object Tape

MICRO-ORDER SUMMARY

APPENDIX

C

Table C-1. Summary of User Micro-orders

MICRO-ASSEMBLER → CARD COLUMN → BITS (ROM) →	OP 10 20-23	SPECIAL 15 0-4	ALU 20 15-19	JMP COND 20 15-19	IMMEDIATE MODIFIER 20 18-19	STORE 25 10-14	RJS 25 14	S-BUS 30 5-9
Corresponding Bit Pattern								
00000	*NOP	IOFF	INC	TBZ	HIGH	TAB	RJS	TAB
00001	ARS	SRG2	OP1	ONES	LOW	CAB		CAB
00010	CRS	L1	OP2	COUT	CMHI	T		T
00011	LGS	L4	ZERO	AL0	CML0	L		CIR
00100	MPY	R1	OP3	AL15		IOO		IOI
00101	DIV	ION	OP4	NMLS		CNTR		CNTR
00110	LWF	SRG1	SUB	CNT8		DSPL		DSPL
00111	WRTE	RES2	OP5	FPSP		DSPI		DSPI
01000	ASG	STFL	OP6	FLAG		IR		ADR
01001	READ	CLFL	ADD	E		M		M
01010	ENV	FTCH	OP7	OVFL		B		B
01011	ENVE	SOV	OP8	RUN		A		A
01100	JSB	COV	OP9	NHOI		MEU		LDR
01101	JMP	RPT	OP10	SKPF		CM		RES2
01110	IMM	SRGE	OP11	ASGN		PNM		MEU
01111	(not used)	*NOP	DEC	IR2		*NOP		*NOP
10000		MESP	CMPS	NLDR		S1		S1
10001		MPCK	NOR	NSNG		S2		S2
10010		IOG	NSAL	NINC		S3		S3
10011		ICNT	OP12	NDEC		S4		S4
10100		SHLT	NAND	NRT		S5		S5
10101		INCI	CMPL	NLT		S6		S6
10110		RES1	XOR	NSTR		S7		S7
10111		SRUN	SANL	NRST		S8		S8
11000		**UNCD	NSOL	NSTB		S9		S9
11001		CNDX	XNOR	NSFP		S10		S10
11010		JIO	PASL	INT		S11		S11
11011		JTAB	AND	SRGL		S12		S12
11100		J74	ONE	RUNE		X		X
11101		J30	SONL	*NOP		Y		Y
11110		RTN	IOR	CNT4		P		P
11111	(not used)	JEAU	*PASS	NMEU		S		S

*default micro-order

**JMP default

(10% screen) not normally used by user microprogrammer.

(20% screen) included here for completeness only; reserved for exclusive use of system microprogrammers.

FUNCTIONAL BLOCK DIAGRAM

APPENDIX

D

ARITHMETIC AND LOGIC SECTION

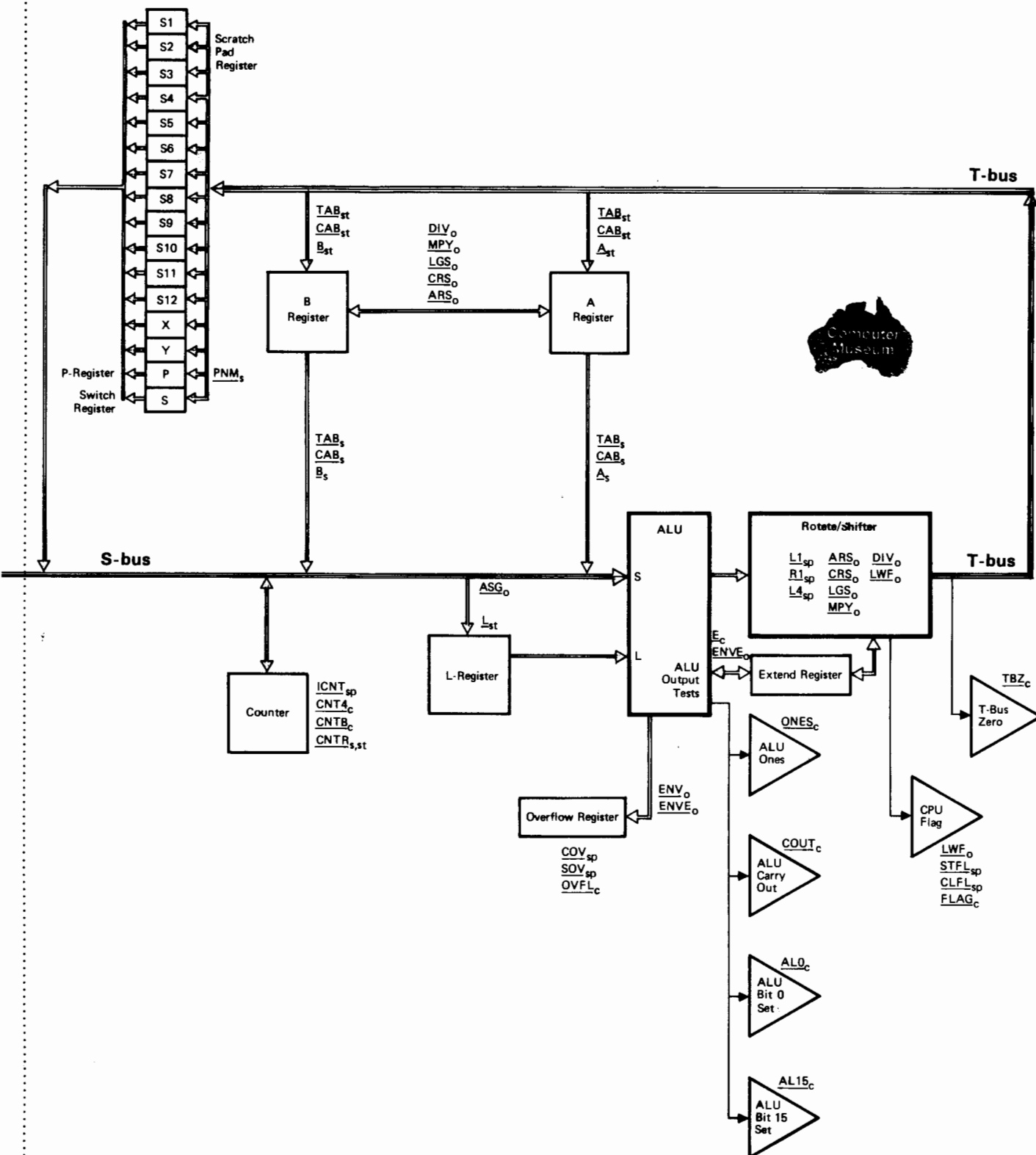
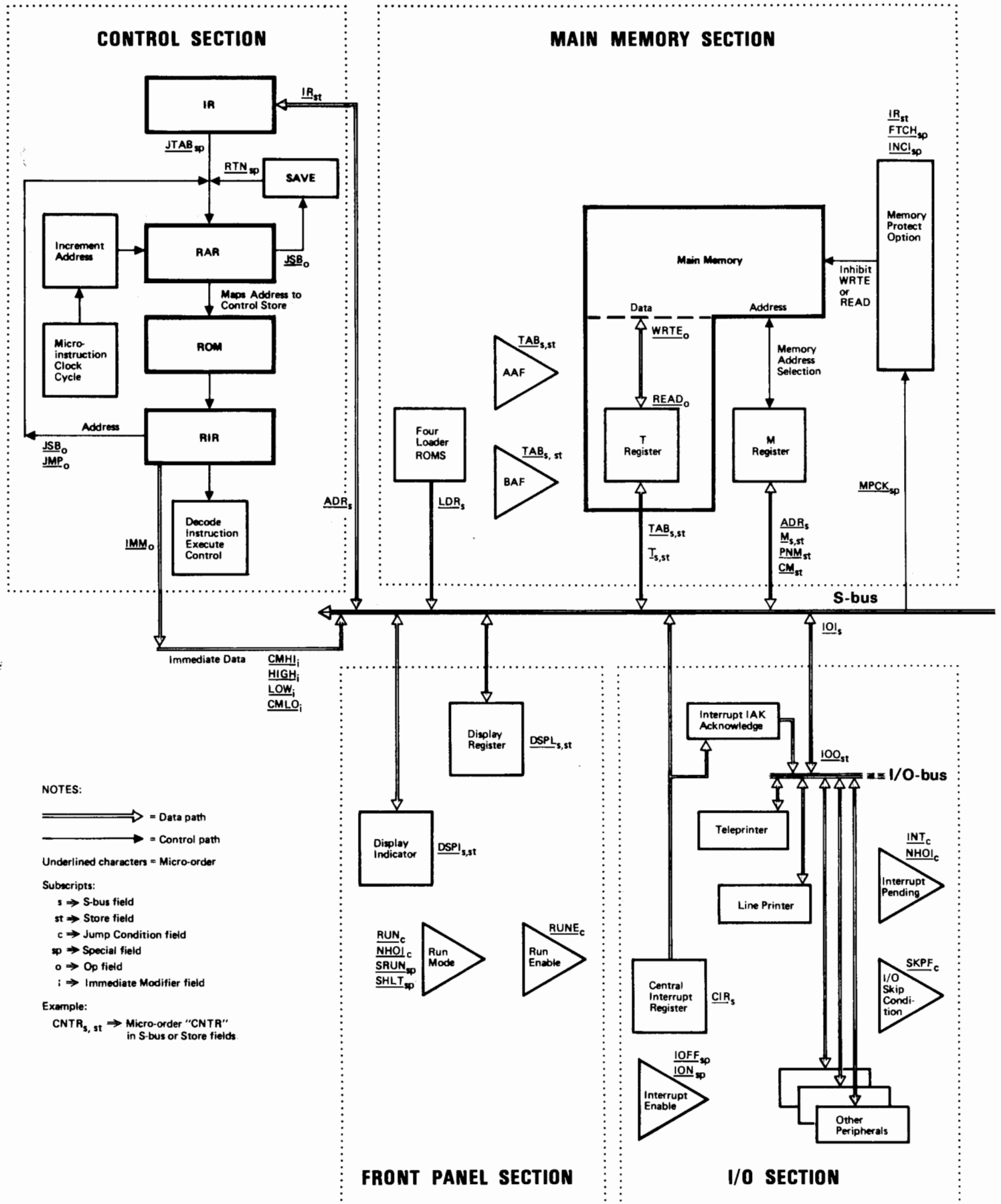


Figure D-1. Functional Block Diagram

CONTROL SECTION

MAIN MEMORY SECTION



BASIC INSTRUCTION SET MICROPROGRAM LISTING

APPENDIX

E

This appendix holds a full micro-assembly listing of the 21MX Computer Basic Instruction Set microprogram. Due to the size of this microprogram, a special micro-assembler was used. Minor differences can be seen between this micro-assembly listing and a listing produced by the micro-assembler described in this manual. The major difference to be noted is that octal numbers are preceded by a "%" symbol in this listing. Other differences are self explanatory.

```

0001          $ORIGIN=0
0002          *****
0003          *
0004          *      21MX MICRO-CODE
0005          *      MODULE 0
0006          *
0007          *****
0008          HALT      EQU          %0400
0009          FADD      EQU          %7125
0010          FSUB      EQU          %7126
0011          FMPY      EQU          %7221
0012          FDIV      EQU          %7262
0013          IFIX      EQU          %7000
0014          FLOAT      EQU          %7025
0015          *****
0016          *      FETCH ROUTINE
0017          *****
0018 0000 220 074712  FETCH      READ FTCH INC PNM P      M<=P; PC=PC+1; READ NEW INSTR
0019 0001 017 136745          ISN          ENABLE INTERRUPT RECOGNITION
0020 0002 017 100411          CLFL PASS IR TAB      IR<= T/A/B; CLR FLAG FF
0021 0003 220 020673          READ JTAB INC CM ADR      JUMP THRU TABLE; LOAD M IF MRG INSTR
0022          *****
0023 0004 325 120031  HORI      JMP CNDX RUN RJS HALT      RUN MODE IMPLIES AN INTERRUPT
0024          *****
0025          *      INTERRUPT RESPONSE ROUTINE
0026          *****
0027 0005 237 106451  INTERRUPT READ CLFL PASS M CIR      M<=CIR; READ TRAP CELL; CLR FLAG FF
0028 0006 320 000471          JMP CNDX TBZ RJS INTOK      CHECK IF CIR IS VALID
0029 0007 237 106457          READ      PASS M CIR      M<=CIR; READ TRAP CELL
0030 0010 320 040031          JMP CNDX TBZ      FETCH      IF NO INT BY NOW, IGNORE
0031 0011 017 104400  INTOK      IOFF PASS IR T      IR<= TRAP CELL; DISABLE INT RECOG.
0032 0012 220 020673          READ JTAB INC CM ADR      JUMP THRU TABLE; LOAD M IF MRG INSTR.
0033          *****
0034          *      INDIRECT ROUTINE
0035          *****
0036 0013 220 022457  INDLEVEL READ      INC M M      READ NEXT LEVEL
0037 0014 326 001031          JMP CNDX WH01 RJS IND2      HALT OR INTERRUPT?
0038 0015 017 100465  INDIRECT      INCI PASS M TAB      M<=T/A/B; INCR INDIRECT COUNT
0039 0016 322 040571          JMP CNDX ALIS      INDLEVEL      CHECK FOR ANOTHER LEVEL OF INDIRECT
0040 0017 220 022476          READ RTH INC M M      READ EFFECTIVE ADDRESS, RETURN
0041 0020 017 100465  IND2      INCI PASS M TAB      M<=T/A/B; INCR INDIRECT COUNT
0042 0021 330 100731          JMP CNDX WONG RJS INDIRECT+1JUMP BACK FOR SINGLE INSTRUCTION
0043 0022 007 175717          DEC P P      RESET P
0044 0023 320 000230          JMP          HORI      HALT OR INTERRUPT

```



```

0045 *****
0046 *      ALTER-SKIP GROUP
0047 *****
0048 0024 017 102757 ASGNOP          PASS      CAB      SET UP SKIP TEST
0049 0025 327 042431          JMP  CNDX ASSN      ASGNSKP  JUMP IF ASG SKIP NOT MET
0050 0026 200 075717          ASS      INC P      P      P<=P+1; ENABLE ASG HARDWARE
0051 0027 327 100031          JMP  CNDX IR2 RJS  FETCH  DONE IF NOT INA/B
0052 0030 260 002076          ENVE RTN INC CAB  CAB      A/B <= A/B PLUS 1
0053 *
0054 0031 001 136057 ASGCL*          ZERO CAB      CLEAR A/B REGISTER
0055 0032 327 042431          JMP  CNDX ASSN      ASGNSKP  JUMP IF ASG SKIP NOT MET
0056 0033 200 075717          ASS      INC P      P      P<=P+1; ENABLE ASG HARDWARE
0057 0034 327 100031          JMP  CNDX IR2 RJS  FETCH  DONE IF NOT INA/B
0058 0035 260 002076          ENVE RTN INC CAB  CAB      A/B <= A/B PLUS 1
0059 *
0060 0036 010 002057 ASGCH*          CNPS CAB  CAB      A/B <= NOT A/B
0061 0037 327 042431          JMP  CNDX ASSN      ASGNSKP  JUMP IF ASG SKIP NOT MET
0062 0040 200 075717          ASS      INC P      P      P<=P+1; ENABLE ASG HARDWARE
0063 0041 327 100031          JMP  CNDX IR2 RJS  FETCH  DONE IF NOT INA/B
0064 0042 260 002076          ENVE RTN INC CAB  CAB      A/B <= A/B PLUS 1
0065 *
0066 0043 016 036057 ASGCC*          ONE CAB      CLR & COMP A/B REGISTER
0067 0044 327 042431          JMP  CNDX ASSN      ASGNSKP  JUMP IF ASG SKIP NOT MET
0068 0045 200 075717          ASS      INC P      P      P<=P+1; ENABLE ASG HARDWARE
0069 0046 327 100031          JMP  CNDX IR2 RJS  FETCH  DONE IF NOT INA/B
0070 0047 260 002076          ENVE RTN INC CAB  CAB      A/B <= A/B PLUS 1
0071 *
0072 0050 217 136757 ASGNSKP ASS      NO SKIP; ENABLE ASG HARDWARE
0073 0051 327 100031          JMP  CNDX IR2 RJS  FETCH  DONE IF NOT INA/B
0074 0052 260 002076          ENVE RTN INC CAB  CAB      A/B <= A/B PLUS 1
0075 *****
0076 *      SHIFT/ROTATE GROUP
0077 *****
0078 0053 017 102046 SRG          SRG1 PASS CAB  CAB      FIRST SHIFT
0079 0054 017 102056          SRG2 PASS CAB  CAB      CHECK FOR CLEAR E; SET SLA TEST
0080 0055 335 103031          JMP  CNDX SRGL RJS  **3    SRGL IS SLA TEST
0081 0056 017 102041          SRG2 PASS CAB  CAB      SECOND SHIFT
0082 0057 000 075736          RTN INC P      P      P<=P+1, WHEN LSB = 0
0083 0060 017 102041          SRG2 PASS CAB  CAB      SECOND SHIFT
0084 0061 017 136776 RETURN          RTN

0085 *****
0086 *      I/O GROUP
0087 *****
0088 0062 017 136757 IOCNTRL NOP          ALLOW TIME TO GET SKIP FLAG
0089 0063 326 100031          JMP  CNDX SKPF RJS  FETCH  CHECK SKIP FLAG
0090 0064 000 075736          RTN INC P      P      P <= P + 1
0091 0065 017 136757          NOP
0092 *
0093 0066 017 102757 IO.OT*          PASS      CAB      SET UP S-BUS
0094 0067 017 102217          PASS IO0 CAB      I/O-BUS <= A/B
0095 0070 017 102236          RTN PASS IO0 CAB      HOLD I/O-BUS VALID
0096 0071 017 136757          NOP
0097 *
0098 0072 017 136757 IO.LI*          NOP          SYNCHRONIZE IOI PULSE
0099 0073 017 136757          NOP
0100 0074 017 110076          RTN PASS CAB  IOI      A/B <= I/O-BUS
0101 0075 017 136757          NOP
0102 *
0103 0076 017 136757 IO.MI*          NOP          SYNCHRONIZE IOI PULSE
0104 0077 017 102157          PASS L      CAB      L <= A/B FOR ALU OPERATION
0105 0100 017 010076          RTN IO0 CAB  IOI      A/B<= (A/B) + (I/O BUS)
0106 *
0107 *****
0108 *      IO GROUP/ EAU GROUP/ MAC GROUP JUMPS
0109 *****
0110 0101 320 003122 IOG          JMP  IOG          IOCNTRL
0111 0102 320 015437 EAU          JMP  JEQU          EAUTABLE
0112 0103 320 016034 MAC0          JMP  J74          MACTABLE0
0113 0104 320 017034 MAC1          JMP  J74          MACTABLE1

```

```

0114 *****
0115 *      MEMORY REFERENCE GROUP
0116 *****
0117 0105 300 000670 AND,I JSB      INDIRECT
0118 0106 017 126157 AND      PASS L   A      L <= A
0119 0107 015 100576 RTN AND A   TAB      A <= T/A/B AND L
0120 *
0121 0110 300 000670 CP*,I JSB      INDIRECT
0122 0111 017 102157 CP*      PASS L   CAB      L <= A/B
0123 0112 013 000757      XOR      TAB      T-BUS <= T/A/B XOR L
0124 0113 320 040031 JMP CNDX TBZ  FETCH  JUMP TO FETCH IF EQUAL
0125 0114 000 075736 RTN INC P   P      P <= P+1 IF NOT EQUAL
0126 *
0127 0115 300 000670 XOR,I JSB      INDIRECT
0128 0116 017 126157 XOR      PASS L   A      L <= A
0129 0117 013 000576 RTN XOR A   TAB      A <= T/A/B XOR L
0130 *
0131 0120 300 000670 IOR,I JSB      INDIRECT
0132 0121 017 126157 IOR      PASS L   A      A <= T/A/B IOR L
0133 0122 017 000576 RTN IOR A   TAB
0134 *
0135 0123 300 000670 ST*,I JSB      INDIRECT
0136 0124 017 122761 ST*      MPOCK PASS  M      MEM PROTECT CHECK OF ADDRESS
0137 0125 177 102036 WRTN RTN PASS TAB CAB  T/A/B <= A/B; WRITE
0138 *
0139 *
0140 0126 300 000670 AD*,I JSB      INDIRECT
0141 0127 017 102157 AD*      PASS L   CAB      L <= A/B
0142 0130 264 100076 ENVE RTN ADD CAB TAB  A/B <= T/A/B PLUS L
0143 *
0144 0131 300 000640 JSB,I JSB IOFF      INDIRECT  DISABLE INTERRUPT RECOGNITION
0145 0132 017 122761 JSB      MPOCK PASS  M      MEM PROTECT CHECKS THIS ADDR
0146 0133 177 174017 WRTN      PASS TAB  P      T/A/B <= RETURN ADDRESS; WRITE
0147 0134 000 023736 RTN INC P   M      P <= M + 1
0148 *
0149 0135 300 000670 ISZ,I JSB      INDIRECT
0150 0136 017 122761 ISZ      MPOCK PASS  M      MEM PROTECT CHECKS THIS ADDR
0151 0137 000 001017      INC S1  TAB      S1 <= T/A/B + 1
0152 0140 177 140017 WRTN      PASS TAB  S1      T <= S1; WRITE
0153 0141 320 000031 JMP CNDX TBZ  RJS  FETCH  ZERO? NO, DONE.
0154 0142 000 075736 RTN INC P   P      YES, P <= P+1
0155 *
0156 0143 300 000670 LD*,I JSB      INDIRECT
0157 0144 017 100076 LD*      RTN PASS CAB TAB  A/B <= T/A/B

```

```

0158 0145 017 136765 JMP,I      INCI      COUNT ONE INDIRECT LEVEL
0159 0146 017 101090      IOFF PASS S1  TAB  DISABLE INT RECOGNITION; S1<=T/A/B
0160 0147 322 046531 JMP CNDX AL15 JINDL  JMP IF ANOTHER LEVEL OF INDIRECT
0161 0150 017 140761      MPOCK PASS  S1      MEM PROT CHECKS DESTINATION ADDR
0162 0151 017 141736 RTN PASS P   S1      P <= DESTINATION ADDR
0163 *
0164 0152 220 040457 JINDL  READ  INC M   S1      READ NEXT LEVEL
0165 0153 326 007031 JMP CNDX MHD1 RJS  HORICK  JMP IF HALT OR INT
0166 0154 017 101025      INCI PASS S1  TAB      S1 <= T/A/B; COUNT INDIRECT LEVEL
0167 0155 322 046531 JMP CNDX AL15 JINDL  JMP IF ANOTHER LEVEL OF INDIRECT
0168 0156 017 140761      MPOCK PASS  S1      MEM PROT CHECKS DESTINATION ADDR
0169 0157 017 141736 RTN PASS P   S1      P <= DESTINATION ADDR
0170 *
0171 0160 017 101025 HORICK  INCI PASS S1  TAB      S1 <= T/A/B; COUNT INDIRECT LEVEL
0172 0161 330 106671 JMP CNDX HSHG RJS  JINDL+3  JUMP BACK FOR SINGLE INSTRUCTION
0173 0162 007 175717      DEC P   P      RESET P
0174 0163 320 000230 JMP      HORI      HALT OR INTERRUPT
0175 0164 017 121021      MPOCK PASS S1  ADR      S1<=DESTINATION ADDR; CHECK WITH M.P.
0176 0165 017 141736 RTN PASS P   S1      P <= DESTINATION ADDRESS

```

Appendix E

```

0177 *****
0178 * EAU MICROPROGRAMS
0179 *****
0180 0166 010 021017 RRR CMPS S1 ADR
0181 0167 000 041017 INC S1 S1 S1 <= TWO'S COMP OF SHIFTS
0182 0170 017 140255 RPT PASS CNTR S1 SET UP COUNTER FOR REPEAT
0183 0171 057 124504 CRS R1 PASS B B DOUBLE-WORD SHIFT REPEAT
0184 0172 017 136776 RTN
0185 *
0186 0173 010 021017 ASR CMPS S1 ADR
0187 0174 000 041014 COV INC S1 S1 S1 <= TWO'S COMP OF SHIFTS
0188 0175 017 140255 RPT PASS CNTR S1 SET UP COUNTER FOR REPEAT
0189 0176 037 124504 ARS R1 PASS B B DOUBLE-WORD SHIFT REPEAT
0190 0177 017 136776 RTN
0191 *
0192 0200 010 021017 LSR CMPS S1 ADR
0193 0201 000 041017 INC S1 S1 S1 <= TWO'S COMP OF SHIFTS
0194 0202 017 140255 RPT PASS CNTR S1 SET UP COUNTER FOR REPEAT
0195 0203 077 124504 LGS R1 PASS B B DOUBLE-WORD SHIFT REPEAT
0196 0204 017 136776 RTN
0197 *
0198 0205 010 021017 RRL CMPS S1 ADR
0199 0206 000 041017 INC S1 S1 S1 <= TWO'S COMP OF SHIFTS
0200 0207 017 140255 RPT PASS CNTR S1 SET UP COUNTER FOR REPEAT
0201 0210 057 124502 CRS L1 PASS B B DOUBLE-WORD SHIFT REPEAT
0202 0211 017 136776 RTN
0203 *
0204 0212 010 021017 ASL CMPS S1 ADR
0205 0213 000 041014 COV INC S1 S1 S1 <= TWO'S COMP OF SHIFTS
0206 0214 017 140255 RPT PASS CNTR S1 SET UP COUNTER FOR REPEAT
0207 0215 037 124502 ARS L1 PASS B B DOUBLE-WORD SHIFT REPEAT
0208 0216 017 136776 RTN
0209 *
0210 0217 010 021017 LSL CMPS S1 ADR
0211 0220 000 041017 INC S1 S1 S1 <= TWO'S COMP OF SHIFTS
0212 0221 017 140255 RPT PASS CNTR S1 SET UP COUNTER FOR REPEAT
0213 0222 077 124502 LGS L1 PASS B B DOUBLE-WORD SHIFT REPEAT
0214 0223 017 136776 RTN
0215 *
0216 0224 220 074457 DLD READ INC M P READ MEMORY ADDRESS
0217 0225 300 000640 JSB IOFF INDIRECT JSB TO GET M<=ADDR OF FIRST WORD
0218 0226 000 023017 INC S1 M S1 <= ADDRESS OF SECOND WORD
0219 0227 017 100557 PASS A TAB A <= FIRST DATA WORD
0220 0230 220 040457 READ INC M S1 M<=ADDR OF SECOND WORD; READ
0221 0231 000 075717 INC P P P <= P + 1
0222 0232 017 100536 RTN PASS B TAB B <= SECOND DATA WORD
0223 *
0224 0233 220 074457 DST READ INC M P READ MEMORY ADDRESS
0225 0234 300 000640 JSB IOFF INDIRECT JSB TO GET M <= ADDR OF FIRST WORD
0226 0235 000 023021 MPCK INC S1 M MP CHECK FIRST ADDR; S1<=SECOND ADDR
0227 0236 177 126017 WRT PASS TAB A STORE A INTO FIRST LOCATION
0228 0237 000 040461 MPCK INC M S1 MP CHECK S1; M<=S1
0229 0240 177 124017 WRT PASS TAB B STORE B INTO SECOND LOCATION
0230 0241 000 075736 RTN INC P P UPDATE P

0231 0242 220 074457 MPY READ INC M P M <= P; READ
0232 0243 300 000640 JSB IOFF INDIRECT JSB TO GET M <= ADDR OF OPERAND
0233 0244 000 075717 INC P P UPDATE P
0234 0245 017 101057 PASS S2 TAB S2 <= MULTIPLIER
0235 0246 017 127114 MPYX COV PASS S3 A S3<=A<MULTIPLICAND>; CLEAR OVFL
0236 0247 001 136517 ZERO B CLEAR B FOR MULTIPLY
0237 0250 017 142157 PASS L S2 L <= S2 <MULTIPLIER>
0238 0251 017 124255 RPT PASS CNTR B CLEAR COUNTER; SET REPEAT FF
0239 0252 104 124504 MPY R1 ADD B B MPY STEP (X16); (B,A)<=A TIMES L PLUS B
0240 0253 017 144757 PASS S3 TEST MULTIPLICAND
0241 0254 322 012731 JMP CNDX AL15 RJS **2 JUMP IF POSITIVE
0242 0255 003 024517 SUB B B UNDO LAST MPY STEP IF NEGATIVE
0243 0256 017 142757 PASS S2 TEST MULTIPLIER
0244 0257 322 003071 JMP CNDX AL15 RJS RETURN JUMP IF POSITIVE
0245 0260 017 144157 PASS L S3 L <= MULTIPLICAND
0246 0261 003 024536 RTN SUB B B B<=B MINUS L (CORRECTS FOR NEG. MULT)
0247 *
0248 *

```

```

0249 0262 220 074457 DIV READ INC M P M <= P; READ
0250 0263 300 000640 JSB IOFF INDIRECT JSB TO GET M <= ADDR OF OPERAND
0251 0264 000 075717 INC P P UPDATE P
0252 0265 010 001157 CMPS S4 TAB S4 <= DVSR(CM)::SAVE ORIG SIGN
0253 0266 010 047017 CMPS S1 S4 S1 <= DVSR
0254 0267 322 013471 JMP CNDX AL15 RJS **2 JMP IF DVSR NEGATIVE
0255 0270 000 047017 INC S1 S4 S1 <= DVSR(2CM)
0256 0271 017 140157 PASS L S1 L <= ABS VALUE(DVSR)
0257 0272 010 025117 CMPS S3 B S3 <= DVNDHI(2CM)
0258 0273 322 054071 JMP CNDX AL15 DIVS JMP IF DVND POSITIVE
0259 0274 017 144517 PASS B S3 IF DVND IS NEGATIVE...
0260 0275 010 027057 CMPS S2 A FORM DVND(2CM)
0261 0276 000 042557 INC A S2 IN B,A-REGISTER
0262 0277 321 014071 JMP CNDX COUT RJS DIVS *
0263 0300 000 044517 INC B S3 *
0264 0301 003 024753 DIVS SOV SUB B CHECK FOR DVSR TOO SMALL
0265 0302 322 000031 JMP CNDX AL15 RJS FETCH (<DVND TOO LARGE)
0266 0303 077 124502 LGS L1 PASS B B SHIFT OUT SIGN BIT OF FULL WORD
0267 0304 001 137054 COV ZERO S2 CLEAR OVFL,S2,& CNTR
0268 0305 017 142255 RPT PASS CNTR S2 AND SET RPTFF
0269 0306 123 024502 DIV L1 SUB B B DIV(16X); A<=QUO(POS); B<=REM*2
0270 0307 157 144142 LMF L1 PASS L S3 L <= FLG <= DVND SIGN(CM)
0271 0310 010 027017 CMPS S1 A S1 <= QUO(CM)
0272 0311 320 155071 JMP CNDX ONES QZERO IF QUO=0, THEN NO FURTHER TESTING
0273 0312 013 047057 XOR S2 S4 S2(15) <= EXPECTED SIGN OF QUO
0274 0313 322 014671 JMP CNDX AL15 RJS **2 JMP IF POSITIVE WAS EXPECTED
0275 0314 000 040557 INC A S1 ELSE A <= QUO(2CM)
0276 0315 017 142157 PASS L S2 L(15) <= EXPECTED SIGN OF QUO
0277 0316 013 026757 XOR A A COMPARE TO FINAL SIGN OF QUO
0278 0317 322 015071 JMP CNDX AL15 RJS **2 JMP IF OK
0279 0320 017 136753 SOV ELSE INDICATE OVERFLOW
0280 0321 017 124504 R1 PASS B B B <= (REM*2)/2
0281 0322 324 040031 JMP CNDX FLAG FETCH CHECK SIGN OF DVND
0282 0323 010 024517 CMPS B B IF NEG, THEN FORM 2-COMP OF
0283 0324 000 024536 RTN INC B B REM & STORE IN B

```



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0284 $ORIGIN=3308
0285 *****
0286 * EAU TABLE
0287 *****
0288 0330 320 007330 EAU TABLE JMP RRR
0289 0331 320 007570 JMP ASR
0290 0332 320 010030 JMP LSR
0291 0333 320 000030 JMP FETCH ILLEGAL JR CODE FOR EAU GROUP
0292 0334 320 010270 JMP RRL
0293 0335 320 010530 JMP ASL
0294 0336 320 010770 JMP LSL
0295 0337 320 012130 JMP MPY
0296 *****
0297 * MAC TABLE
0298 *****
0299 0340 321 145270 MACTABLO JMP FADD / FLOATING POINT
0300 0341 321 145330 JMP FSUB / FLOATING POINT
0301 0342 321 151070 JMP FMPY / FLOATING POINT
0302 0343 321 153130 JMP FDIV / FLOATING POINT
0303 0344 321 140030 JMP IFIX / FLOATING POINT
0304 0345 321 141270 JMP FLOAT / FLOATING POINT
0305 0346 320 060030 JMP %1400 *** PROBABLE FUTURE HP USE
0306 0347 320 060035 JMP J30 %1400 *** PROBABLE FUTURE HP USE
0307 0350 320 100030 JMP %2000 *** PROBABLE FUTURE HP USE
0308 0351 320 100035 JMP J30 %2000 *** PROBABLE FUTURE HP USE
0309 0352 320 120030 JMP %2400 *** PROBABLE FUTURE HP USE
0310 0353 320 120035 JMP J30 %2400 *** PROBABLE FUTURE HP USE
0311 0354 320 140030 JMP %3000 *** PROBABLE FUTURE HP USE
0312 0355 320 140035 JMP J30 %3000 *** PROBABLE FUTURE HP USE
0313 0356 320 160030 JMP %3400 *** PROBABLE FUTURE HP USE
0314 0357 320 160035 JMP J30 %3400 *** PROBABLE FUTURE HP USE
0315 *****

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0316 0360 321 000030 MACTABL1 JMP          X4000    *** PROBABLE FUTURE HP USE
0317 0361 321 000035      JMP J30      X4000    *** PROBABLE FUTURE HP USE
0318 0362 321 020030      JMP          X4400    *** PROBABLE FUTURE HP USE
0319 0363 321 020035      JMP J30      X4400    *** PROBABLE FUTURE HP USE
0320 0364 321 040030      JMP          X5000    *** PROBABLE FUTURE HP USE
0321 0365 321 040035      JMP J30      X5000    *** PROBABLE FUTURE HP USE
0322 0366 321 060030      JMP          X5400    *** PROBABLE FUTURE HP USE
0323 0367 321 060035      JMP J30      X5400    *** PROBABLE FUTURE HP USE
0324 0370 321 100030      JMP          X6000    *** PROBABLE FUTURE HP USE
0325 0371 321 100035      JMP J30      X6000    *** PROBABLE FUTURE HP USE
0326 0372 321 120030      JMP          X6400    *** PROBABLE FUTURE HP USE
0327 0373 321 120035      JMP J30      X6400    *** PROBABLE FUTURE HP USE
0328 0374 320 040030      JMP          X1000    *** RESERVED FOR CUSTOMER ONLY
0329 0375 320 040035      JMP J30      X1000    *** RESERVED FOR CUSTOMER ONLY
0330 0376 321 160035      JMP          X7400    *** RESERVED FOR CUSTOMER ONLY
0331 0377 321 161035      JMP J30      X7400    *** RESERVED FOR CUSTOMER ONLY
0332                                X7420    / RESERVED FOR HP USE
0333                                X7420    / RESERVED FOR HP USE
                                X7420    / BASE SET EXTENSION
                                X7420    / BASE SET EXTENSION.
*****
SEND
** NO ERRORS**

```

```

0001 $ORIGIN=4008
0002 *****
0003 *
0004 *      21MX MICRO-CODE
0005 *      MODULE 1
0006 *
0007 *****
0008 DISPLAYA EQU          X376
0009 DISPLAYT EQU          X367
0010 DISPLAYS EQU          X337
0011 INTERRUPT EQU         X0005
0012 *****
0013 *      MEMORY INITIALIZATION ROUTINE
0014 *****
0015 0400 322 161171 HALT      JMP CHDX MMLS      NGOOD      JUMP IF MEMORY NOT LOST
0016 0401 341 004617      IMM      HIGH NEU      X102      ENABLE SYSTEM MAP
0017 0402 347 101017      IMM      LOW S1        X340      S1 <= 2'S COMP OF 32
0018 0403 001 137057      PASS A      S2          CLR S2 (MAP ADDR)
0019 0404 017 142557      PASS B      S2          CLR A-REG
0020 0405 017 142517      PASS T      S2          CLR B-REG
0021 0406 017 142117      PASS T      S2          CLR T REG
0022 0407 353 077117      IMM      CMH1 S3      X337      S3 <= "LOAD ADDR REG" COMMAND
0023 0410 347 076264 LOSTLOOP IMM SHLT      LOW CNTR      X337      CNTR<=COMP OF 32; CLEAR RUN FF
0024 0411 017 144617      PASS NEU      S3          LOAD 0 INTO ADDR REG ON NEU
0025 0412 017 142620 MAPLOOP      MESP PASS NEU      S2          LOAD MAP IN NEU
0026 0413 000 043063      ICNT INC      S2          INC MAP ADDR
0027 0414 323 020531      JMP CHDX CNTB RJS      MAPLOOP      LOOP(*32)
0028 0415 001 137717      ZERO P          CLR P REG
0029 0416 160 074717      WRTE      INC PHM      P          M<=P; P<=P+1; WRITE ZERO DATA
0030 0417 322 020731      JMP CHDX AL15 RJS      *-1      LOOP UNTIL M=077777
0031 0420 000 041017      INC S1          S1          INC MAP CNTR
0032 0421 320 020431      JMP CHDX TBZ RJS      LOSTLOOP      LOOP (*32)
0033 0422 341 000617      IMM      HIGH NEU      X100      DISABLE ALL MAPS NOW...
0034 *****
0035 *      FRONT PANEL STANDARD SCAN ROUTINES
0036 *****
0037 0423 334 165531 NGOOD     JMP CHDX NSFP      CONTFP      JUMP IF NON-STANDARD FRONT PANEL
0038 0424 017 115752      FTCH PASS S      DSPL      S<=DISPLAY; INITIALIZE MEM. PROTECT
0039 0425 330 121371      JMP CHDX NSMG RJS      WAIT      JUMP IF "INSTR STEP" PRESSED
0040 0426 347 156357      IMM      LOW DSP1      DISPLAYT      ACTIVATE "T" INDICATOR IN DSP1
0041 0427 300 024270 WAIT      JSB          UPDATE      UPDATE DISPLAY WITH PROPER DATA
0042 0430 334 021431      JMP CHDX NSTB RJS      *          WAIT FOR BUTTON RELEASES
0043 0431 325 164231      JMP CHDX RUN          RUN
0044 0432 334 061471      JMP CHDX NSTB      *-1
0045 0433 017 136757 SCAN     NOP          SCAN FOR SWITCH PRESSED
0046 *                               NOP ONE CYCLE TO SET SWITCH CONDITIONS
0047 0434 332 122571      JMP CHDX MLT RJS      LEFT
0048 0435 331 023431      JMP CHDX MINC RJS      INC.M
0049 0436 331 123531      JMP CHDX MDEC RJS      DEC.M
0050 0437 333 025471      JMP CHDX MSTR RJS      STOREX
0051 0440 333 121371      JMP CHDX MRST RJS      WAIT
0052 0441 332 032231 SCANRT  JMP CHDX MRT RJS      RIGHT      JUMP IF "RIGHT" TO TEST FOR ENTRY
0053 *                               INTO SPECIAL DISPLAY ROUTINE.

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0054 0442 330 026171      JMP CHDX HLD RJS  LOADER
0055 0443 325 164231      JMP CHDX RUN      RUN
0056 0444 330 161431      JMP CHDX HSHG     WAIT+1    JMP IF "INSTR STEP" NOT PRESSED
0057 0445 335 040271      JMP CHDX INT      INTERRUPT SERVICE ANY PENDING INTERRUPT
0058 0446 017 176317      PASS DSPL S      DISPLAY <= S
0059 0447 220 074712      READ FTCH INC PHM P      DO STANDARD FETCH ROUTINE
0060 0450 017 136745      IOM
0061 0451 017 100411      CLFL PASS IR      TAB
0062 0452 220 020673      READ JTAB INC CM      ADR

0063 *****
0064 *      DISPLAY INDICATOR SHIFT ROUTINES
0065 *****
0066 0453 017 117004      LEFT      R1      PASS S1      DSPI      S1<=DSPI SHIFTED RIGHT ONE
0067 0454 321 122771      JMP CHDX ALD RJS LEFTA      JUMP IF DSPI WRAP-AROUND REQUIRED
0068 0455 017 140357      LEFTB      PASS DSPI S1      DSPI <= DSPI SHIFTED RIGHT ONE
0069 0456 320 021370      JMP          WAIT          JUMP TO STANDARD SCAN ROUTINES
0070 0457 347 076357      LEFTA      IMM          LOW DSPI DISPLAYS DSPI WRAP-AROUND A TO S
0071 0460 320 021370      JMP          WAIT          JUMP TO STANDARD SCAN ROUTINES
0072 0461 347 076157      RIGHT     IMM          LOW L      DISPLAYS
0073 0462 017 016750      STFL IOR      DSPI      SET FLAG; TEST DSPI
0074 0463 320 123331      JMP CHDX ONES RJS RIGHTA     JUMP IF WRAP-AROUND OF DSPI REQD
0075 0464 157 117002      LWF L1      PASS S1      DSPI      S1<=DSPI SHIFTED LEFT ONE
0076 0465 320 022670      JMP          LEFTB
0077 0466 347 174357      RIGHTA     IMM          LOW DSPI DISPLAYA DSPI WRAP-AROUND S TO A
0078 0467 320 021370      JMP          WAIT          JUMP TO STANDARD SCAN ROUTINES
0079 *****
0080 *      INC M, DEC M ROUTINES
0081 *****
0082 0470 000 023017      INC.M      INC S1      M      S1 <= M + 1
0083 0471 320 023570      JMP          DEC.M+1
0084 0472 007 123017      DEC.M      DEC S1      M      S1 <= M - 1
0085 0473 017 140457      PASS M      S1      M <= S1
0086 0474 320 021370      JMP UNCD      WAIT          JUMP TO STANDARD SCAN ROUTINE
0087 *****
0088 *      SPECIAL TEST TO EXIT SPECIAL DISPLAY LOOP
0089 *****
0090 0475 017 116417      LEFTR      PASS IR      DSPI      CHECK FOR "M" DSPI
0091 0476 327 122571      JMP CHDX IR2 RJS LEFT      JUMP IF "M" TO LEAVE SPECIAL CODE
0092 0477 347 166357      IMM          LOW DSPI X373 DSPI <= "M" (SHIFT FROM "T")
0093 0500 017 142317      PASS DSPL S2      SHOW POINTER ON DISPLAY
0094 0501 320 033130      JMP UNCD      WAITR      WAIT FOR BUTTON RELEASE IN SPECIAL CODE

0095 *****
0096 *      STORE AND UPDATE ROUTINES
0097 *****
0098 *      THE REGISTER INDICATED IN DSPI IS THE BIT POSITION WHICH IS
0099 *      LOW. ALL OTHER BITS ARE 1. THE ORDER (MSB TO LSB) IS
0100 *      S P T M B A
0101 *      THE INDICATED REGISTER IS DETERMINED BY LOADING DSPI INTO
0102 *      THE IR, AND JUMPING USING J30 TO GET TO THE APPROPRIATE
0103 *      STORE OR UPDATE ROUTINE. OTHER CODE IS INTERSPERSED
0104 *      FOR MAXIMUM CONTROL STORE EFFICIENCY
0105 0502 017 116417      STORE     PASS IR      DSPI
0106 0503 320 024035      JMP J30      X0500      JMP TO STORE SELECTED REGISTER
0107 0504 347 076357      RUN      IMM          LOW DSPI DISPLAYS DSPI <= "S". THE SAVE REGISTER IS
0108 *      ZERO AT THIS POINT SO THE NEXT RTN
0109 *      WILL INITIATE THE FETCH ROUTINE
0110 *****
0111 0505 017 116417      UPDATE     PASS IR      DSPI
0112 0506 320 025035      JMP J30      X520      JMP TO DISPLAY SELECTED REGISTER
0113 *****
0114 0507 177 114017      WRT      PASS TAB      DSPL      STORE T
0115 0510 000 023017      INC S1      M
0116 0511 000 040457      INC M      S1      INCREMENT M, SET TAB LOGIC
0117 0512 320 021430      JMP UNCD      WAIT+1
0118 0513 000 014476      RTN      INC M      DSPL      STORE M
0119 0514 017 115776      STORES    RTN      PASS S      DSPL      STORE S
0120 0515 017 114536      RTN      PASS B      DSPL      STORE B
0121 0516 017 114576      RTN      PASS A      DSPL      STORE A
0122 0517 347 136157      IMM          LOW L      X357      P OR S TO BE DISPLAYED
0123 0520 017 016757      IOR      DSPI      MASK OUT "S"
0124 0521 320 164631      JMP CHDX ONES STORES      JUMP IF "S" INDICATED
0125 0522 017 115736      RTN      PASS P      DSPL      STORE P

```

0126
0127
0128 0523 017 115013
0129 0524 321 165331
0130 0525 017 136754
0131 0526 017 136776
0132
0133 0527 220 022457
0134 0530 017 100336
0135
0136 0531 300 024130
0137 0532 320 021370
0138
0139 0533 017 122336
0140 0534 017 176336
0141 0535 017 124336
0142 0536 017 126336
0143 0537 347 136157
0144 0540 017 016757
0145 0541 320 165631
0146 0542 017 174336

```
*****
***** OVFL REG. STORE--PART OF SPECIAL DISPLAY ROUTINES *****
STORDO      SOV PASS S1 DSPL      CHECK DISPLAY
              JMP CNDX ALO      **2      CLEAR OVERFLOW
              COV
              RTH
*****
              READ      INC M      M      UPDATE T, READ M, SET TAB LOGIC
              RTH PASS DSPL TAB      DSPL <= MEM DATA
*****
              STOREX JSB      STORE      STORE ROUTINES END WITH RTH
              CONTFP JMP      WAIT      JUMP TO STANDARD SCAN ROUTINES
*****
              RTH PASS DSPL M      UPDATE M
              RTH PASS DSPL S      UPDATE S
              RTH PASS DSPL B      UPDATE B
              RTH PASS DSPL A      UPDATE A
              IMM      LOW L      357B      P OR S INDICATED
              IOR      DSPI      MASK OUT "S"
              JMP CNDX ONES      UPDATES
              RTH PASS DSPL P      UPDATE P
*****
```

0147
0148
0149
0150 0543 341 177053
0151 0544 353 137017
0152
0153 0545 347 000157
0154 0546 015 143717
0155 0547 010 075217
0156 0550 000 051217
0157 0551 017 142457
0158 0552 320 161371
0159 0553 177 150117
0160 0554 017 140157
0161 0555 223 043057
0162 0556 017 150157
0163 0557 013 004757
0164 0560 320 026271
0165
0166 0561 347 000157
0167 0562 347 164257
0168 0563 017 176417
0169 0564 017 177155
0170 0565 017 147144
0171 0566 013 147157
0172 0567 347 160157
0173 0570 004 147153
0174 0571 322 061371
0175
0176 0572 344 000257
0177 0573 017 174454
0178
0179 0574 017 131003
0180 0575 017 140163
0181 0576 015 131003
0182 0577 017 140163
0183 0600 015 131003
0184 0601 017 140163
0185 0602 012 031017
0186 0603 177 140117
0187 0604 000 023063
0188 0605 017 142457
0189 0606 344 000157
0190 0607 017 022757
0191 0610 320 127631

```
*****
***** 21MX ROM BOOTSTRAP MEMORY LOADER ROUTINE *****
*****
***** DETERMINE MEMORY SIZE, STARTING ADDR FOR LOADER *****
SIZE      IMM      LOW L      X300      FORM 0111111111111111 (MAX ADDR)
              AND P      S2      FORM 0001000000000000 (10K) IN S1
              CMPS S5      P      ADDR FOR LOADER *****
              INC S5      S5      FORM 1111111110000000 IN L
              PASS M      S2      FORM STARTING ADDR IN P
              JMP CNDX ONES      WAIT      FORM TWO'S COMP
              WRTE      PASS T      S5      OF SA IN S5
              PASS L      S1      PUT LAST ADDR INTO M
              READ      SUB S2      S2      TEST FOR NO READ/WRTE CAPABILITY
              PASS L      S5      PASS INTO T
              XOR      T      UPDATE LAST ADDR WHILE WAITING
              JMP CNDX TBZ RJS SIZE      TO RETRIEVE DATA
***** CHECK SELECT CODE IN S REG. *****      COMPARE WHAT WAS READ FROM MEM.
              IMM      LOW L      X300      TO DATA WRITTEN (S3)
              IMM      LOW CNTR X372      IF IT CHECKS, WE HAVE CORRECT STRT ADDR
              PASS IR      S      FORM 1111111110000000 IN L
              RPT      PASS S4      S      CNTR GETS -6
              R1      PASS S4      S4      SET UP LOADER SELECT BIT
              SANL S4      S4      SET UP S-REG FOR SHIFT
              IMM      LOW L      X370      SHIFT SELECT CODE INTO BITS(0-5)
              SOV ADD S4      S4      MASK OFF SEL. CODE
              JMP CNDX AL15      WAIT      FORM 11111111111000 (-10B) IN L
***** PREPARE FOR LOADER TRANSFER *****      SUB 10B FROM SEL CODE; SAVE IN SJ
              IMM      LOW CNTR X0      IF NEG RESULT, SCB < 10B; RTH W/ OVFL ON
              COV PASS M      P      CLEAR CNTR (ROM ADDR REG)
***** TRANSFER CONTENTS OF LOADER ROM TO MEMORY *****      PUT SA IN M; CLR OVFL = NO OPER ERR
LOOP1      L4      PASS S1      LDR      PASS XXXXXXXXXXXXXXXX INTO S1; CNTR=X00
              ICNT PASS L      S1      CNTR=X01
              L4      AND S1      LDR      FORM XXXXAAAABBBBXXXX IN S1; CNTR=X01
              ICNT PASS L      S1      CNTR=X10
              L4      AND S1      LDR      FORM AAAABBBBCCCCXXXX IN S1; CNTR=X10
              ICNT PASS L      S1      CNTR=X11
              HAND S1      LDR      FORM AAAABBBBCCCCDDDD (CMPL FORM)
              WRTE      PASS T      S1      WRITE INTO MEMORY
              ICNT INC S2      M      UPDATE MEM ADDR; CNTR=X00
              PASS M      S2      PASS NEW ADDR INTO M
              IMM      LOW L      X0      FORM 1111111000000000 IN L
              IOR      M      MASK M TO SEE IF LAST WORD OF LDR
              JMP CNDX ONES RJS LOOP1      IF M(0-8)=11111111, DON'T LOOP
*****
```

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0192 *****
0193 IMM LOW CNTR X300 SET UP COUNT TO FIND LAST WORD
0194 IMM STFL LOW S3 X037
0195 LWF L1 PASS S3 S3 FORM 11111100011111 IN S3
0196 0614 017 175017 PASS S1 P PASS SA INTO S1
0197 ***** CHECK INSTRUCTION IN MEMORY FOR I/O TYPE *****
0198 NUWRD READ PASS M S1 PASS SA INTO M & READ FIRST INSTR
0199 IMM HIGH L X013 FORM COMP OF 1111010000000000 IN L
0200 0617 017 105057 PASS S2 T SAVE WORD IN S2
0201 0620 013 143017 SAML S1 S2 MASK UPPER BITS FOR I/O TYPE
0202 0621 341 166157 IMM HIGH L X173 FORM 0111101111111111 IN L
0203 0622 013 040757 XOR S1 S1 NOW CHECK FOR I/O TYPE
0204 0623 320 171531 JMP CNDX ONES HTST IF MATCH OCCURS, JUMP OUT OF LOOP
0205 *****
0206 0624 000 023023 UPDT ICNT INC S1 M OTHERWISE UPDATE M IN S1
0207 0625 323 030671 JMP CNDX CNT8 RJS NUWRD LOOP BACK
0208 0626 017 146154 COV PASS L S4 PASS (SCB-10B) INTO L
0209 0627 004 143051 CLFL ADD S2 S2 CHNG SC OF DCPC CNTRL WORD
0210 0630 177 142117 WRTE PASS T S2 SAVE IN MEM
0211 0631 320 021370 JMP WAIT RETURN TO SCAN ROUTINE
0212 ***** UPDATE SELECT CODE IN I/O INSTRUCTION *****
0213 0632 017 144157 HTST PASS L S3 PASS 11111100011111 INTO L
0214 0633 014 042757 NSOL S2 BLEND TO CHECK FOR...000...OF HLT
0215 0634 320 171231 JMP CNDX ONES UPDT IF FOUND GET NEXT INSTR
0216 0635 347 016157 IMM LOW L X307 FORM 111111111000111 IN L
0217 0636 013 142757 SAML S2 S2 MASK BITS TO CHECK FOR SC < 10B
0218 0637 320 071231 JMP CNDX TBZ UPDT IF SO, RTN TO LOOP
0219 0640 017 146157 PASS L S4 PASS (SCB-10B) INTO L
0220 0641 004 143057 ADD S2 S2 ADD TO SC FROM INSTR
0221 0642 177 142117 WRTE PASS T S2 PASS INTO T AND WRITE INTO MEMORY
0222 0643 320 031230 JMP UPDT RTN TO LOOP

0223 *****
0224 * SPECIAL DISPLAY ROUTINES
0225 *****
0226 0644 017 116417 RIGHTR PASS IR DSP1 "RIGHT" PRESSED: IR <= DSP1
0227 0645 327 163071 JMP CNDX IR2 RIGHT JUMP IF M NOT SELECTED BY DSP1
0228 0646 017 115057 PASS S2 DSPL S2 <= DSPL (POINTER)
0229 0647 322 023071 JMP CNDX AL15 RJS RIGHT JUMP IF DSPL BIT 15 WASNT SET
0230 0650 347 156357 IMM LOW DSP1 X367 DSP1 <= "T"
0231 *****
0232 0651 006 042157 UPDATR OP3 L S2 CHECK DSPL BIT 14, STORE S2 IN L
0233 0652 322 074671 JMP CNDX AL15 NEUMAPS JUMP IF S2 BIT 14 = 1 TO UPDATE MEU
0234 0653 350 007093 IMM L4 CMH1 S1 X003 S1 <= MASK FOR REGISTERS = 140017B
0235 0654 015 141057 AND S2 S1 S2 <= S2 MASK OUT UNUSED BITS
0236 0655 017 142417 PASS IR S2 SET REGISTER SELECTION
0237 0656 333 073071 JMP CNDX NSTR READREG JUMP IF STORE BUTTON NOT PRESSED
0238 0657 300 037035 JSB J30 STOREG SELECTED REGISTER <= DISPLAY
0239 0660 320 033130 JMP UNCD WAITR WAIT FOR NEXT BUTTON
0240 0661 300 036035 READREG JSB J30 DSPLREG DISPLAY <= SELECTED REGISTER
0241 *****
0242 0662 334 033131 WAITR JMP CNDX NSTB RJS * WAIT FOR BUTTON RELEASE
0243 0663 325 164231 JMP CNDX RUN RUN JUMP IF RUN INDICATOR LIT
0244 0664 334 073171 JMP CNDX NSTB *-1 JUMP BACK IF NO BUTTON PRESSED
0245 *
0246 0665 017 136757 NOP WAIT ONE CYCLE FOR SETTING SWITCH CONDIT
0247 0666 332 123671 JMP CNDX NLT RJS LEFTR JUMP IF "LEFT" PRESSED
0248 0667 331 073531 JMP CNDX NINC NOTINC JUMP IF "INCH" NOT PRESSED
0249 0670 000 043057 INC S2 S2 INCREMENT POINTER
0250 0671 320 034030 JMP UNCD DECMR+1 DECMR+1
0251 0672 331 174231 JMP CNDX NDEC NOTDEC NOTDEC
0252 0673 340 000157 IMM HIGH L X000 CHECK FOR
0253 0674 015 142757 AND S2 S2 DECREMENT OF
0254 0675 320 033771 JMP CNDX TBZ RJS DECMR ZERO COUNT
0255 0676 000 143057 OP1 S2 S2 S2 OR L PLUS 1 (WRAP AROUND COUNT + 1)
0256 0677 007 143057 DEC S2 S2 DECREMENT POINTER
0257 0700 017 116417 PASS IR DSP1 IR <= DSP1
0258 0701 327 172471 JMP CNDX IR2 UPDATR JUMP IF M NOT INDICATED
0259 0702 017 142317 PASS DSPL S2 UPDATE DISPLAY
0260 0703 320 033130 JMP UNCD WAITR WITH NEW POINTER VALUE AND JUMP
0261 0704 333 134031 JMP CNDX NRST RJS DECMR+1 JUMP IF "DISPLAY" PRESSED
0262 0705 333 074471 JMP CNDX NSTR *-4 JUMP IF STORE NOT PRESSED
0263 0706 017 116417 PASS IR DSP1
0264 0707 327 125471 JMP CNDX IR2 RJS STOREX JUMP IF M SELECTED. LEAVE SPECIAL MODE
0265 0710 320 032470 JMP UNCD UPDATR M NOT SELECTED
0266 0711 320 022070 JMP UNCD SCANRT JUMP TO STD ROUTINES

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0267 *****
0268 0712 347 172417 STOREE IMM LOW IR X375 SET UP SRC TYPE ER* SHIFT
0269 0713 017 114741 SRG2 PASS DSPL SET E ACCORDING TO DSPL BIT 0
0270 0714 017 136776 RTN
0271 *****
0272 ***** MEU MAP MANIPULATIONS *****
0273 0715 346 000157 MEUMAPS IMM LOW L X200 S1 <= MASK OF LOW 7 BITS
0274 0716 013 143017 SANL S1 S2
0275 0717 340 176157 IMM HIGH L X077 L <= 037777B
0276 0720 016 141057 SONL S2 S1 S2 <= MASK OUT BITS 13 TO 8
0277 0721 343 076157 IMM HIGH L X337 OR IN BIT 13
0278 0722 016 141017 SONL S1 S1
0279 0723 017 140617 PASS MEU S1 SEND MAP NO. TO MEU
0280 0724 333 075371 JMP CHDX NSTR READMAP JUMP IF STORE NOT PRESSED
0281 0725 017 114620 MESP PASS MEU DSPL MEU MAP <= DISPLAY
0282 0726 320 033130 JMP UNCD WAITR
0283 0727 017 134320 READMAP MESP PASS DSPL MEU DISPLAY <= MEU MAP
0284 0730 320 033130 JMP UNCD WAITR
0285 ***** SIMULATED LIA 4 I/O INSTRUCTION *****
0286 0731 344 010417 DSPLCIR IMM LOW IR X004 TO READ CIR *****
0287 0732 017 136762 IOG SET UP SEL CODE 4 IN IR
0288 0733 017 136757 NOP INITIATE I/O CYCLE AT TIME T2
0289 0734 017 110336 RTN PASS DSPL IOI WAIT FOR TIME T4
CIR TO DISPLAY. DONT ISSUE IAK

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0290 $ORIGIN=7408
0291 *****
0292 * SHORT SUBROUTINES TO STORE/DISPLAY SELECTED REGISTERS
0293 *****
0294 0740 017 170336 DSPLREG RTN PASS DSPL X PASS REG TO FRONT PANEL AND RETURN
0295 0741 017 172336 RTN PASS DSPL Y
0296 0742 017 112336 RTN PASS DSPL CNTR
0297 0743 017 144336 RTN PASS DSPL S3
0298 0744 017 146336 RTN PASS DSPL S4
0299 0745 017 150336 RTN PASS DSPL S5
0300 0746 017 152336 RTN PASS DSPL S6
0301 0747 017 154336 RTN PASS DSPL S7
0302 0750 017 156336 RTN PASS DSPL S8
0303 0751 017 160336 RTN PASS DSPL S9
0304 0752 017 162336 RTN PASS DSPL S10
0305 0753 017 164336 RTN PASS DSPL S11
0306 0754 017 166336 RTN PASS DSPL S12
0307 0755 320 035470 JMP UNCD DSPLCIR
0308 0756 343 176336 IMM RTN HIGH DSPL 377B
0309 0757 343 176336 IMM RTN HIGH DSPL 377B
0310 *****
0311 0760 017 115636 STOREG RTN PASS X DSPL STORE INTO REG FROM FRONT PANEL
0312 0761 017 115676 RTN PASS Y DSPL
0313 0762 017 114276 RTN PASS CNTR DSPL
0314 0763 017 115136 RTN PASS S3 DSPL
0315 0764 017 115176 RTN PASS S4 DSPL
0316 0765 017 115236 RTN PASS S5 DSPL
0317 0766 017 115276 RTN PASS S6 DSPL
0318 0767 017 115336 RTN PASS S7 DSPL
0319 0770 017 115376 RTN PASS S8 DSPL
0320 0771 017 115436 RTN PASS S9 DSPL
0321 0772 017 115476 RTN PASS S10 DSPL
0322 0773 017 115536 RTN PASS S11 DSPL
0323 0774 017 115576 RTN PASS S12 DSPL
0324 0775 017 106336 RTN PASS DSPL CIR
0325 *
0326 0776 320 025170 JMP UNCD STOROO LOAD CIR FROM INT. REQUEST LINES
0327 0777 320 034530 JMP UNCD STOREE AND ISSUE INTERRUPT ACKNOWLEDGE
0328 *****
0329 $END *****
** NO ERRORS**

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0001 $ORIGIN=7000B
0002 *****
0003 *
0004 *      21MX MICRO-CODE
0005 *      MODULE 14: FLOATING POINT INSTRUCTIONS
0006 *
0007 *****
0008 INDIRECT EQU          X0015
0009 MPYX      EQU          X0246
0010 *****
0011 *
0012 7000 017 125414 IFIX      COV PASS S9      B      CLEAR THE OVFL AND PUT EXP IN S9
0013 7001 321 100171      JMP CNDX AL0 RJS  **2      TEST FOR NEG EXP
0014 7002 001 136576      RTN ZERO A      IF EXP<0 WE CAN'T FIX
0015 7003 017 126517      PASS B      A      PUT HIBITS IN B-REG
0016 7004 344 000157      IMM      LOW L      X000      PUT 'UP-B' MASK IN L
0017 7005 015 160557      AND A      S9      MASK LEAST SIG. 8 BITS INTO A
0018 7006 015 161457      AND S10 S9      SAVE BITS FOR ROUND-OFF
0019 7007 013 161404      R1 SANL S9      S9      MASK EXP INTO S9 WITHOUT SIGN
0020 7010 347 140157      IMM      LOW L      X360      PUT -20(BB) INTO L
0021 7011 004 161413      SOV ADD S9      S9      CHECK TO SEE IF EXP TOO LARGE
0022 7012 320 140771      JMP CNDX ONES      NOSHIFT      OR IF NO SHIFT REQUIRED
0023 7013 322 005231      JMP CNDX AL15 RJS OVER      IF SO THEN WE CAN'T FIX
0024 7014 000 061417      INC S9      S9      START LOOP TO SHIFT DIGITS
0025 7015 017 160255      RPT PASS CNTR S9      PASS 8 OF SHIFTS INTO CNTR
0026 7016 037 124504      ARS R1 PASS B      B      32-BIT SHIFT
0027 7017 017 126157      NOSHIFT      PASS L      A      HOLD LEFTOVER BITS IN L
0028 7020 017 124554      COV PASS A      B      PUT INTEGER INTO A-REG
0029 7021 322 017631      JMP CNDX AL15 RJS RTNFP      TEST FOR NEG INTEGER
0030 7022 017 062757      IOR      S10      IF NEG THEN CHECK FOR TRUNC. BITS
0031 7023 320 057631      JMP CNDX TBZ      RTNFP      IF ALL ZEROS WE ARE DONE
0032 7024 000 024576      RTN INC A      B      OTHERWISE INC THE INTEGER & RTN
0033 *****
0034 *****
0035 7025 017 126517 FLOAT      PASS B      A      PUT INTEGER IN B-REG
0036 7026 001 136557      ZERO A      CLEAR A-REG
0037 7027 357 141417      IMM      CHLO S9      X360      STORE +15(B10) IN EXP REG
0038 7030 321 142530      JMP      PACK
0039 *
0040 *****
0041 *
0042 7031 017 101317 FLD      PASS S7      TAB      STORE HIBITS IN S7
0043 7032 000 023017      INC S1      M      INC ADDR8 FOR NEXT READ
0044 7033 340 000157      IMM      HIGH L      X000      STORE 'LO-B' MASK IN L
0045 7034 220 040457      READ      INC M      S1      READ SECOND HALF OF WRD
0046 7035 015 125417      AND S9      B      MEANWHILE, MASK EXP OF WRD1 INTO S9
0047 7036 017 101217      PASS S5      TAB      STORE WRD2 LOBITS/EXP IN S5
0048 7037 013 125457      SANL S10 B      MASK LOBITS OF WRD1 INTO S10
0049 7040 013 151257      SANL S6      S5      MASK LOBITS OF WRD2 INTO S6
0050 7041 015 151204      R1 AND S5      S5      MASK EXP OF WRD2 INTO S11 WITHOUT SGN
0051 7042 321 102271      JMP CNDX AL0 RJS **3      IF SIGN WAS POS, JMP
0052 7043 346 000157      IMM      LOW L      X200      OTHERWISE PUT -200(BB) INTO L
0053 7044 004 151217      ADD S5      S5      ADD TO EXP OF WRD2
0054 7045 017 161404      R1 PASS S9      S9      MASK EXP OF WRD1 INTO S9 WITHOUT SIGN
0055 7046 321 102471      JMP CNDX AL0 RJS **3      IF SIGN WAS POS, JMP
0056 7047 346 000157      IMM      LOW L      X200      OTHERWISE PUT -200(BB) INTO L
0057 7050 004 161417      ADD S9      S9      ADD TO EXP OF WRD1
0058 7051 017 127536      RTN PASS S11 A      PUT HIBITS OF WRD1 INTO S3 & RTN
0059 *****
0060 *****
0061 7052 017 126154 PACK      COV PASS L      A      CLR OVFL AND PUT WRD1 LOBITS INTO L
0062 7053 001 137457      ZERO S10      CLEAR COUNTER REG
0063 7054 017 024757      IOR      B      PASS THRU ALU WITH HIBITS
0064 7055 320 057631      JMP CNDX TBZ      RTNFP      IF A/B IS ZERO, RTN
0065 7056 346 003517      IMM      LOW S11 X201      STORE -177(BB) IN S11
0066 7057 037 124742      ARS L1 PASS B      B      TEST IF NUMBER IS NORMALIZED
0067 7060 325 043231      JMP CNDX OVFL      RND      IF SO, JMP TO ROUNDING ROUTINE
0068 7061 077 124502      LGS L1 PASS B      B      IF NOT, DO 32-BIT LEFT-SHIFT
0069 7062 000 063457      INC S10      S10      INC THE EXP CNTR
0070 7063 321 142770      JMP      HRMLZ      GO BACK TO CHECK FOR NORMAL NUMBER
0071 7064 322 043331      RND      JMP CNDX AL15 **2      SINCE B WAS JUST PASSED THRU ALU
0072 7065 007 165517      DEC S11      S11      CHECK SGN & ADJUST ROUND OFF
0073 7066 017 164154      COV PASS L      S11      PUT 'ROUND' INTO L
0074 7067 003 026557      SUB A      A      ACTUALLY, ADD 200(BB) TO LOBITS
0075 7070 321 004171      JMP CNDX COUT RJS XPMT      IF NO COUT FROM LOBITS, OK, JMP
0076 7071 340 000157      IMM      HIGH L      X0      CLR L(15) FOR OVERFLOW
0077 7072 240 024517      ENV      INC B      B      IF COUT, INC HIBITS AND CHECK FOR OVFL
0078 7073 325 003771      JMP CNDX OVFL RJS **4      IF NO OVFL, OK, JMP
0079 7074 017 124504      R1 PASS B      B      OVFL IMPLIES B/A= 1000...
0080 7075 000 061414      COV INC S9      S9      SO WE SHIFT 8 TO FORM 0100...&

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0081	7076	321	144170	JMP			XPNT	BUMP EXP, THEN JMP
0082	7077	037	124742	ARS	L1	PASS	B	IF B NEQ 100...CHECK IF B=111...
0083	7100	325	044171	JMP	CNDX	OVFL	XPNT	IF NOT, JMP
0084	7101	077	124502	LGS	L1	PASS B	B	RE-NORMALIZE
0085	7102	000	063457			INC S10	S10	
0086	7103	017	162153	XPNT	SOV	PASS L	S10	CLR OVFL AND PUT EXP INTO L
0087	7104	003	061417			SUB S9	S9	SUB CALC EXP FROM ORIG EXP
0088	7105	346	000157	IMM		LOW L	X200	PUT -200(B8) INTO L
0089	7106	003	060757			SUB	S9	TEST FOR EXP UNDERFLO
0090	7107	322	045131	JMP	CNDX	AL15	UNFLO	IF S0, JMP
0091	7110	004	160757			ADD	S9	TEST FOR EXP OVERFLOW
0092	7111	322	017671	JMP	CNDX	AL15	RJS	OVFLO
0093	7112	157	160742	LWF	L1	PASS	S9	IF S0, JMP (TO 7375)
0094	7113	157	161402	LWF	L1	PASS S9	S9	PASS EXP SIGN INTO FLAG-REG
0095	7114	340	000157	IMM		HIGH L	X000	SHIFT EXP WITH SIGN
0096	7115	015	161457			AND S10	S9	STORE 'LO-B' MASK IN L
0097	7116	013	127417			SANL S9	A	MASK EXP INTO S10
0098	7117	017	124557			PASS A	B	MASK LOBITS INTO S9
0099	7120	017	160154		COV	PASS L	S9	PUT HIBITS INTO A-REG
0100	7121	017	062536		RTN	10R B	S10	PUT LOBITS INTO L
0101	7122	001	136557	UNFLO		ZERO A		COMBINE WITH EXP AND STORE IN B-REG
0102	7123	001	136536		RTN	ZERO B		CLEAR A-REG; OVFL=1
0103				*				NOW CLR B-REG AND RTN
0104	7124	341	176576	OVER	IMM	RTN	HIGH A	X177
0105								SET UP ERROR CONDITION IN A

0106 *****

0107 7125 017 136750 FADD STFL

0108 *

0109	7126	220	074457	FSUB	READ	INC	M	P	PASS P INTO M TO READ ADDR OF WRD2
0110	7127	300	000670		JSB			INDIRECT	CHECK FOR INDIRECTS
0111	7130	301	141470		JSB			FLD	UNPACK WRDS INTO SCRATCH REGS
0112	7131	017	154517			PASS B		S7	CHECK FOR WRD2=0
0113	7132	320	005631		JMP	CNDX	TBZ	RJS	IF NOT,CONTINUE
0114	7133	346	001217		IMM	LOW	S5	X200	IF SO,MAKE EXP MOST NEG (-200,B8)
0115	7134	017	164757			PASS		S11	CHECK FOR WRD1=0
0116	7135	320	005771		JMP	CNDX	TBZ	RJS	IF NOT,CONTINUE
0117	7136	346	001417		IMM	LOW	S9	X200	IF SO,MAKE EXP MOST NEG (-200,B8)
0118	7137	324	046531		JMP	CNDX	FLAG	DIFR	IF DOING ADD,SKIP AHEAD
0119	7140	010	024517			CMPS B		B	FORM 2-COMP OF HIBITS IN B
0120	7141	010	053257			CMPS S6		S6	FORM 2-COMP OF
0121	7142	000	053257			INC S6		S6	LOBITS OF WRD2
0122	7143	321	006531		JMP	CNDX	COUT	RJS	IF COUT OCCURS
0123	7144	000	024517			INC B		B	BUMP HIBITS
0124	7145	322	006531		JMP	CNDX	AL15	RJS	CHECK SGN; IF POS,JMP
0125	7146	017	124742			L1	PASS	B	IF NEG,CHECK FOR MOST
0126	7147	320	006531		JMP	CNDX	TBZ	RJS	NEG #(100...)
0127	7150	017	124504			R1	PASS B	B	IF SO,SHIFT BACK (010...)
0128	7151	000	051217			INC S5		S5	&BUMP EXP
0129	7152	017	152557	DIFR		PASS A		S6	
0130	7153	017	150157			PASS L		S5	FIND DIFF IN EXPS
0131	7154	003	061351			CLFL	SUB	S8	&STORE IN S8; FLC=0
0132	7155	320	047731		JMP	CNDX	TBZ		ADD2
0133	7156	322	047131		JMP	CNDX	AL15		RVR5
0134	7157	010	057357			CMPS S8		S8	IF NEG,WRD2>WRD1
0135	7160	000	057357			INC S8		S8	FORM -DIFF
0136	7161	321	147430		JMP				& STORE -DIFF IN S8
0137	7162	017	124157	RVR5		PASS L		B	SWANPCHK
0138	7163	017	164517			PASS B		S11	HOLD B IN L
0139	7164	017	162557			PASS A		S10	WRD1<WRD2; FILL B,A
0140	7165	015	037517			PASL S11			WITH S11,S10
0141	7166	017	153457			PASS S10		S6	ALSO FILL S11,S10,S9
0142	7167	017	151417			PASS S9		S5	WITH B,S6,S5
0143	7170	347	120157	SWANPCHK	IMM	LOW	L	X350	FORM -30(B8) IN L
0144	7171	003	056757			SUB		S8	IF -DIFF>-31,RTN WITH LARGER #
0145	7172	322	050671		JMP	CNDX	AL15		OUT
0146	7173	037	124504	SHIFT	ARS	R1	PASS B	B	JMP TO RESTORE A,B
0147	7174	000	057357			INC S8		S8	NOW START SHIFT LOOP
0148	7175	320	007571		JMP	CNDX	TBZ	RJS	INC COUNTER
0149				*					LOOP UNTIL DONE
0150									

***** CONTINUED ON NEXT PAGE *****

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0151
0152 7176 017 162154 *
0153 7177 004 126537 ADD2 COV PASS L S10 PASS LOBITS INTO L
0154 7200 321 010171 JMP CNDX COUT RJS *+3 ADD & CHECK FOR COUT
0155 7201 340 000157 IMM HIGH L X0 IF NOT JMP
0156 7202 240 024517 ENV INC B B CLR L(15) FOR OVFL
0157 7203 017 164151 CLFL PASS L S11 IF SO, INC HIBITS & ENABLE OVFL
0158 7204 244 124517 ENV ADD B B FLG=0
0159 7205 325 010571 JMP CNDX OVFL RJS PKSUB ADD HIBITS AND ENABLE OVFL
0160 7206 322 050431 JMP CNDX AL15 *+2 IF NO OVFL, RETURN
0161 7207 017 136750 STFL SO FLG=U IF ALU15=0
0162 7210 157 124504 LWF R1 PASS B B DO FULLWRD SHIFT
0163 7211 157 126544 LWF R1 PASS A A USING FLG REG TO INJECT SGH
0164 7212 000 061417 INC S9 S9 BUMP EXP
0165 7213 301 142530 PKSUB JSB RTN INC P P REPACK A,B REGS
0166 7214 000 075736 OUT PASS B S11 INC P AND RETURN
0167 7215 017 164517 PASS A S10 PASS MUCH LARGER WRD INTO B,A
0168 7216 017 162557 JSB RTN INC P P
0169 7217 301 142530
0170 7220 000 075736
0171
*****
0172
0173 7221 220 074457 FMPY READ INC M P PASS P INTO M TO READ ADDR OF WRD2
0174 7222 300 000670 JSB INDIRECT CHECK FOR INDIRECTS
0175 7223 301 141470 JSB FLD STORE ARGS IN SCRATCH REGS
0176 7224 000 061417 INC S9 S9
0177 7225 017 150157 PASS L S5 FORM EXP1+EXP2+1
0178 7226 004 161417 ADD S9 S9 AND SAVE IN S9
0179 7227 017 162544 R1 PASS A S10 FORM (WRD1 LOBITS)/2 IN A
0180 7230 017 155057 PASS S2 S7 PASS WRD2 HIBITS INTO S2
0181 7231 300 012330 JSB MPYX JMP TO MPY SUB & RTN WITH
0182 7232 017 125217 PASS S5 S8 HIBITS IN B; SAVE IN S5
0183 7233 017 163057 PASS S2 S11 PASS WRD1 HIBITS INTO S2
0184 7234 017 127517 PASS S11 A LOBITS INTO A; SAVE IN S11
0185 7235 017 152544 R1 PASS A S6 FORM (WRD2 LOBITS)/2 IN A
0186 7236 300 012330 JSB MPYX JMP TO MPY SUB & RTN WITH
0187 7237 017 126157 PASS L A LOBITS IN A; PASS INTO L
0188 7240 004 164557 ADD A S11 ADD BOTH LOBITS & CHK FOR COUT
0189 7241 321 012171 JMP CNDX COUT RJS *+2 (ELSE TRUNCATE DIGITS)
0190 7242 000 024517 INC B B IF COUT, BUMP HIBITS
0191 7243 017 124157 PASS L B ADD HIBITS AND SAVE IN S11
0192 7244 004 151517 ADD S11 S5
0193 7245 017 154557 PASS A S7 PASS WRD2 HIBITS INTO A
0194 7246 300 012330 JSB MPYX JMP TO MPY SUB & RTN WITH
0195 7247 017 126544 R1 PASS A A LOBITS IN A; SAVE LOBITS/2
0196 7250 017 126154 COV PASS L A ADD LOBITS/2 TO HIBITS SUM &
0197 7251 244 164542 ENV L1 ADD A S11 SHFT L1 TO REORIENT
0198 7252 322 012671 JMP CNDX AL15 RJS *+3 CHECK FOR CARRY INTO OR
0199 7253 325 052771 JMP CNDX OVFL *+4 BORROW FROM HIBITS &
0200 7254 007 124517 DEC B B ADJUST ACCORDINGLY
0201 7255 301 142530 JSB PACK
0202 7256 000 075736 RTN INC P P
0203 7257 000 024517 INC B B CAN'T OVFL FROM HIBITS
0204 7260 301 142530 JSB PACK
0205 7261 000 075736 RTN INC P P
0206
*****
0207
0208 7262 220 074457 FDIV READ INC M P PASS P INTO M TO READ ADDR OF WRD2
0209 7263 300 000670 JSB INDIRECT CHECK FOR INDIRECTS
0210 7264 301 141470 JSB FLD
0211 7265 010 054554 COV CNPS A S7 PASS WRD2 HIBITS & CHECK
0212 7266 320 156131 JMP CNDX ONES DBYZR FOR DIV BY ZERO
0213 7267 322 053471 JMP CNDX AL15 *+2 SINCE WE USE SAME DVSR, MAKE POS
0214 7270 000 027313 SOV INC S7 A NOW & SAVE SGH IN OVFL
0215 7271 017 150157 PASS L S5 FORM EXP1-EXP2+1
0216 7272 003 061417 SUB S9 S9 & SAVE IN S9
0217 7273 000 061417 INC S9 S9
0218 7274 017 162557 PASS A S10 FILL B,A WITH WRD1 AS DVND
0219 7275 017 164517 PASS B S11 & PRESIFT TO AVOID OVFL
0220 7276 037 124504 ARS R1 PASS B B
0221 7277 301 156370 JSB DIVX JMP TO SPECIAL DIV SUB
0222 7300 017 127217 PASS S5 A SAVE QUO1 IN S5
0223 7301 017 124757 PASS B A PASS QUO & CHECK FOR ODD/EVEN
0224 7302 321 114231 JMP CNDX ALD RJS *+2 TO SIMULATE FIRST
0225 7303 007 124517 DEC B B LEFT SHIFT IN DIV ROUTINE
0226 7304 001 136557 ZERO A CLR DVND LOBITS; DVSR SAME

```

```

0227 7305 301 156370
0228 7306 017 127517
0229 7307 017 152504
0230 7310 017 124504
0231 7311 001 136557
0232 7312 301 156370
0233 7313 010 026557
0234 7314 000 026557
0235 7315 017 151057
0236 7316 300 012330
0237 7317 017 125317
0238 7320 001 136517
0239 7321 017 164757
0240 7322 322 015231
0241 7323 016 036517
0242 7324 017 154757
0243 7325 322 015371
0244 7326 007 124517
0245 7327 017 155302
0246 7330 017 154542
0247 7331 017 126157
0248 7332 004 164557
0249 7333 321 015671
0250 7334 000 024517
0251 7335 077 124502
0252 7336 017 150157
0253 7337 004 124517
0254 7340 301 142530
0255 7341 000 075736
0256 7342 001 136557
0257 7343 352 000517
0258 7344 017 125417
0259 7345 301 142530
0260 7346 000 075736
0261

```

```

JSB      PASS S11      DIVX
R1      PASS B        A
R1      PASS B        S6
        ZERO A        8
JSB      CMPS A        DIVX
        INC A        A
        PASS S2      S5
        MPYX        B
JSB      PASS S7      S11
        ZERO B      *+2
        PASS        S7
JMP      CHDX AL15 RJS *+2
        ONE B        B
        PASS        S7
JMP      CHDX AL15 RJS *+2
        DEC B        B
        L1 PASS S7    S7
        L1 PASS A     S7
        PASS L        A
        ADD A        S11
JMP      CHDX COUT RJS *+2
        INC B        B
        LGS L1 PASS B  B
        PASS L        S5
        ADD B        B
        JSB          PACK
        RTN INC P      P
DBYZR    IMM          X200
        CMHI B        B
        PASS S9      PACK
        JSB          PACK
        RTN INC P      P
        *****
JMP TO SPEC DIV SUB
SAVE QUO2 IN S11
FORM (WRD2 LOBITS)/4 IN
B(=DVND HIBITS)
CLR DVND LOBITS; DVSR SAME
JMP TO SPEC DIV SUB
FORM 2-COMP OF QUO3
AS MPLR
PASS QUO1 AS MCND
JMP TO MPY SUB
SAVE PRDD HIBITS IN S5
PRE-CLR B
CHECK SGN OF QUO2
& EXTEND AS ALL 3'S(POS)
OR ALL 1'S(NEG)
CHECK SGN OF -QUO1+QUO3
IF NEG,SUB 1 FROM B
REDIRECT PROD (ADJUST EXP, REALLY)
ADD TO QUO2
IF COUT OCCURRED
BUMP HIBITS OF RESULT
SHIFT FULLWRD TO ORIENT RESULT
ADD QUO1 TO HIBITS
CLR LOBITS
FORM 0111111100000000 IN HIBITS
ALSO PASS INTO EXP
*****

```

```

0262
0263
0264
0265
0266
0267
0268
0269
0270
0271
0272 7347 344 000257
0273 7350 157 124742
0274 7351 322 016771
0275 7352 010 024517
0276 7353 010 026557
0277 7354 000 026557
0278 7355 321 016771
0279 7356 000 024517
0280 7357 017 154155
0281 7360 123 024502
0282 7361 017 124504
0283 7362 324 017271
0284 7363 010 024517
0285 7364 000 024517
0286 7365 325 057471
0287 7366 324 017631
0288 7367 010 026557
0289 7370 000 026576
0290 7371 324 057631
0291 7372 010 026557
0292 7373 000 026576
0293
0294
0295 7374 017 136776
0296
0297 7375 016 036544
0298 7376 347 174536
0299
0300
** NO ERRORS**

```

```

*****
*
*   THIS IS A SPECIAL SUB FOR F.P. DIV
*   IT ASSUMES THAT DVSR IS ALWAYS POS
*   THAT ORIG DVSR SGN IS IN OVFL REG
*   THAT YOU HAVE PREVIOUSLY DONE FIRST LEFT SHIFT
*   AND THAT NO ERROR COND NEED BE CHECKED
*   (BUT IT IS FAST)
*
*****
DIVX    IMM    LOW    CNTR  X0    CLR CNTR
LWF    L1     PASS    B      CHECK FOR NEG DVND & SAVE SGN IN FLG
JMP    CHDX  AL15 RJS  READY  IF POS, WE ARE READY
        CMPS B      B      COMP HIBITS
        CMPS A      A      COMP LOBITS
        INC A      A      FORM 2-COMP OF LOBITS
JMP    CHDX  COUT RJS  READY  IF NO COUT, OK
        INC B      B      ELSE BUMP HIBITS
READY   RPT    PASS L  S7     PASS DVSR INTO L; SET RPTFF
DIV     L1     SUB B    B      PERFORM DIV STEP(16X)
        R1     PASS B  B      FORM REM IN B
JMP    CHDX  FLAG RJS  *+3    IF REM SGN IS TO BE NEG
        CMPS B  B      (DETERMINED BY DVND), THEN
        INC B  B      FORM 2-COMP IN B
JMP    CHDX  OVFL    *+4     CHECK ORIG DVSR SGN; IF POS.
JMP    CHDX  FLAG RJS  RTNFP  LOOK FOR NEG DVND
        CMPS A  A      WHICH MEANS FORM
        RTN    INC A  A      NEG QUO IN A & RTN
JMP    CHDX  FLAG    RTNFP  ELSE IF NEG, LOOK FOR POS DVND
        CMPS A  A      WHICH MEANS FORM
        RTN    INC A  A      NEG QUO IN A & RTN
*****
*
RTNFP    RTN
*
OVFLO    R1     ONE A      PUT MOST POS # AND MOST POS EXP
        IMM    RTN  LOW B  X376  INTO A,B-RECS; OVFLO=1
*****
$END

```

```

0001          $ORIGIN=7400
0002          *****
0003          *
0004          *      21MX MICRO-CODE
0005          *      MODULE 15: EXTENDED INSTRUCTION GROUP
0006          *
0007          *****
0008          FETCH      EQU      X0000
0009          *****
0010          *      JUMP TABLES - ENTERED FROM BASE SET
0011          *****
0012          7400 321 163170          JMP          EADRX      SAX/SBX
0013          7401 017 103636          RTN      PASS X      CAB      CAX/CBX
0014          7402 321 163170          JMP          EADRX      LAX/LBX
0015          7403 321 163030          JMP          EADR      STX
0016          7404 017 170076          RTN      PASS CAB X      CXA/CXB
0017          7405 321 163030          JMP          EADR      LDY
0018          7406 321 163030          JMP          EADR      ADX
0019          7407 321 164070          JMP          XABX      XAX/XBX
0020          7410 321 163630          JMP          EADRY      SAY/SBY
0021          7411 017 103676          RTN      PASS Y      CAB      CAY/CBY
0022          7412 321 163630          JMP          EADRY      LAY/LBY
0023          7413 321 163030          JMP          EADR      STY
0024          7414 017 172076          RTN      PASS CAB Y      CYA/CYB
0025          7415 321 163030          JMP          EADR      LDY
0026          7416 321 163030          JMP          EADR      ADY
0027          7417 321 164230          JMP          XABY      XAY/XBY
0028          7420 321 164370          JMP          ISX
0029          7421 321 164670          JMP          DSX
0030          7422 321 177070          JMP          JLY
0031          7423 321 172530          JMP          LBT
0032          7424 321 171630          JMP          SBT
0033          7425 321 173230          JMP          MBT
0034          7426 321 175130          JMP          CBT
0035          7427 321 174030          JMP          SFB
0036          7430 321 164530          JMP          ISY
0037          7431 321 165030          JMP          DSY
0038          7432 321 177370          JMP          JPY
0039          7433 321 167330          JMP          SBSCBS      SBS
0040          7434 321 167330          JMP          SBSCBS      CBS
0041          7435 321 166630          JMP          TBS
0042          7436 321 170230          JMP          CMW
0043          7437 321 171030          JMP          MVW

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0044          *****
0045          *      INDEX REGISTER INSTRUCTIONS
0046          *****
0047          *      DISPLACEMENT FROM FINISH CORRESPONDS TO
0048          *      DISPLACEMENT FROM 7400B FOR INSTRS. LISTED
0049          *      IN COMMENT FIELD BELOW.
0050          7440 000 022461          FINISH      MPCK INC M      M      SAX/SBX
0051          7441 177 102036          WRTE RTN PASS TAB CAB      CAX/CBX
0052          7442 017 100076          RTN      PASS CAB TAB      LAX/LBX
0053          7443 000 022461          MPCK INC M      M      STX
0054          7444 177 170036          WRTE RTN PASS TAB X      CXA/CXB
0055          7445 017 101636          RTN      PASS X      TAB      LDY
0056          7446 017 100157          PASS L      TAB      ADX
0057          7447 264 171636          ENVE RTN ADD X      X      SAY/SBY
0058          7450 000 022461          MPCK INC M      M      CAX/CBX
0059          7451 177 102036          WRTE RTN PASS TAB CAB      LAX/LBX
0060          7452 017 100076          RTN      PASS CAB TAB      LAY/LBY
0061          7453 000 022461          MPCK INC M      M      STY
0062          7454 177 172036          WRTE RTN PASS TAB Y      CYA/CYB
0063          7455 017 101676          RTN      PASS Y      TAB      LDY
0064          7456 017 100157          PASS L      TAB      ADY
0065          7457 264 173676          ENVE RTN ADD Y      Y
0066          *****
0067          *      EADR IS COMMON TO LD*,ST*,AD*
0068          7460 220 074717          EADR      READ      INC PNM P      READ WORD 2 PC=ADDR OF NEXT INSTR.
0069          7461 301 165630          JSB          INDBIT      CHECK FOR INDIRECT, GET OPERAND
0070          7462 321 162035          JMP J3D          FINISH      JUMP TO COMPLETE INSTRUCTION
0071          *****
0072          *      EADRX DOES EFFECTIVE ADDR FOR
0073          *      SAX,SBX,LAX,LBX INSTRS.
0074          7463 220 074717          EADRX      READ      INC PNM P      READ ADDRESS OF WORD 2
0075          7464 017 170157          PASS L      X
0076          7465 017 100457          PASS M      TAB      M<=CONTENTS OF WORD 2.
0077          7466 322 023471          JMP CNDX AL15 RJS DIRECT      JUMP IF NO INDIRECT.

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0078 *****
0079 *
0080 *      INDIRECT ROUTINE FOR INDEXED INST
0081 7467 220 022457 EADRI  READ      INC  M      M      READ INDIRECT ADDRESS
0082 7470 301 165630 JSB      INDBIT   JSB TO INDIRECT ROUTINE
0083 *****
0084 *      COMPUTE INDEXED ADDRESS THEN JUMP
0085 7471 004 123017 DIRECT  READ      ADD  S1      M      S1<=TARGET ADDR. + X OR Y.
0086 7472 220 040457      INC  M      S1      READ INDEXED ADDRESS.
0087 7473 321 162035      JMP  J30      FINISH   JUMP TO COMPLETE THE INSTRUCTION
0088 *****
0089 *      EADRY COMPUTES EFFECTIVE ADDRESS
0090 *      FOR SAY, SBY, LAY, LBY INSTRS.
0091 7474 220 074717 EADRY  READ      INC  PHM  P
0092 7475 017 172157      PASS L  Y
0093 7476 017 100457      PASS M  TAB  M<= CONTENTS OF WORD 2.
0094 7477 322 023471      JMP  CNDX AL15 RJS DIRECT  JUMP IF NO INDIRECTS.
0095 7500 321 163370      JMP      EADRI   JUMP TO DO INDIRECT ROUTINE
0096 *****
0097 7501 017 103017 XABX      PASS S1  CAB  EXCHANGE A/B WITH X
0098 7502 017 170057      PASS CAB X
0099 7503 017 141636      RTN  PASS X  S1
0100 *****
0101 7504 017 103017 XABY      PASS S1  CAB  EXCHANGE A/B WITH Y
0102 7505 017 172057      PASS CAB Y
0103 7506 017 141676      RTN  PASS Y  S1
0104 *****
0105 7507 000 071617 ISX      INC  X      X      INCREMENT X, SKIP IF ZERO
0106 7510 320 067271      JMP  CNDX TBZ      SKIP
0107 7511 017 136776 RETURN      RTN
0108 *****
0109 7512 000 073657 ISY      INC  Y      Y      INCREMENT Y, SKIP IF ZERO.
0110 7513 320 067271      JMP  CNDX TBZ      SKIP
0111 7514 017 136776      RTN
0112 *****
0113 7515 007 171617 DSX      DEC  X      X      DECREMENT X, SKIP IF ZERO.
0114 7516 320 067271      JMP  CNDX TBZ      SKIP
0115 7517 017 136776      RTN
0116 *****
0117 7520 007 173657 DSY      DEC  Y      Y      DECREMENT Y, SKIP IF ZERO.
0118 7521 320 067271      JMP  CNDX TBZ      SKIP
0119 7522 017 136776      RTN

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```

0119 *****
0120 *      GENERAL INDIRECT ROUTINE FOR INDEX BIT INSTR
0121 *      COMMON ROUTINES FOR WORD/BYTE INSTRUCTIONS
0122 *****
0123 *      INITIALIZATION FOR WORD, BYTE
0124 7523 000 075217 INITCH  READ      INC  S5  P      S5<= ADDRESS OF WORD 3.
0125 7524 220 050457      INC  M      S5      READ ADDRESS OF WORD 3.
0126 7525 000 051157      INC  S4  S5      S4<= ADDRESS OF NEXT INSTRUCTION.
0127 7526 017 101117      PASS S3  TAB  S3<= CONTENTS OF WORD 3.
0128 7527 320 065531      JMP  CNDX TBZ      *+3      JUMP IF WORD 3 = 0 (NO INTERRUPT)
0129 7530 000 075717      INC  P      P      P<=ADDRESS OF WORD 3 (FOR EXIT)
0130 7531 001 155336      RTN  ZERO S7  S7      S7<=0 AND RETURN TO CALLER.
0131 7532 220 074710      READ STFL INC PHM P      READ ADDRESS OF WORD 2. P<=P+1.
0132 7533 001 155317      ZERO S7  S7      S7 <= 0.
0133 *****
0134 *      COMMON INDIRECT INBEDDED IN INITCH
0135 7534 017 100457 INDBIT      PASS M  TAB  M <= CONTENTS OF LAST READ ADDRESS.
0136 7535 322 026271      JMP  CNDX AL15 RJS CONTBIT  JUMP IF NO INDIRECT.
0137 7536 220 022465 INDLBIT  READ INCI INC  M      M      READ ADDRESS IN M
0138 7537 326 065631      JMP  CNDX NHO1  INDBIT  JUMP IF NO HALT OR INTERRUPT PENDING
0139 7540 017 100457 IN2BIT      PASS M  TAB  M<= CONTENTS OF LAST READ ADDRESS
0140 7541 330 125671      JMP  CNDX MSGC RJS INDBIT+1 JUMP IF SINGLE-INSTRUCT. MODE
0141 7542 007 175717 DEC2      DEC  P      P
0142 7543 007 175717      DEC  P      P      P <= ADDRESS OF WORD 1.
0143 7544 320 000030      JMP      FETCH      ATTEMPT JUMP TO FETCH ROUTINE.
0144 7545 324 066371 CONTBIT  JMP  CNDX FLAG      *+2      FLAG IDENTIFIES CALLER TO INDBIT
0145 7546 220 022476      READ RTN  INC  M      M      READ ADDRESS AND RETURN.
0146 *****
0147 7547 220 022451      READ CLFL INC  M      M      CALLER=INITCH--RESET FLAG, READ COUNT
0148 7550 017 101257      PASS S6  TAB  S6 <= COUNT FOR THIS INSTRUCTION
0149 7551 320 026571      JMP  CNDX TBZ  RJS  RTNCNT  JUMP IF COUNT NOT ZERO.
0150 7552 301 176730 JSB      EXIT      END THE INSTRUCTION.
0151 7553 017 153136 RTNCNT      RTN  PASS S3  S6      S3 <= COUNT, RETURN TO CALLER.

```

```

0152
0153
0154
0155 7554 220 074717
0156 7555 301 165630
0157 7556 017 100157
0158 7557 220 074457
0159 7560 301 165630
0160 7561 015 101017
0161 7562 013 041017
0162 7563 320 067271
0163 7564 000 075717
0164 7565 000 075736
0165
0166 7566 220 074717
0167 7567 301 165630
0168 7570 017 100157
0169 7571 220 074457
0170 7572 301 165630
0171 7573 327 170031
0172 7574 017 001017
0173 7575 000 022461
0174 7576 177 140017
0175 7577 000 075736
0176 7600 013 101017
0177 7601 000 022461
0178 7602 177 140017
0179 7603 000 075736
0180
0181
0182
0183 7604 301 165170
0184 7605 220 026457
0185 7606 017 100157
0186 7607 220 024457
0187 7610 000 024517
0188 7611 003 001017
0189 7612 320 036431
0190 7613 000 026557
0191 7614 007 145117
0192 7615 320 076731
0193 7616 335 030271
0194 7617 321 176130
0195
0196 7620 301 165170
0197 7621 220 026457
0198 7622 000 026557
0199 7623 017 101017
0200 7624 000 024457
0201 7625 017 122761
0202 7626 177 140017
0203 7627 000 024517
0204 7630 007 145117
0205 7631 320 076731
0206 7632 335 031071
0207 7633 321 176130

*****
* BIT INSTRUCTIONS
*****
TBS READ INC PMN P
JSB INDBIT GET MASK
PASS L TAB L <= MASK.
READ INC M P
JSB INDBIT GET WORD TO BE TESTED
AND S1 TAB LOGICAL AND OF MASK, WORD UNDER TEST.
XOR S1 S1 S1 <= 0 IF ALL MASK BITS SET IN WORD
JMP CNDX TBZ SKIP SKIP IF ALL MASK BITS SET IN WORD.
INC P P SKIP NEXT MACHINE INSTRUCTION.
RTN INC P P ADJUST P, JUMP TO FETCH ROUTINE.
*****
SBSCBS READ INC PMN P
JSB INDBIT OBTAIN BIT MASK
PASS L TAB L <= BIT MASK
READ INC M P
JSB INDBIT OBTAIN WORD TO BE OPERATED ON.
JMP CNDX IR2 CBS JUMP IF INSTRUCTION IS CBS.
IOR S1 TAB SET BITS IN WORD, PUT IN S1
MPCK INC M M MEMORY PROTECT CHECK.
WRTS PASS TAB S1 REWRITE WORD TO MEMORY. RETURN TO FETCH.
RTN INC P P
CBS SANL S1 TAB S1 <= MEMORY WORD WITH BITS CLEARED
MPCK INC M M
WRTS PASS TAB S1 REWRITE WORD TO MEMORY.
RTN INC P P RETURN TO FETCH ROUTINE.
*****
* WORD INSTRUCTIONS
*****
CMW JSB INITCH INITIALIZE
READ INC M A READ FROM ARRAY A.
PASS L TAB L <= WORD FROM ARRAY A
READ INC M B READ ADDRESS IN ARRAY B.
INC B B INCREMENT ARRAY B POINTER.
SUB S1 TAB SUBTRACT ARRAY WORDS B - A.
JMP CNDX TBZ RJS CHALIS JUMP IF UNEQUAL.
INC A A INCREMENT ARRAY A POINTER.
DEC S3 S3 DECREMENT COUNT.
JMP CNDX TBZ EXIT JUMP TO EXIT IF COUNT IS ZERO.
JMP CNDX INT RJS CMW+1 JUMP IF NOT INTERRUPTED.
JMP INTPEND
*****
MVW JSB INITCH INITIALIZE.
READ INC M A READ FROM ARRAY A
INC A A INCREMENT ARRAY A POINTER.
PASS S1 TAB S1 <= CONTENTS OF WORD OF ARRAY A
INC M B M <= ADDRESS FROM ARRAY B
MPCK PASS M MEMORY PROTECT CHECK-- BIT 15 LOW.
WRTS PASS TAB S1 WRITE WORD INTO ARRAY B.
INC B B ADVANCE ARRAY B POINTER.
DEC S3 S3 DECREMENT COUNT.
JMP CNDX TBZ EXIT EXIT IF COUNT IS ZERO.
JMP CNDX INT RJS MVW+1 JUMP IF NOT INTERRUPTED.
JMP INTPEND

0208
0209
0210
0211 7634 340 000157
0212 7635 015 127017
0213 7636 157 125244
0214 7637 220 052461
0215 7640 324 032171
0216 7641 013 101417
0217 7642 321 172330
0218 7643 015 101417
0219 7644 017 141093
0220 7645 017 141093
0221 7646 017 160157
0222 7647 017 041017
0223 7650 000 024517
0224 7651 177 140036

*****
* BYTE INSTRUCTIONS
*****
SBT IMM HIGH L X000 L <= 0003778.
STBYTE AND S1 A S1 <= RIGHT BYTE OF A REG.
LWF R1 PASS S6 B S6 <= WORD ADDRESS. FLAG SET IF BYTE ODD
READ MPCK INC M S6 READ WORD ADDRESS, CHECK FOR MP VIOLATION
JMP CNDX FLAG RJS STEVEN JUMP IF STORE TO EVEN BYTE.
SANL S9 TAB MASK OUT EVEN BYTE OF MEMORY WORD.
JMP MERGE
STEVEN AND S9 TAB MASK OUT ODD BYTE OF MEMORY WORD.
L4 PASS S1 S1 EXCHANGE BYTES IN REGISTER CONTAINING
L4 PASS S1 S1 BYTE TO BE STORED.
MERGE PASS L S9 L <= MEMORY WORD WITH TARGET BYTE CLEARED
IOR S1 S1 S1 <= WORD WITH BYTES MERGED.
INC B B INCREMENT BYTE ADDRESS.
WRTS RTN PASS TAB S1 WRITE NEW WORD BACK INTO WORD ADDRESS.

```


Appendix E

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0225 *****
0226 7652 017 125057 LBT PASS S2 B S2 <= BYTE ADDRESS.
0227 7653 000 024517 INC B B INCREMENT BYTE ADDRESS FOR NEXT INSTRUCT.
0228 7654 340 000151 LDBYTE IMM CLFL HIGH L %000 L <= 000377B. CLEAR CPU FLAG.
0229 7655 157 143244 LWF R1 PASS S6 S2 S6 <= WORD ADDRESS OF BYTE. SET FLAG IF 0
0230 7656 220 052457 READ INC M S6 ODD BYTE. READ WORD ADDRESS.
0231 7657 324 073171 JMP CNDX FLAG LODD JUMP IF BYTE IS ODD.
0232 7660 013 100543 L4 SANL A TAB MASK OUT EVEN BYTE AND MOVE ODD BYTE
0233 7661 017 126543 L4 PASS A A TO EVEN BYTE OF A REG.
0234 7662 017 136776 RTN AND A TAB RETURN TO CALLER OR FETCH.
0235 7663 015 100776 LODD RTN AND A TAB MASK OUT EVEN BYTE, LOAD INTO A. RETURN.
0236 *****
0237 7664 301 165170 MBT JSB INITCM INITIALIZE.
0238 7665 017 127057 PASS S2 A S2 <= ADDRESS START OF ARRAY A.
0239 7666 301 172630 JSB CLFL INC S2 LDBYTE LOAD BYTE FROM ARRAY A.
0240 7667 000 043051 JSB STBYTE S2 RESET FLAG, S2 <= NEXT BYTE ADDR IN ARRAY
0241 7670 301 171670 DEC S3 S3 STORE BYTE INTO BYTE ADDRESS IN B REG.
0242 7671 007 145117 JMP CNDX TBZ RJS S3 DECREMENT COUNT
0243 7672 320 033671 PASS A S2 JUMP IF COUNT NOT ZERO.
0244 7673 017 142557 JMP CNDX INT RJS S2 A <= 1 + LAST ARRAY A BYTE ADDRESS MOVED
0245 7674 321 176730 JMP PASS A S2 JUMP IF NOT INTERRUPTED.
0246 7675 335 033331 NOTDAM S2 A <= 1 + LAST ARRAY A ADDRESS MOVED
0247 7676 017 142557 JMP INTPEND
0248 7677 321 176130

0249 *****
0250 7700 340 000157 SFB IMM HIGH L %000 L <= 000377B.
0251 7701 015 127117 AND S3 A S3 <= TEST BYTE IN LOW-ORDER BYTE
0252 7702 013 127143 L4 SANL S4 A S4 <= TERMINATION BYTE IN
0253 7703 017 147143 L4 PASS S4 S4 LOW-ORDER BYTE
0254 7704 017 127357 PASS S8 A S8 <= SAVE ORIGINAL CONTENTS OF A REG.
0255 7705 301 172530 CONTSFB JSB LBT LOAD BYTE INTO A REG FROM ADDRESS IN B.
0256 7706 017 126157 PASS L A L <= BYTE TO TESTED IN LOW BYTE.
0257 7707 013 045257 XOR S6 S3 COMPARE BYTE TO TEST BYTE.
0258 7710 320 034571 JMP CNDX TBZ RJS NOMATCH JUMP IF UNEQUAL.
0259 7711 017 156557 PASS A S8 RESTORE ORIGINAL CONTENTS OF A REG
0260 7712 007 124536 RTN DEC B B B <= BYTE ADDRESS OF MATCH. GO TO FETCH
0261 7713 013 047017 NOMATCH XOR S1 S4 COMPARE BYTE TO TERMINATION BYTE.
0262 7714 320 034771 JMP CNDX TBZ RJS INTTST JUMP IF UNEQUAL.
0263 7715 017 156557 PASS A S8 RESTORE ORIGINAL CONTENTS OF A REG.
0264 7716 000 075736 RTN INC P P SKIP NEXT MACHINE INSTRUCTION AND FETCH.
0265 7717 335 034271 INTTST JMP CNDX INT RJS CONTSFB JUMP IF NOT INTERRUPTED.
0266 7720 017 156557 PASS A S8 RESTORE ORIGINAL CONTENTS OF A REG.
0267 7721 007 175736 RTN DEC P P P <= INSTRUCTION ADDRESS, GO TO FETCH.
0268 *****
0269 7722 301 165170 CBT JSB INITCM INITIALIZE.
0270 7723 017 127357 PASS S8 A S8 <= POINTER FOR ARRAY A.
0271 7724 017 157057 PASS S2 S8 S2 <= NEXT BYTE ADDRESS IN ARRAY A.
0272 7725 301 172630 JSB LDBYTE LOAD ARRAY A BYTE INTO A REG.
0273 7726 017 127417 PASS S9 A S9 <= ARRAY A BYTE.
0274 7727 000 043357 INC S8 S2 INCREMENT ARRAY A POINTER.
0275 7730 301 172530 JSB LBT A <= BYTE FROM ARRAY B (ADDRESS IN B REG)
0276 7731 017 160157 PASS L S9 L <= BYTE FROM ARRAY A
0277 7732 003 027017 SUB S1 A A SUBTRACT BYTE FROM ARRAY B - A.
0278 7733 320 036331 JMP CNDX TBZ RJS CHALISE JUMP IF BYTES NOT EQUAL.
0279 7734 007 145117 DEC S3 S3 DECREMENT THE COUNT.
0280 7735 320 036031 JMP CNDX TBZ RJS CONTCBT JUMP IF COUNT IS NOT ZERO.
0281 7736 017 156557 PASS A S8 EQUAL EXIT... A <= 1+LAST MOVED BYTE ADDR
0282 7737 321 176730 JMP EXIT
0283 7740 335 035231 CONTCBT JMP CNDX INT RJS CBT+2 JUMP IF NOT INTERRUPTED.
0284 7741 017 156557 PASS A S8 A <= NEXT BYTE ADDRESS OF ARRAY A TO TEST

```

```

0285
0286
0287
0288
0289 7742 000 050457
0290 7743 177 144017
0291 7744 007 175717
0292 7745 007 175736
0293
0294
0295 7746 007 156557
0296 7747 017 141017
0297 7750 322 036531
0298 7751 000 047157
0299 7752 000 047157
0300 7753 007 145117
0301 7754 017 144157
0302 7755 004 124517
0303
0304
0305 7756 000 050457
0306 7757 177 154017
0307 7760 017 147736
0308
0309
0310
0311 7761 220 074717
0312 7762 301 165630
0313 7763 340 120417
0314 7764 017 122461
0315 7765 017 175657
0316 7766 017 123736
0317
0318 7767 220 074717
0319 7770 017 172157
0320 7771 004 101017
0321 7772 340 120417
0322 7773 017 140457
0323 7774 017 122761
0324 7775 017 123736
0325
0326 7776 377 177777
0327 7777 377 177777
0328
** NO ERRORS**

*****
*      COMMON ROUTINES TO MOVE, COMPARE INSTRUCTIONS
*****
*      INTERRUPT EXIT
INTPEND      INC M      S5      M <= ADDRESS OF WORD 3
             WRT      PASS TAB S3      WRITE REMAINING COUNT INTO WORD 3.
             DEC P      P
             RTN      DEC P      P      P <= ADDRESS OF WORD 1, GO TO FETCH.
*****
*      EXIT TESTS FOR CBT, CMW
CHAL158      DEC A      S8      A <= BYTE ADDRESS OF MISMATCH.
             PASS S1      S1      CHECK RESULT OF COMPARE
CHAL15      JMP CNDX AL15 RJS SKIP1      JUMP IF SIGN BIT IS ZERO.
             INC S4      S4      SKIP ONE MACHINE INSTRUCTION.
SKIP1      INC S4      S4      SKIP ONE MACHINE INSTRUCTION.
             DEC S3      S3      DECREMENT THE COUNT.
             PASS L      S3      L <= COUNT REMAINING
             ADD B      B      B <= FIRST ADDR. IN ARRAY B + COUNT.
*****
*      COMPLETION EXIT
EXIT      INC M      S5      M <= ADDRESS OF WORD 3
             WRT      PASS TAB S7      WORD 3 <= ZERO.
             RTN      PASS P      S4      P <= NEXT MACHINE INSTR. TO EXECUTE. FETCH
*****
*      JUMP INSTRUCTIONS
*****
JLY      READ      INC PHM P      READ ADDRESS OF WORD 2.
             JSB      INDBIT      CHECK FOR INDIRECT, M<= DESTINATION ADDR.
             IMM      HIGH IR      X050      MACHINE JNP INTO IR TO SET LOW MP BOUNDS
             MPCK PASS M      M      DO MP CHECK ON JUMP TARGET ADDRESS.
             PASS Y      P      Y <= ADDRESS OF FOLLOWING MACHINE INSTR.
             RTN      PASS P      M      P <= DESTINATION ADDRESS, JUMP TO FETCH.
*****
JPY      READ      INC PHM P      READ ADDRESS OF WORD 2.
             PASS L      Y      L <= INDEX REG. Y.
             ADD S1      TAB      S1 <= INDEXED JUMP ADDRESS.
             IMM      HIGH IR      X050      MACHINE JNP INTO IR TO SET LOW MP BOUNDS
             PASS M      S1      M<= INDEXED ADDRESS, WITH BIT 15 LOW
             MPCK PASS M      M      MP CHECK ON 15-BIT DESTINATION ADDRESS.
             RTN      PASS P      M      P <= DESTINATION ADDRESS, GO TO FETCH.
*****
ONES
ONES
$END

```


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